

Decompiling for Constant-Time Analysis

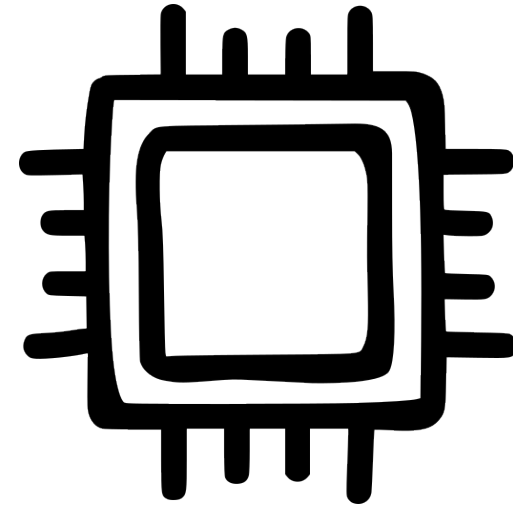
Decompilers disturb Formal Side Channel Security!

**Santiago Arranz Olmos, Gilles Barthe, Lionel Blatter,
Youcef Bouzid, Sören van der Wall, Zhiyuan Zhang**

**PriSC 2026
Rennes**

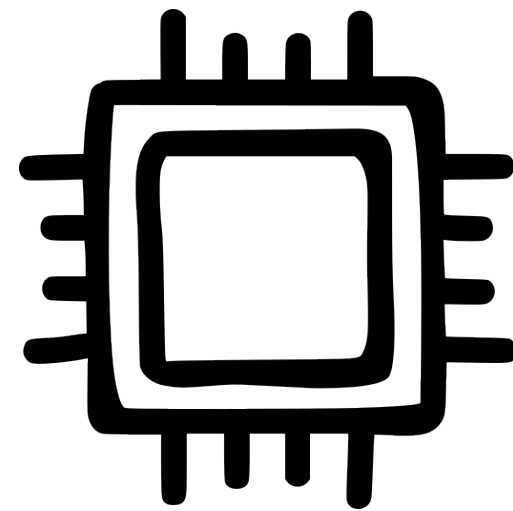
Formal Side-Channel Security 101

Attacker Model



Formal Side-Channel Security 101

Attacker Model



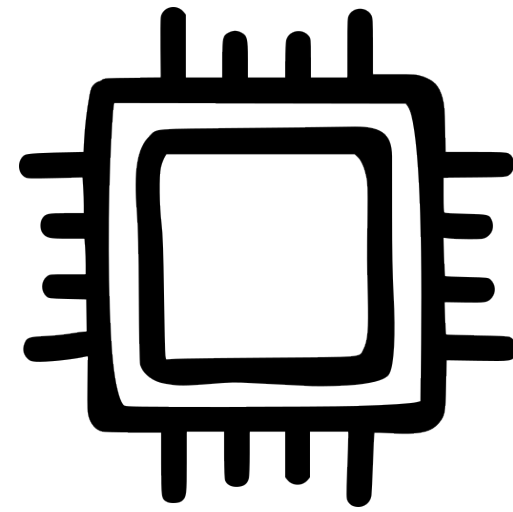
mem

crypto.o

attacker.o

Formal Side-Channel Security 101

Attacker Model



mem

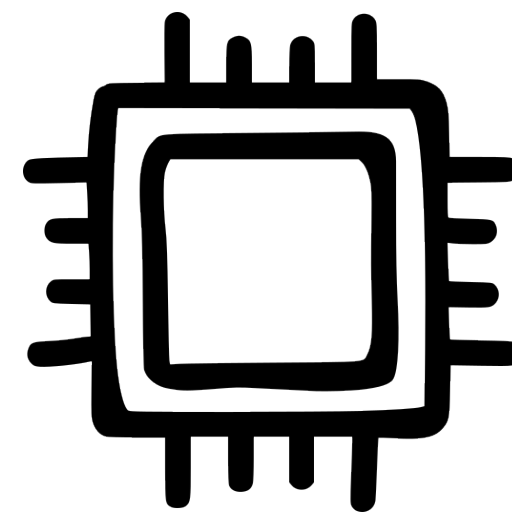
crypto.o

os

attacker.o

Formal Side-Channel Security 101

Attacker Model



cache

branch predictor

mem

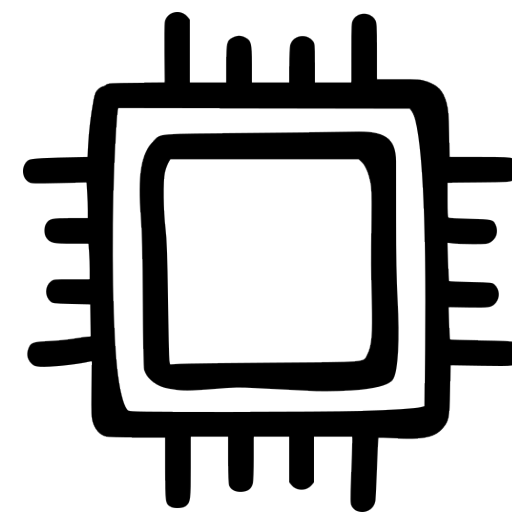
crypto.o

os

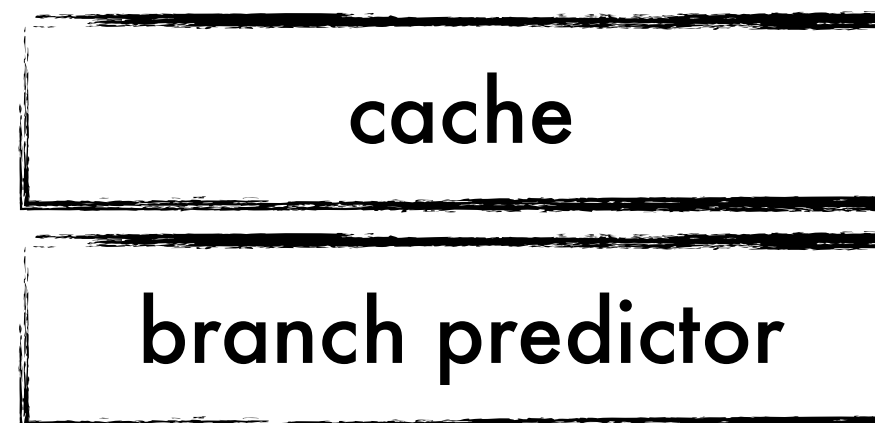
attacker.o

Formal Side-Channel Security 101

Attacker Model



leak



mem

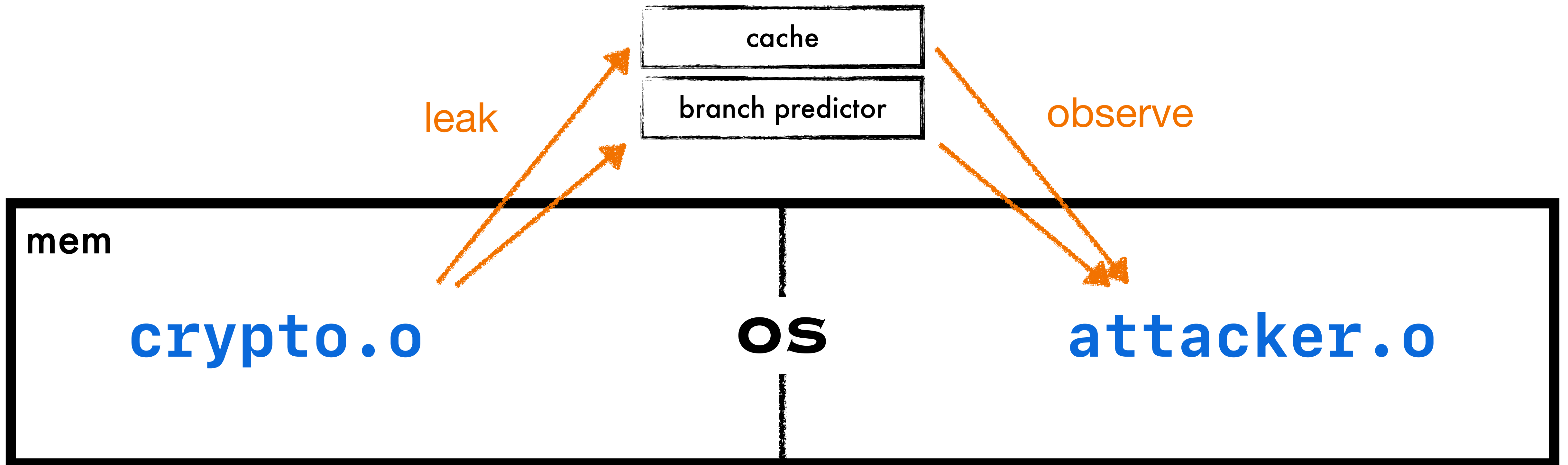
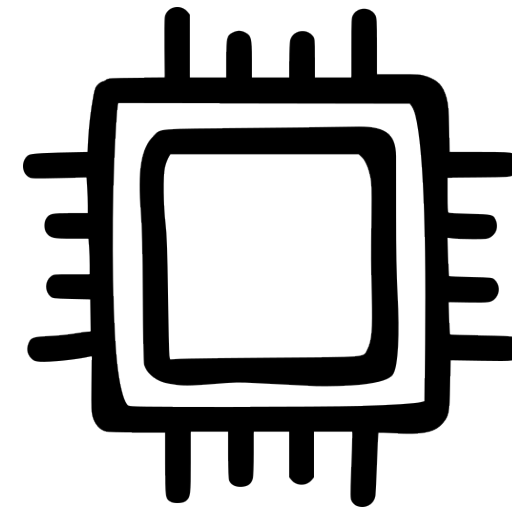
crypto.o

