

Introduction to Scientific Computing

(Lecture 11: Overview)

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What we have learned?



CONDITION NUMBER



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Condition number

This number predicts the relative error in output y if the relative error in input x is known. Or better to say represents sensitivity of the output on perturbations of input data.

$$k(x) := \frac{\partial f}{\partial x}(\mathbf{x}) \frac{x}{f(\mathbf{x})}$$
$$\epsilon_y = \frac{\Delta y}{y} = k(x)\epsilon_x = k(x)\frac{\Delta x}{x}$$

Example: substruction is ill-conditioned operation, e.g.

$$\frac{1}{x_1-x_2}, \quad x_1 \approx x_2$$

DIFFERENCE AND DIFFERENTIAL EQUATIONS



Difference vs. Differential equations

Discrete system	Continuous system
Difference	Differential
$x_{n+1} - x_n = 2h^2n$	$\frac{dx}{dt} = 2t$
	Discretise in time
	$x_{n+1} - x_n = 2h^2n$

Differential equations become difference when we apply numerical integration (time discretisation)!

Homogeneous equations (zero rhs)

Difference	Differential
$x_{n+1}-x_n=0$	$\frac{dx}{dt} = 0$
Solution form	Solution form
$x_n = C \rho^n$	$x(t) = Ce^{ ho t}$
Characteristic	Characteristic
ho - 1 = 0	ho = 0
$x_n = C(1)^n = C$	$x = Ce^{0t} = C$

Homogeneous solution wrt roots

Difference	Differential	
Single roots ρ_i	Single roots $ ho_i$	
$x_n = \sum_i C_i \rho_i^n$	$x(t) = \sum_{i} C_{i} e^{ ho_{i} t}$	
Double roots ρ_1	Double roots $ ho_1$	
$x_n = C_1 \rho_1^n + n C_2 \rho_1^n$	$x(t) = C_1 e^{\rho_1 t} + t C_2 e^{\rho_1 t}$	
Roots and eigenvalues are equal!!		

Homogeneous solution wrt roots

Differential
$Complex\ \rho = \alpha \pm \beta i$
$e^{(lpha\pmeta i)t}=e^{lpha t}e^{eta it}$
$x = e^{\alpha t} (A \cos(\beta t) + iB \sin(\beta t))$

Т

Note that in one case is the modulus and in another real part of the root.

Non-homogeneous- Particular solution



STABILITY



Stability

Difference	Differential
$x_{n+1} = F(x_n)$	$\frac{dx}{dt} = f(x)$
$x_* = F(x_*)$	$f(x_*)=0$
roots of charact. equation or	roots of charact. equation or
$[\lambda, v] = \mathrm{eig} \ DF(x_*)$	$[\lambda, v] = \mathrm{eig} \ Df(x_*)$
$(\lambda)^n \Rightarrow \mathrm{abs}\; (\lambda) \leq 1 \; (real)$	$e^{\lambda t} \Rightarrow (\lambda) \leq 0$ (real)
$\mathrm{abs}\;(\mu)\leq 1\;(complex)$	$(\mathit{Re}(\lambda)) \leq 0 \; (complex)$

ī.

If there is just one root equal to 1 (i.e. 0) \Rightarrow stable but not asympt.

$$(1)^n = 1, \quad e^{0t} = 1, \quad n \to \infty \quad \& \quad t \to \infty$$

If there are more roots equal to 1(i.e. 0) \Rightarrow not stable

$$c_1(1)^n + c_2 n(1)^n \to \infty, \quad c_1 e^{0t} + c_2 t e^{0t} \to \infty, \quad n \to \infty \quad \& \quad t \to \infty$$

Stability with respect to the intial conditions. Apply numerical method to the ODE

$$\dot{x} = 0$$

such as for example Euler:

$$x_{n+1}-x_n=0$$

then test stability of this difference equation.

Zero stability is important due to the process of selecting initial conditions. They are not always correct.