os

attacker.o

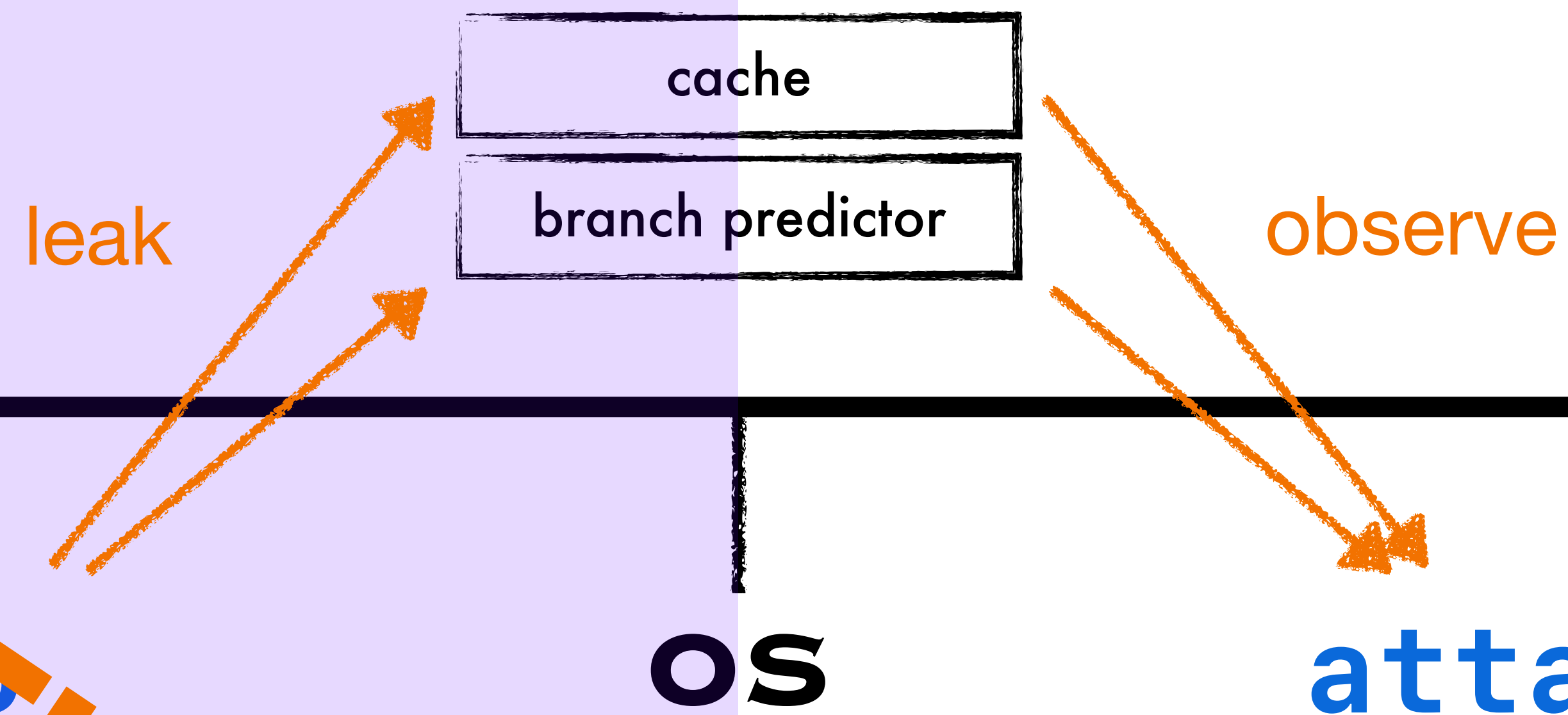
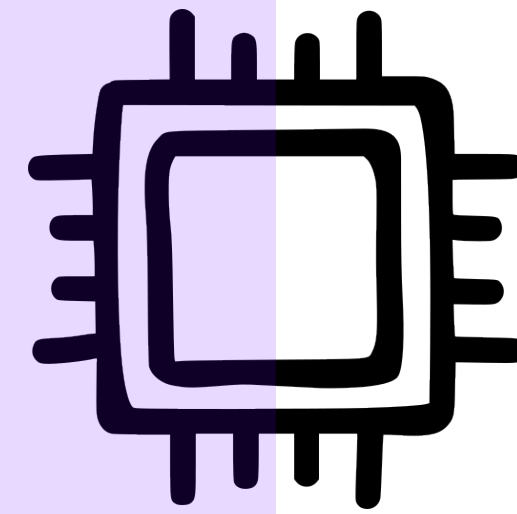
Formal Side-Channel Security 101

Attacker Model



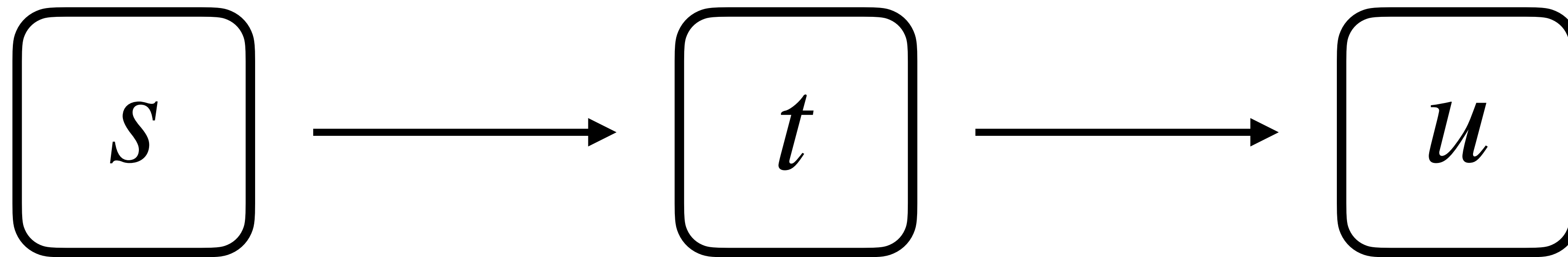
Formal Side-Channel Security 101

Attacker Model

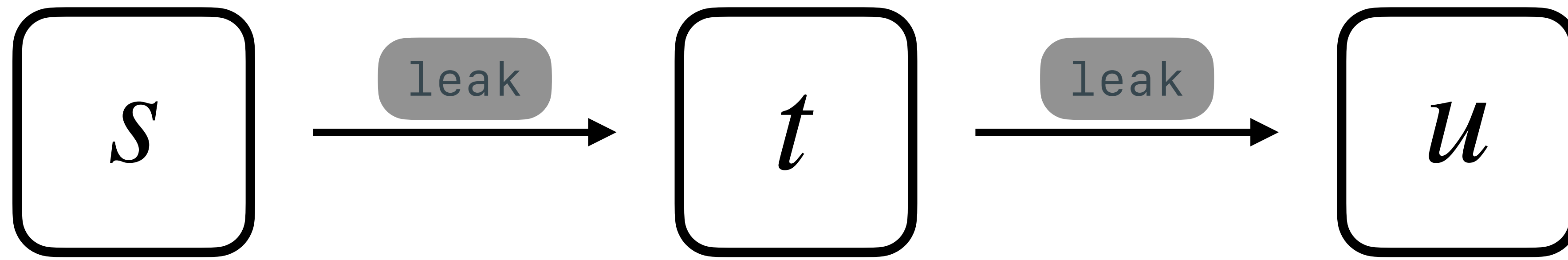


Model This!

Model: CT-Leakage Semantics



operational semantics



operational semantics

+

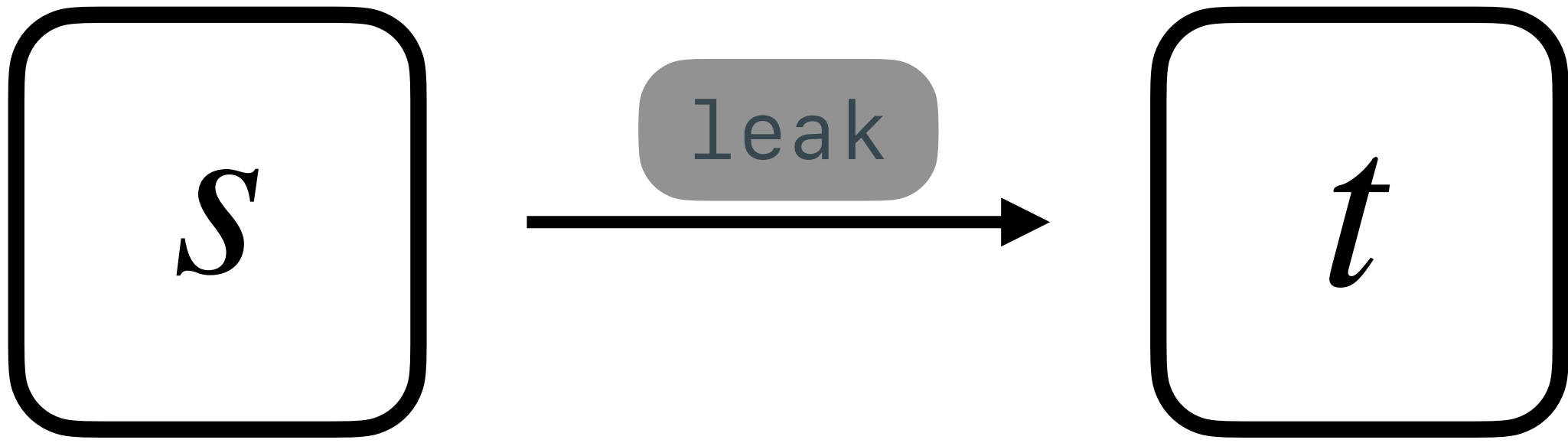
leakage observations

CT-Leakage



leak depends on
executed instruction

CT-Leakage

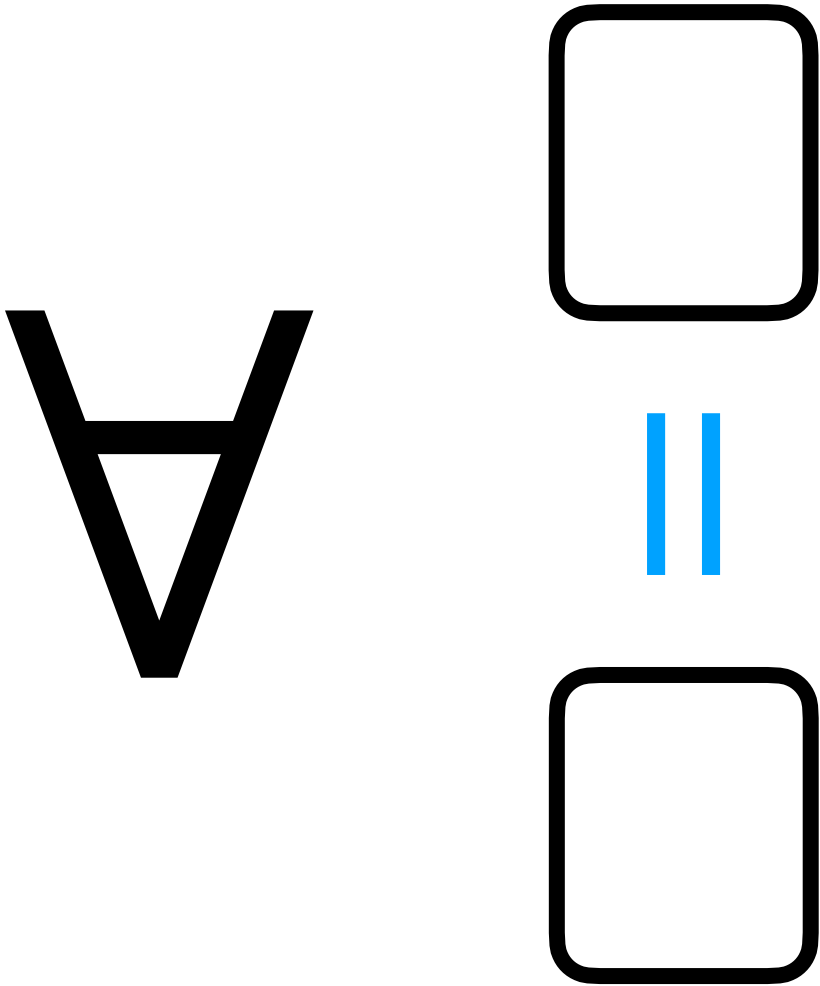


leak depends on
executed instruction

instruction	leak
load / store	address
conditional branch	branch outcome

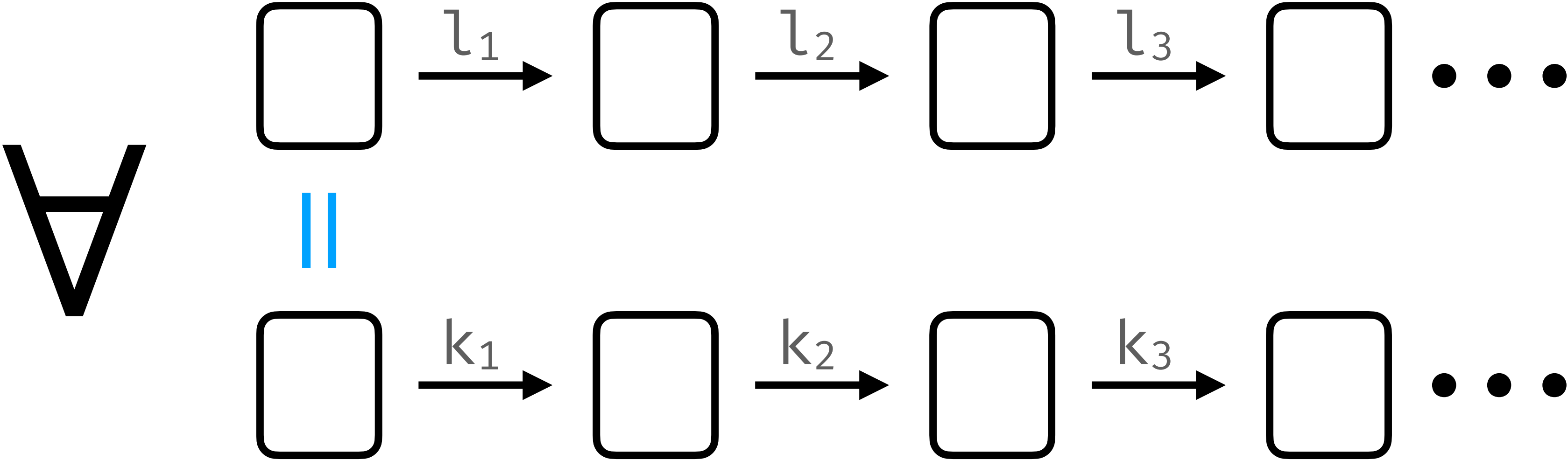
prog.c is **CT** **WHEN**

prog.c is CT WHEN



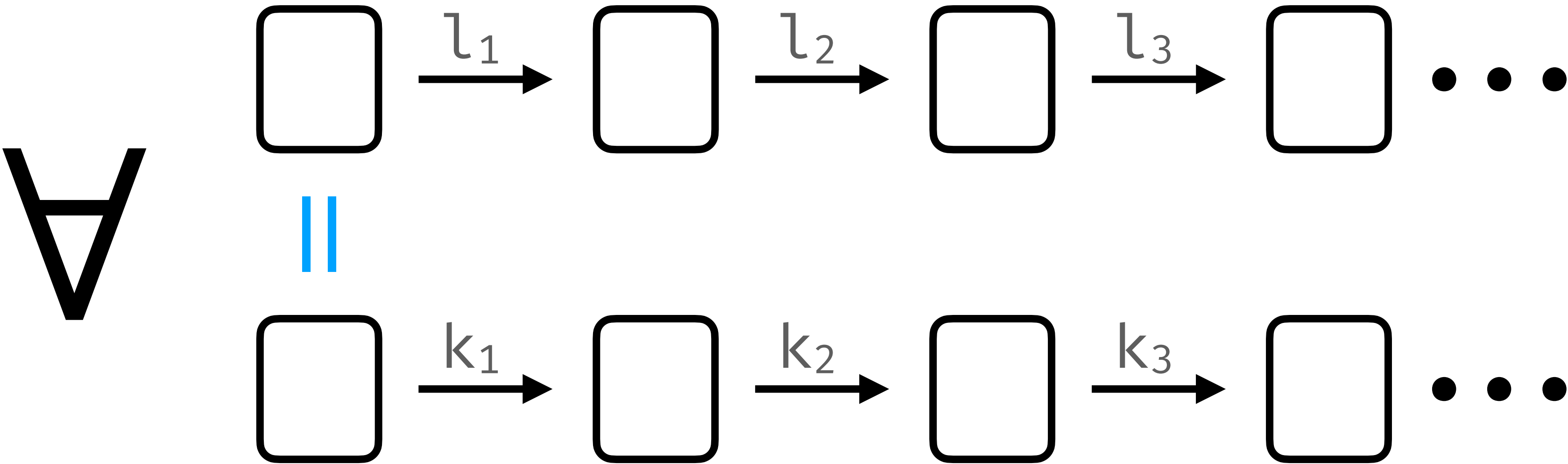
= attacker-indistinguishable:
equal except on secrets.

prog.c is **CT** WHEN



$=$ attacker-indistinguishable:
equal except on secrets.

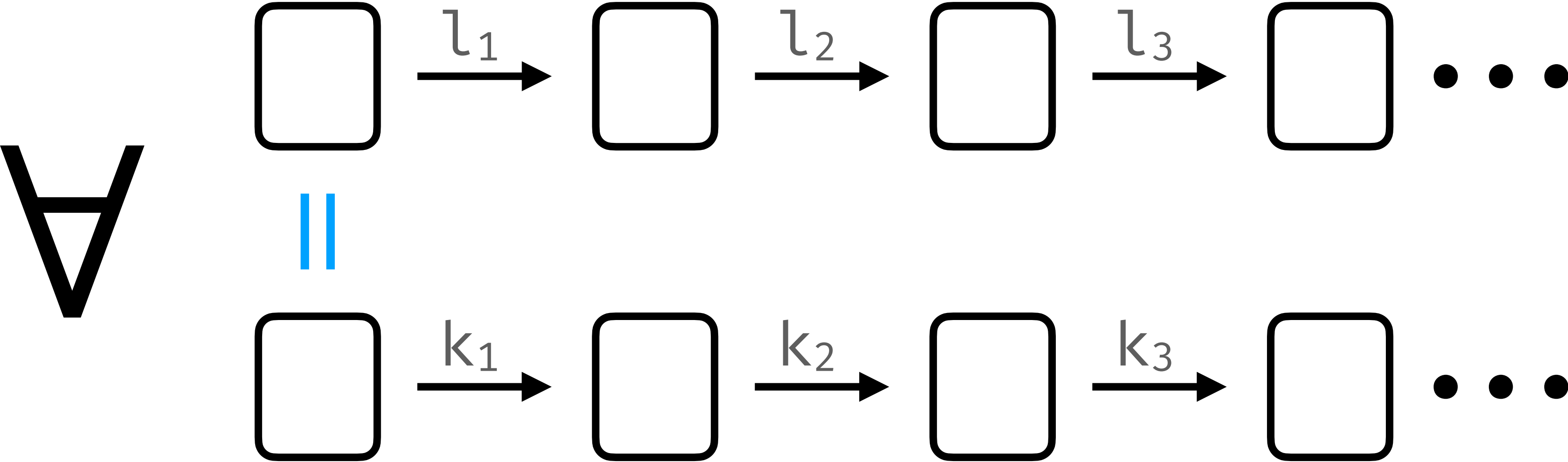
prog.c is **CT** **WHEN**



have $k_1 k_2 k_3 \dots = l_1 l_2 l_3 \dots$

$\equiv$ attacker-indistinguishable:
equal except on secrets.

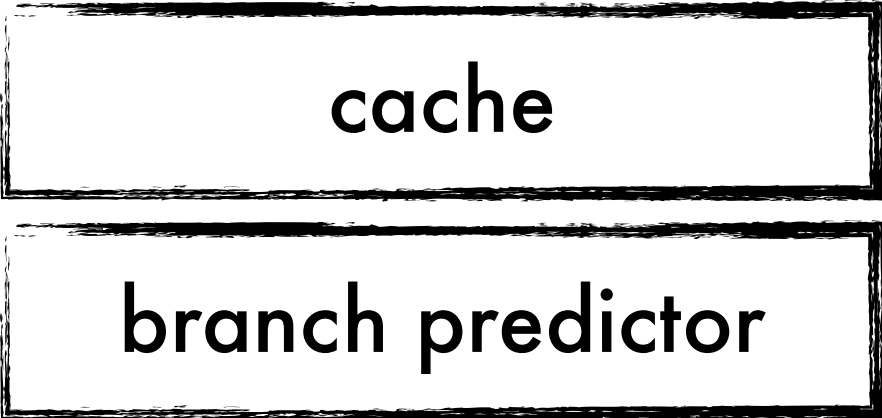
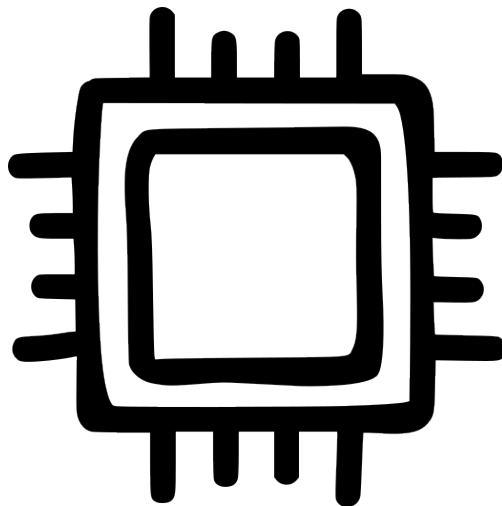
prog.c is **CT** **WHEN** “Leakages don’t depend on secrets.”



have $k_1k_2k_3... = l_1l_2l_3...$

= attacker-indistinguishable:
equal except on secrets.