COMPUTING THE ROOTS—SOLVING SYSTEMS OF EQUATIONS

F(X) = 0



Solving linear system of equations

System:

$$Ax = b$$

Gauss-Seidel method:

$$\mathsf{L}\mathsf{x}^{(k+1)} = \mathsf{b} - \mathsf{U}\mathsf{x}^{(k)}$$

Jacobi method:

$$\mathbf{D}\mathbf{x}^{(k+1)} = \mathbf{b} - \mathbf{R}\mathbf{x}^{(k)}$$

Successive over-relaxation (SOR) method:

$$\mathbf{x}^{(k+1)} = (\mathbf{D} + \omega \mathbf{L})^{-1} (\omega \mathbf{b} - [\omega \mathbf{U} + (\omega - 1)\mathbf{D}]\mathbf{x}^{(k)})$$

Fixed-point iterations

We want to solve

$$F(x) = x$$

such that

- **F** has fixed point (solution)
- the fixed point is unique
- can be obtained by iterative process
- **F** is also called mapping.

Fixed-point iterations

Banach fixed point

Lipshitz continuity (for $q \ge 0$):

$$\|oldsymbol{F}(\mathbf{x}) - oldsymbol{F}(\mathbf{y})\| \leq q\|\mathbf{x} - \mathbf{y}\|$$

If this holds when $0 \leq q < 1$ then

- $F(\mathbf{x}_*) = \mathbf{x}_*$ has unique solution
- for any initial value sequence x_{n+1} = F(x_n) converges to solution x_{*} (this means that iterative method is convergent and gives the solution)
- speed of convergence are given by apriori and aposteriori estimates

We have studied two error estimates

• A posteriori error estimate:

$$\|\mathbf{x}_* - \mathbf{x}_n\| \leq \frac{q}{1-q} \|\mathbf{x}_n - \mathbf{x}_{n-1}\|$$

• A priori error estimate:

$$\|\mathbf{x}_* - \mathbf{x}_n\| \leq rac{q^n}{1-q} \|\mathbf{x}_1 - \mathbf{x}_0\|.$$

Newton method

$$F(x) = 0$$

 $F(x) = F(x_0) + F'(x_0)(x - x_0) + h.o.t. = 0$

$$\mathsf{J}(\mathsf{x})=\mathsf{F}'(\mathsf{x})$$

with elements:

Jacobian:

$$J_{ij} = \frac{\partial f_i}{\partial x_j}$$

Newton iteration:

$$\mathbf{x}_k = \mathbf{x}_{k-1} - \frac{\mathbf{F}(\mathbf{x}_{k-1})}{\mathbf{J}(x_{k-1})}$$

NUMERICAL INTEGRATION

$$\int_a^b f(t)dt = \dots$$

To compute the solution of integral equation one may use the fixed point iteration (also called Picard-Lindelöf iteration):

$$x^{(k+1)}(t) = x_0 + \int_{t_0}^t f(x^{(k)}(s), s) ds = F(x^{(k)}, s)$$

by starting from $x^{(0)} = x_0$.

Newton-Cotes integration

Approximate function (rhs of ODE) by a polynomial

$$\int_{t_0}^T f(t) dt \approx \int_{t_0}^T P(t) dt = \int_{t_0}^T \sum_{j=0}^n a_j t_i^j dt$$

Possible polynomial approximations (global or local):

- piecewise constant interpolation (rectangle rule)
- linear interpolation (trapezoidal rule)
- quadtratic interpolation (Simpsons rule)

In Newton-Cotes formula one chooses the points t_i in which the value of function will be evaluated. This is not optimal as the result may lead to large errors. The Gauss quadrature has for a goal to vary the placements t_i such that the integration is more accurate. In general Gauss formula approximates:

$$\int_{-1}^{1} g(t) dt \approx \int_{-1}^{1} P_n(t) dt = \sum_{i=0}^{n} g(t_i) w_i, \quad w_i = \int_{-1}^{1} \ell_i(t) dt$$

in the same way as Newton-Cotes. But, t_i are unknown and have to be found. Additionally to them are also unknown the coefficients of a polynomial $P_n(x)$.

NUMERICAL SCHEMES TO SOLVE ODE

$$\frac{dx}{dt} = f(x,t) \Rightarrow x(t_{n+1}) = \dots$$



Multistep methods

Numerical methods for solving ODE

$$\dot{x}=f(x,t)$$

can be classified as

• **one step methods** : these methods use only information from one time point to compute the next

 $x_{n+1} = x_n + hf(t_n, x_n)$ requires knowledge on x_n

 multistep methods: these methods require knowledge on more than one time point

$$x_{n+1} = x_n + h\left(rac{3}{2}f_n - rac{1}{2}f_{n-1}
ight)$$
 requires knowledge on x_n, x_{n-1}