Formal Side Channel Security 101

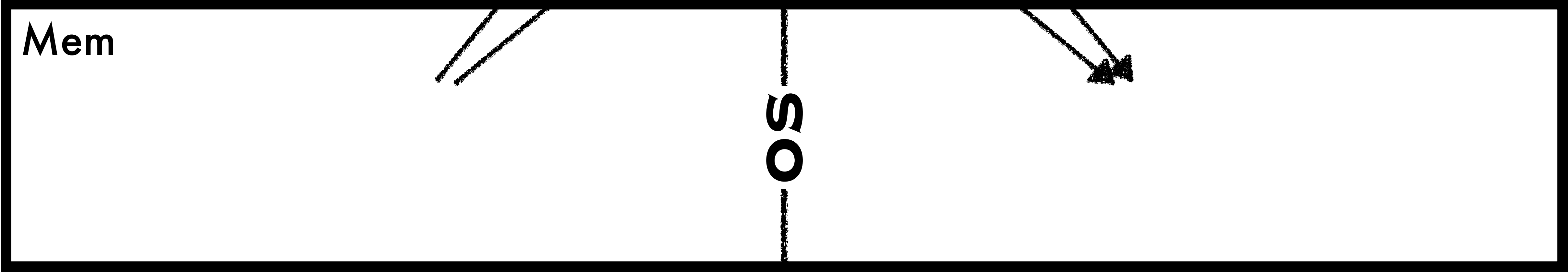


leak

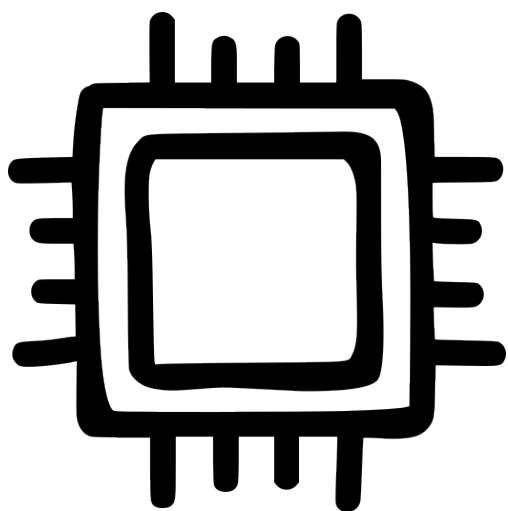
observe

Mem

OS

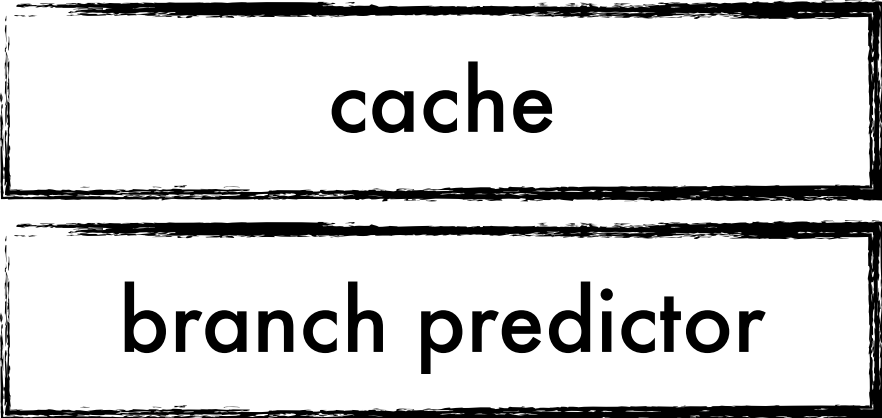


Formal Side Channel Security 101



leak

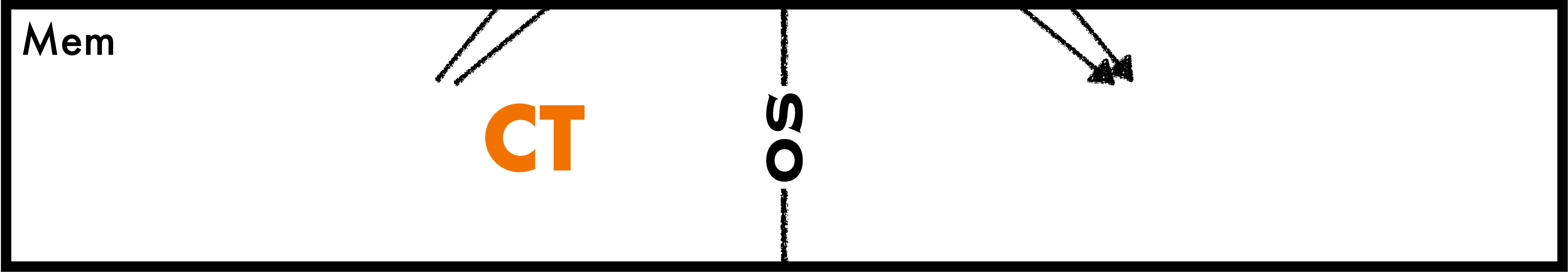
observe



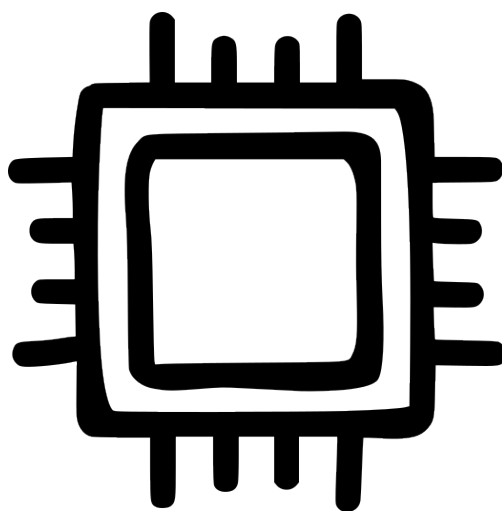
Mem

CT

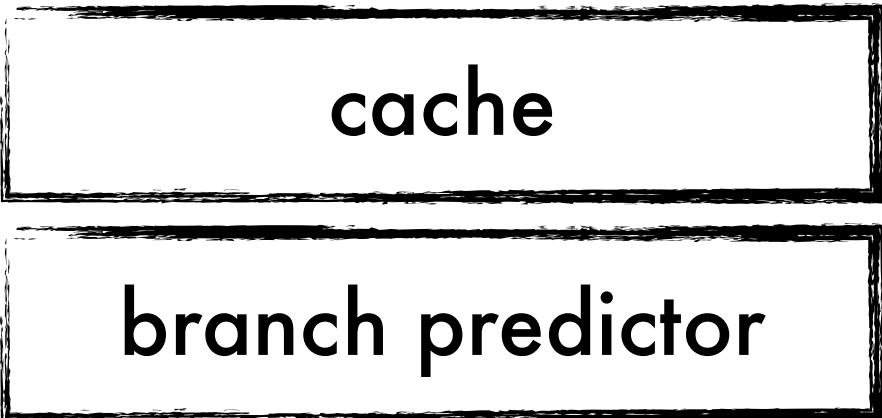
OS



Formal Side Channel Security 101



harmless
leak

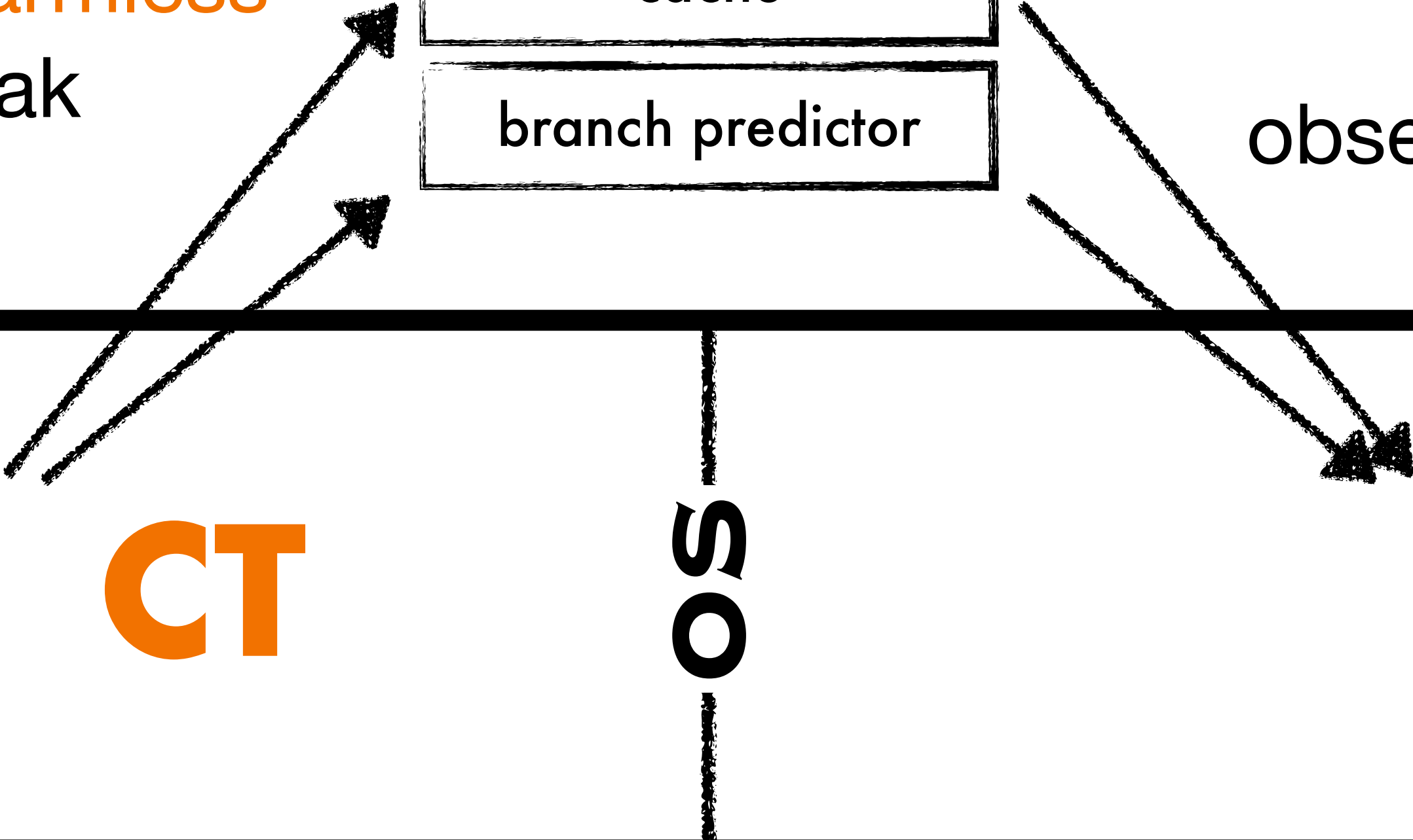


observe

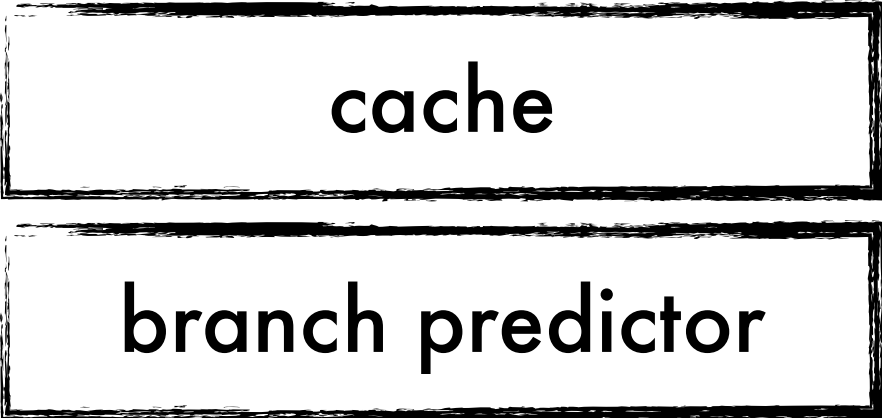
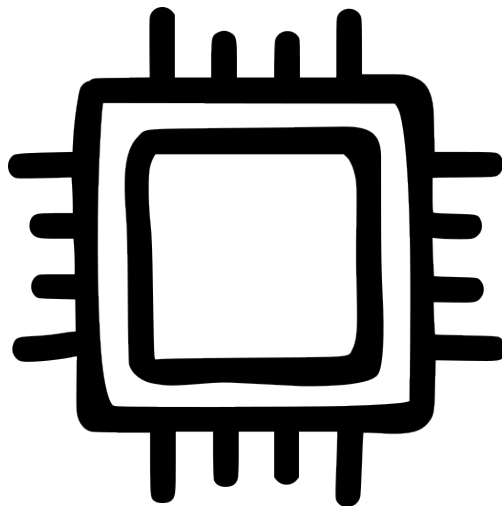
Mem

CT

OS

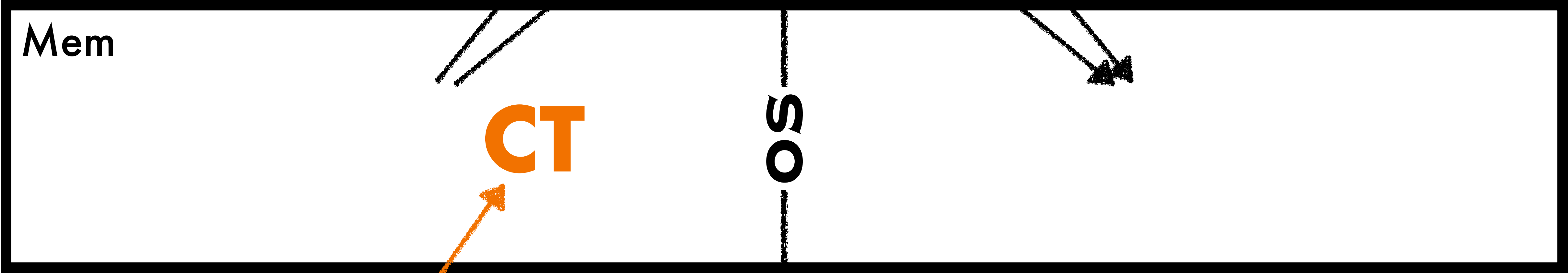


Formal Side Channel Security 101



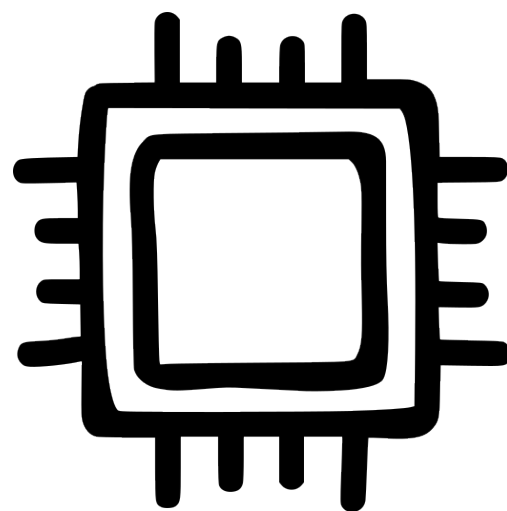
harmless
leak

observe

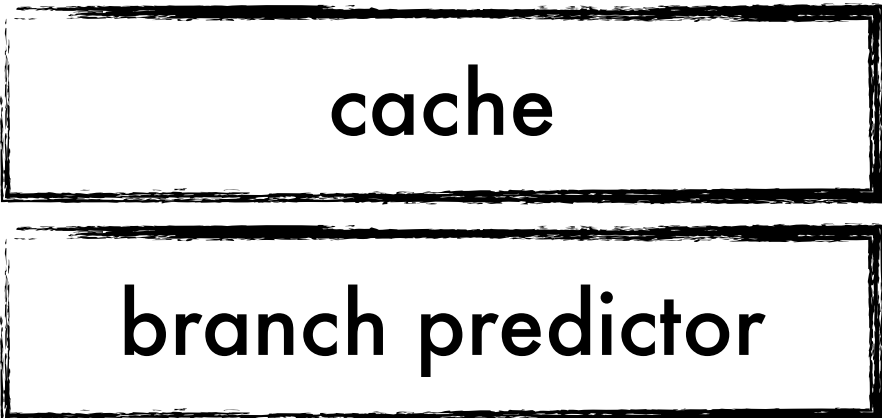


CT-Analysis

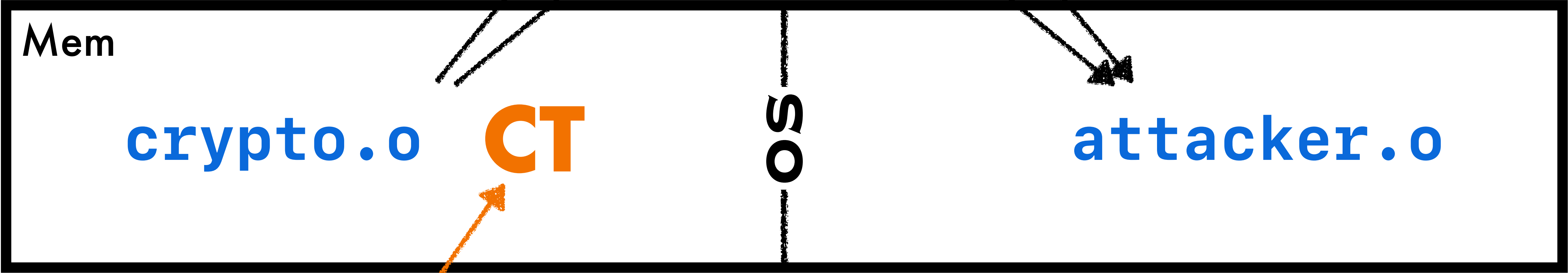
Formal Side Channel Security 101



harmless
leak



observe

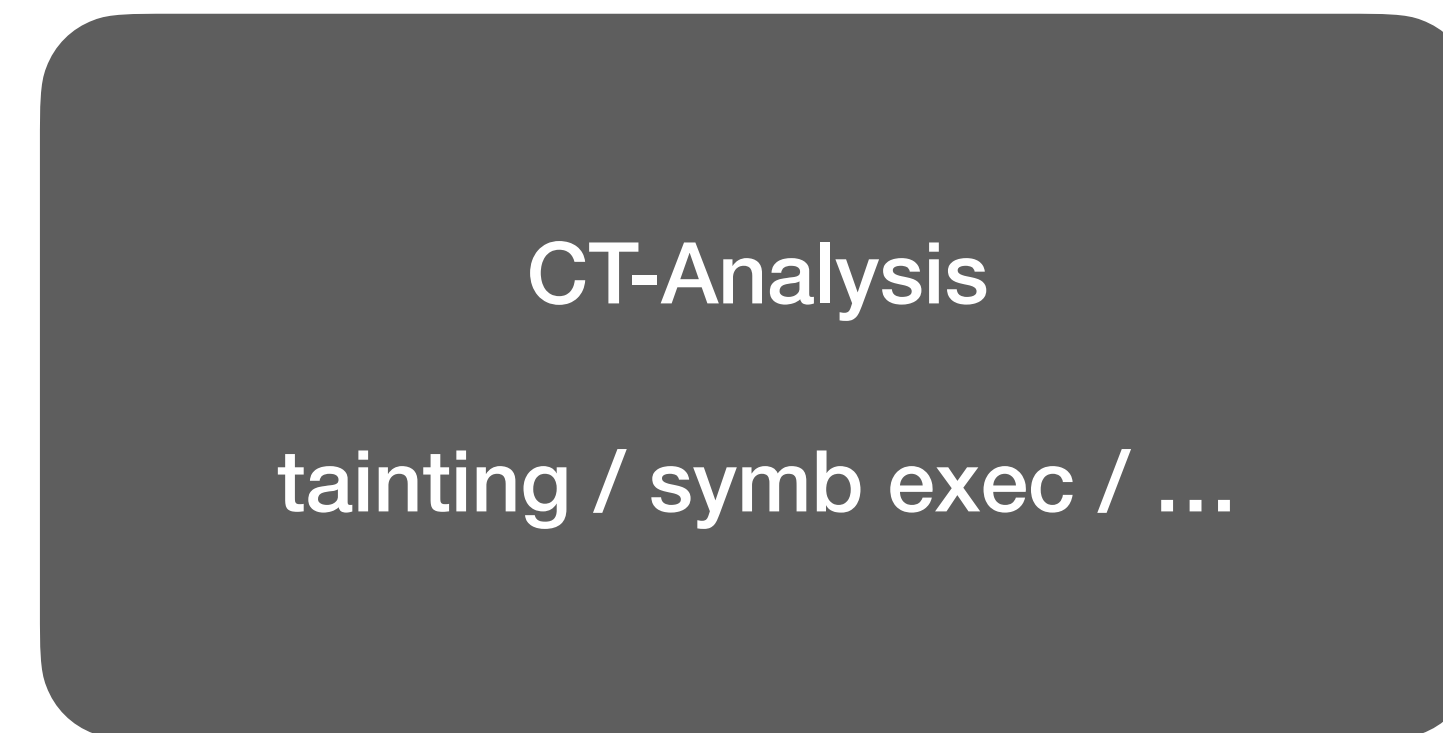


CT-Analysis

Constant Time Analysis

Remove functions

```
<> trusted-crypto-lib.c Raw
1  encrypt(secret_data, key);
2  create_hmac(secret_data, key);
3  ...
```