Linear multistep methods of the form (only linear combination of x's and f's)

$$\sum_{l=0}^{k} a_{l} x_{n+l} = h \sum_{l=0}^{k} b_{l} f(t_{n+l}, x_{n+l})$$

Constraining ourselves to one step, one obtains

- explicit Euler method $x_{n+1} = x_n + hf(t_n, x_n)$
- implicit Euler method $x_{n+1} = x_n + hf(t_{n+1}, x_{n+1})$
- trapezoidal rule $x_{n+1} = x_n + \frac{h}{2}(f(t_n, x_n) + f(t_{n+1}, x_{n+1}))$
- explicit midpoint rule $x_{n+1} = x_n + hf(t_n + \frac{h}{2}, x(t_n) + \frac{h}{2}f(t_n, x_n))$
- implicit midpoint rule $x_{n+1} = x_n + hf(t_n + \frac{h}{2}, \frac{1}{2}(x_n + x_{n+1}))$

ACCURACY OF NUMERICAL SCHEMES

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Many students actually look forward to Mr. Atwadder's math tests.

Consistency of numerical methods for ODE

The method is said to be consistent if the local (truncation) error converges to zero when the step-size becomes arbitrarily small, i.e.

$$\lim_{h\to 0} T_h = \lim_{h\to 0} \frac{\epsilon_{loc}}{h} = 0$$

The scheme is consistent of order p if

$$\max \|\frac{\epsilon_{loc}}{h}\| \le Ch^p$$

The local error ϵ_{loc} is the difference between the exact $x_a(t_{n+1})$ and numerical $x(t_{n+1})$ solutions at t_{n+1}

$$\epsilon_{loc} = x_a(t_{n+1}) - x(t_{n+1})$$

where $x(t_{n+1})$ is computed by a numerical scheme starting from $x_a(t_n)$.

Convergence

Linear multistep method is said to be convergent if

$$\max_{t_n\in T} \|\epsilon_{glob}(t_n,h)\| \to 0, \quad h \to 0$$

and it has the order of convergence q if

$$\max_{t_n \in T} \|\epsilon_{glob}(t_n, h)\| \leq Ch^q$$

where

• the global error is

$$\epsilon_{glob} = x_a(t) - x(t)$$

- $x_a(t)$ is the exact solution of the differential equation $\dot{x} = f(t, x)$
- x(t) is the approximate solution
- C is a constant independent of h



Convergence implies consistency. Consistency does not imply convergence.

STABILITY OF NUMERICAL SCHEMES



Global error is defined as:

$$\epsilon_{glob}^{n+1} = x_a(t_{n+1}) - x(t_{n+1})$$

in which x_a is analytical and x(t) full numerical solution. Having in mind that

$$x(t_{n+1}) = x(t_n) + hf(t_n, x_n) = x(t_n) + h\lambda x_n, \quad n = 0, 1, 2, ...$$

with λ being -1 or -100, one may write

$$\epsilon_{glob}^{n+1} = x_a(t_{n+1}) - x(t_n) - h\lambda x_n$$

Absolute Stability

To judge if the numerical method produces desired numerical accuracy for given h > 0, one introduces the notion of absolute stability. Absolute stability is defined as a stability of numerical scheme applied on the Dahlquist problem:

$$\dot{x} = \lambda x, \quad x(0) = 1$$

whose exact solution is given by

$$x = \exp(\lambda t)$$

and depends on the value of λ

$$\lim_{t \to \infty} |x(t)| = \begin{cases} 0, & \text{if } \lambda < 0 \Rightarrow \text{stable ODE} \\ 1, & \text{if } \lambda = 0 \Rightarrow \text{not considered} \\ \infty, & \text{if } \lambda > 0 \Rightarrow \text{not considered} \end{cases}$$

Absolute Stability

The difference equation $\sum_{j=0}^{k} a_j x_{n+j} = h \sum_{j=0}^{k} b_j f(t_{n+j}, x_{n+j})$ is stable if roots

$$\rho(\xi) - z\sigma(\xi) = 0$$

are smaller by amplitude than 1, or eventually only one of them is equal to 1. With respect to this one defines

Definition

The region of absolute stability for the LMM is the set of points $z = h\lambda$ in the complex plane for which the polynomial $\rho(\xi) - z\sigma(\xi) = 0$ satisfies the root condition.

$$\mathit{G_s}:=\{z\in\mathbb{C}\;:\;|\xi(z)|\leq 1\}$$

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Thats all folks

"Any fool can know. The point is to understand." Albert Einstein



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