CT



~~**CT**~~

Decompile-then-Analyze

crypto.o

CT?

binary level

Decompile-then-Analyze

crypto.o

CT?

binary level
source level

CT-Analysis

Decompile-then-Analyze

crypto.o

CT?

binary level
source level

decompile

crypto.c

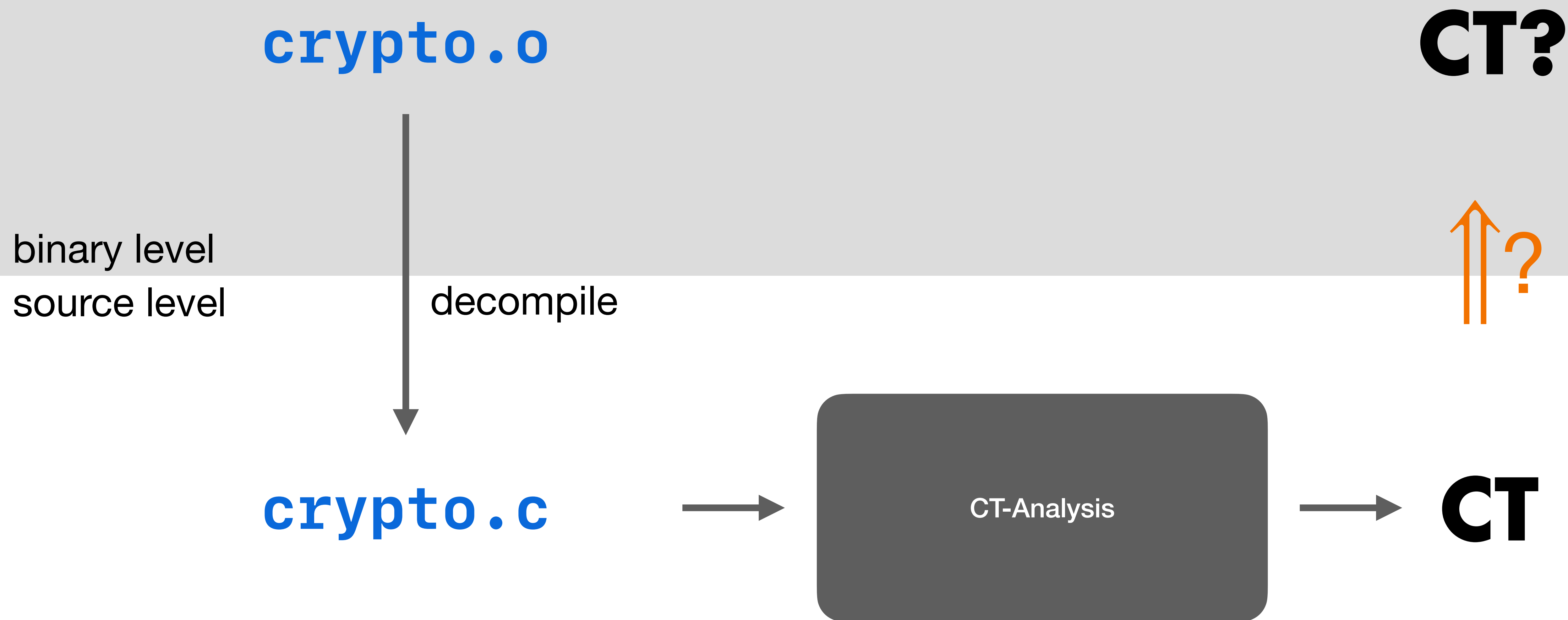


CT-Analysis



CT

Decompile-then-Analyze



Decompile-then-Analyze

CLangover (CVE-2024-37880)

ML-KEM.o

```
.LBB0_2: ; for(j = 0; j < 8; j++)  
    movzx    r8d, byte ptr [rsi + rax]  
    xor      edx, edx  
    bt       r8d, ecx  
    jae      .LBB0_4 ; leaks bit of msg  
    mov      edx, 1665  
.LBB0_4:  
    mov      word ptr [rdi + 2*rcx], dx  
    inc      rcx  
    cmp      rcx, 8  
    jne      .LBB0_2
```

Decompile-then-Analyze

CLangover (CVE-2024-37880)

ML-KEM.o

```
.LBB0_2: ; for(j = 0; j < 8; j++)  
    movzx    r8d, byte ptr [rsi + rax]  
    xor      edx, edx  
    bt       r8d, ecx  
    jae      .LBB0_4 ; leaks bit of msg  
    mov      edx, 1665  
.LBB0_4:  
    mov      word ptr [rdi + 2*rcx], dx  
    inc      rcx  
    cmp      rcx, 8  
    jne      .LBB0_2
```

Decompile-then-Analyze

CLangover (CVE-2024-37880)



```
.LBB0_2: ; for(j = 0; j < 8; j++)
    movzx    r8d, byte ptr [rsi + rax]
    xor      edx, edx
    bt       r8d, ecx
    jae      .LBB0_4 ; leaks bit of msg
    mov      edx, 1665
.LBB0_4:
    mov      word ptr [rdi + 2*rcx], dx
    inc      rcx
    cmp      rcx, 8
    jne      .LBB0_2
```

Decompile-then-Analyze

CLangover (CVE-2024-37880)

ML-KEM.o $\xrightarrow{\text{RetDec}}$ **ML-KEM.c**

```
.LBB0_2: ; for(j = 0; j < 8; j++)  
    movzx    r8d, byte ptr [rsi + rax]  
    xor      edx, edx  
    bt       r8d, ecx  
    jae      .LBB0_4 ; leaks bit of msg  
    mov      edx, 1665  
.LBB0_4:  
    mov      word ptr [rdi + 2*rcx], dx  
    inc      rcx  
    cmp      rcx, 8  
    jne      .LBB0_2
```

```
for (int64_t j = 0; j < 8; j++) {  
    byte = *(char *) (msg_offset + msg);  
    coef = (1 << (int32_t)j % 32 & (int32_t)byte) == 0 ?  
            0 : 1665; // no leak  
    *(int16_t *) (2 * j + poly_offset) = coef;  
}
```

Decompile-then-Analyze



binary level

source level



Decompile-then-Analyze

~~CT~~ crypto.o

binary level
source level

decompile

CT crypto.c



Decompile-then-Analyze

~~CT~~ crypto.o

binary level
source level

decompile

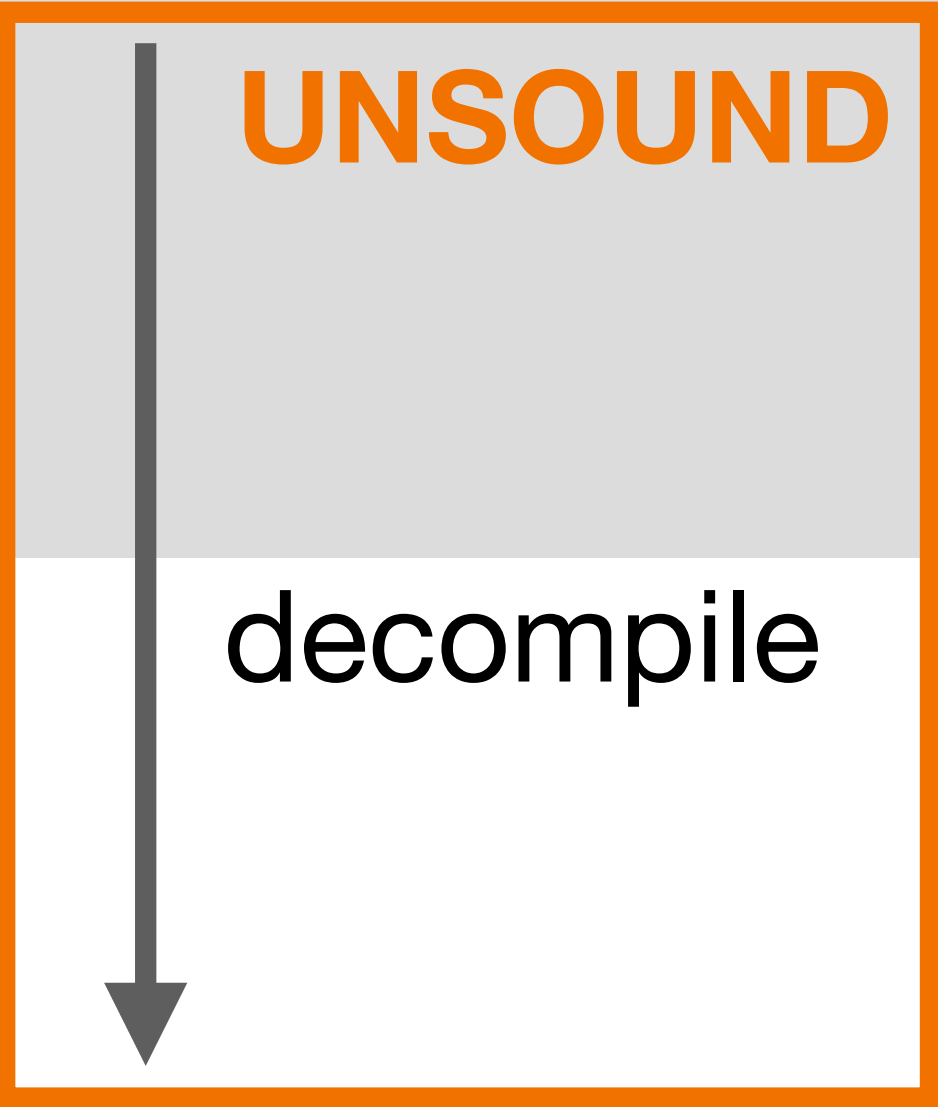
CT crypto.c



CT

Decompile-then-Analyze

~~CT~~ crypto.o



binary level
source level

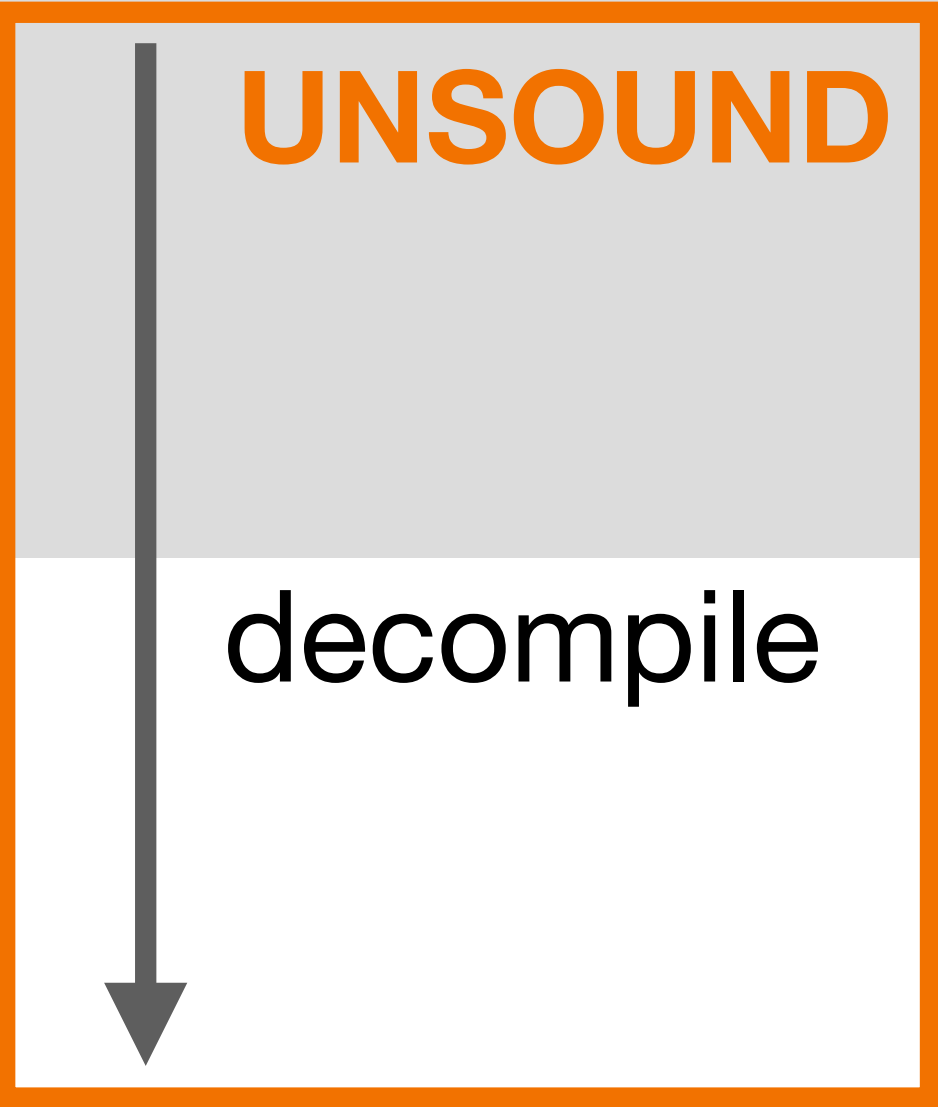
CT crypto.c



CT

Decompile-then-Analyze

~~CT~~ crypto.o



ANGR
BINARY NINJA
GHIDRA
RETDEC
HEX-RAYS

binary level
source level

CT crypto.c



CT

Decompile-then-Analyze

~~cr~~ crypto.o

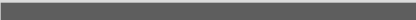


~~cr~~

binary level

Decompile-then-Analyze

crypto.o



~~cr~~

binary level



Decompile-then-Analyze

crypto.o

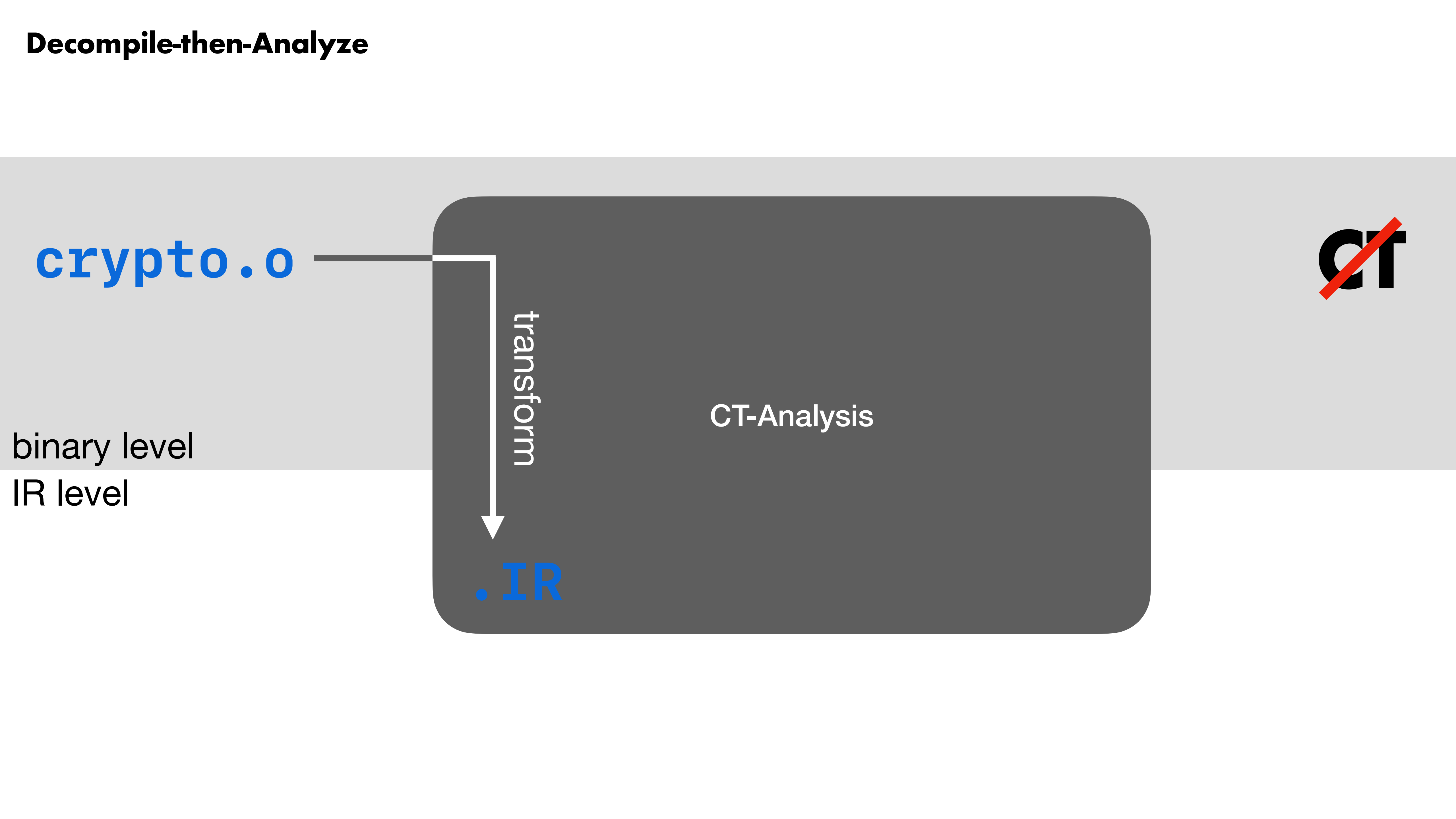
~~cr~~

binary level
IR level

transform

CT-Analysis

.IR

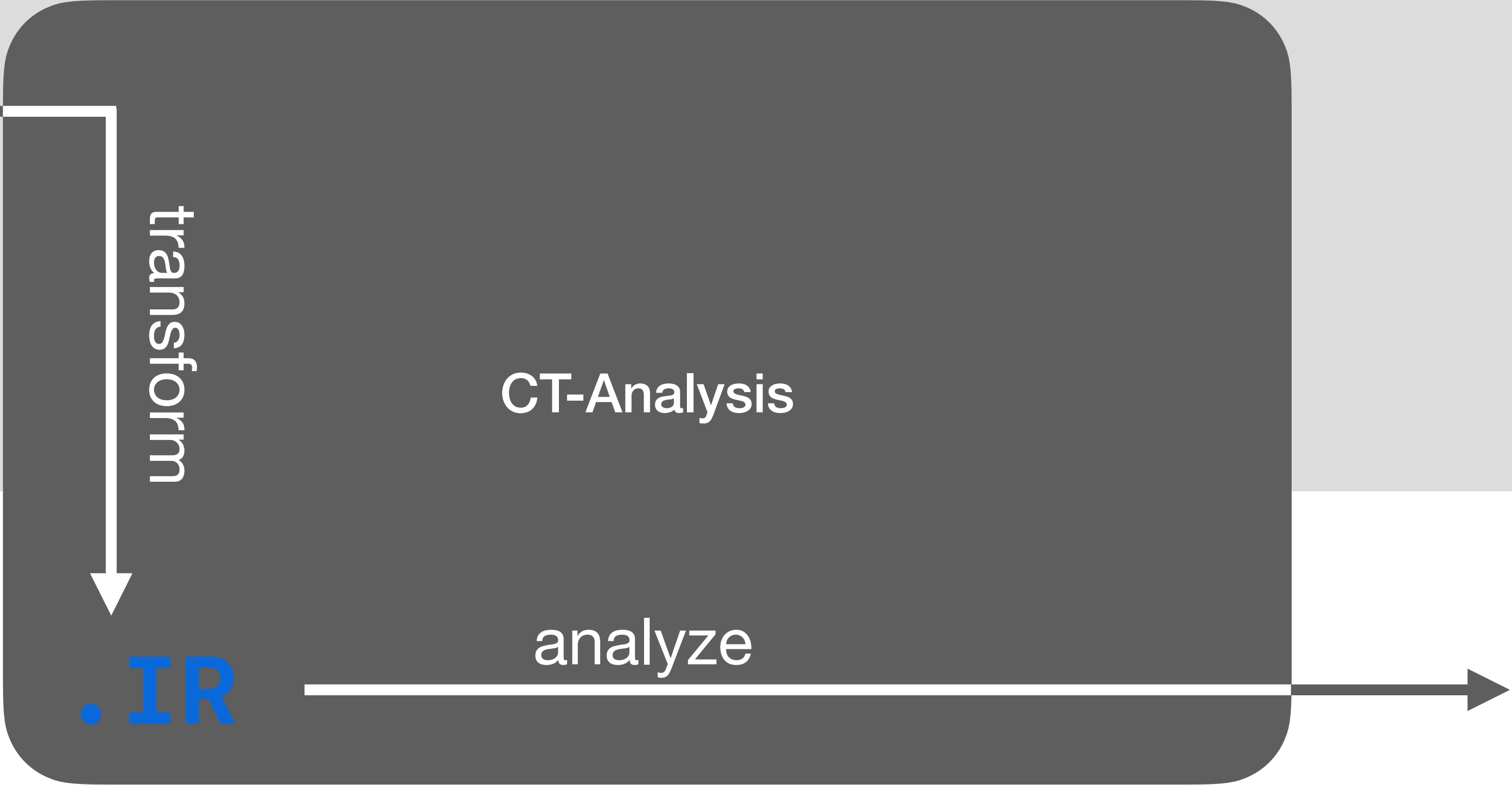


Decompile-then-Analyze

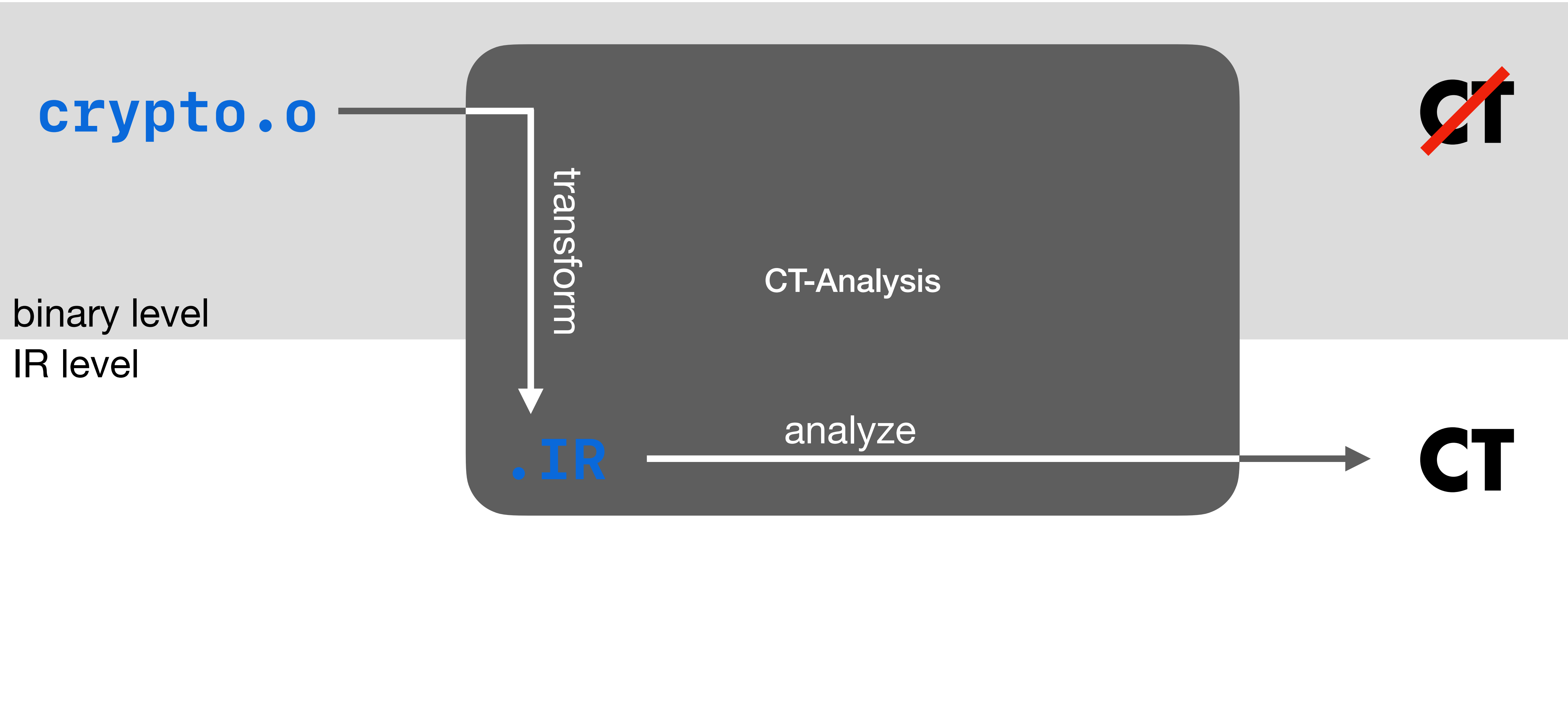
crypto.o

~~cr~~

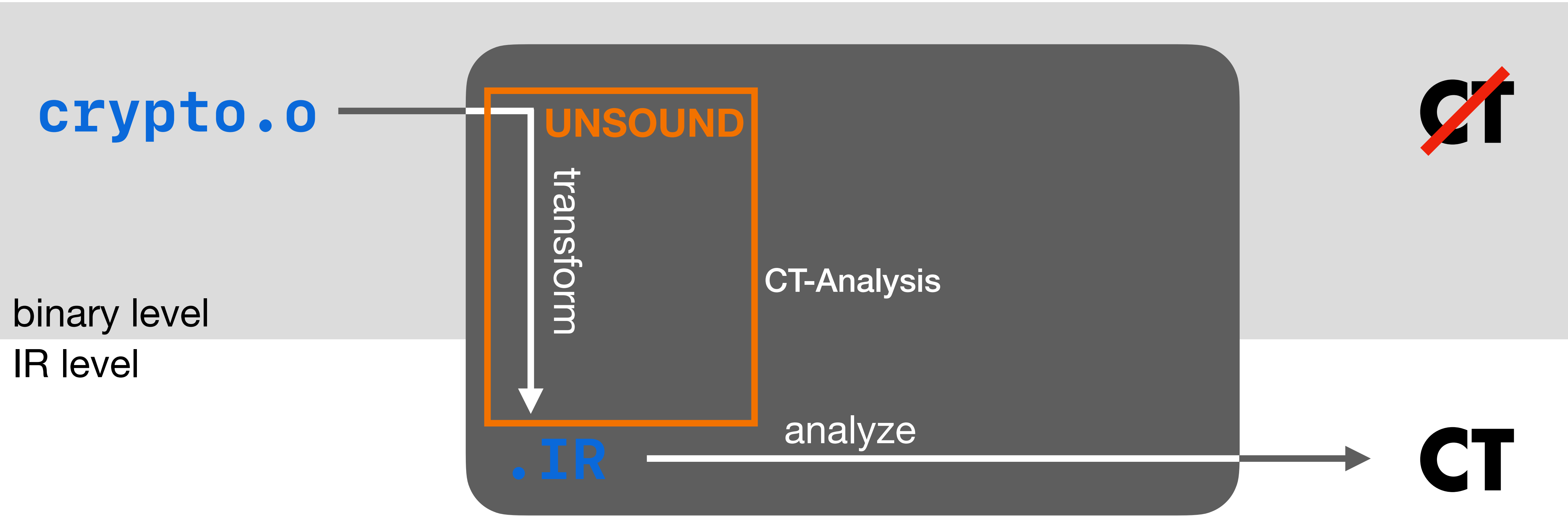
binary level
IR level



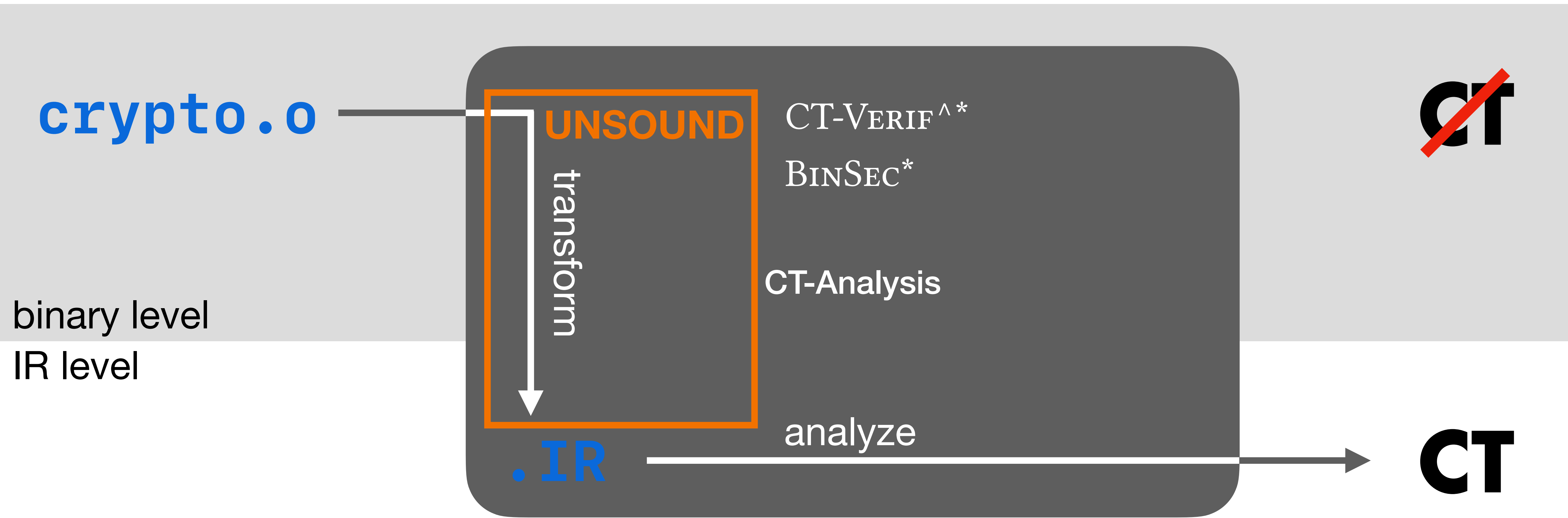
Decompile-then-Analyze



Decompile-then-Analyze



Decompile-then-Analyze



[^]: input is LLVM-IR

^{*}: on handcrafted **example.o**

Properties for Sound Decompilation



Properties for Sound Decompilation



Properties for Sound Decompile



`[ ]` reflects **CT**

`[ ]` preserves **CT**

`[ ]` **CT** transparent

Properties for Sound Decompilation



[] reflects **CT**

WHEN

prog ~~**CT**~~ $\implies$ **[prog]** ~~**CT**~~
“**[]** removes no vulnerabilities”

[] preserves **CT**

[] **CT** transparent

Properties for Sound Decompilation



[] reflects **CT**

WHEN

prog ~~**CT**~~ $\implies$ **[prog]** ~~**CT**~~
“**[]** removes no vulnerabilities”

[] preserves **CT**

WHEN

prog **CT** $\implies$ **[prog]** **CT**
“**[]** introduces no vulnerabilities”

[] **CT** transparent

Properties for Sound Decompile



`[ ]` reflects CT

WHEN

`prog` ~~CT~~ $\implies$ `[prog]` ~~CT~~
“`[ ]` removes no vulnerabilities”

`[ ]` preserves CT

WHEN

`prog` CT $\implies$ `[prog]` CT
“`[ ]` introduces no vulnerabilities”

`[ ]` CT transparent

WHEN

`prog` CT $\iff$ `[prog]` CT

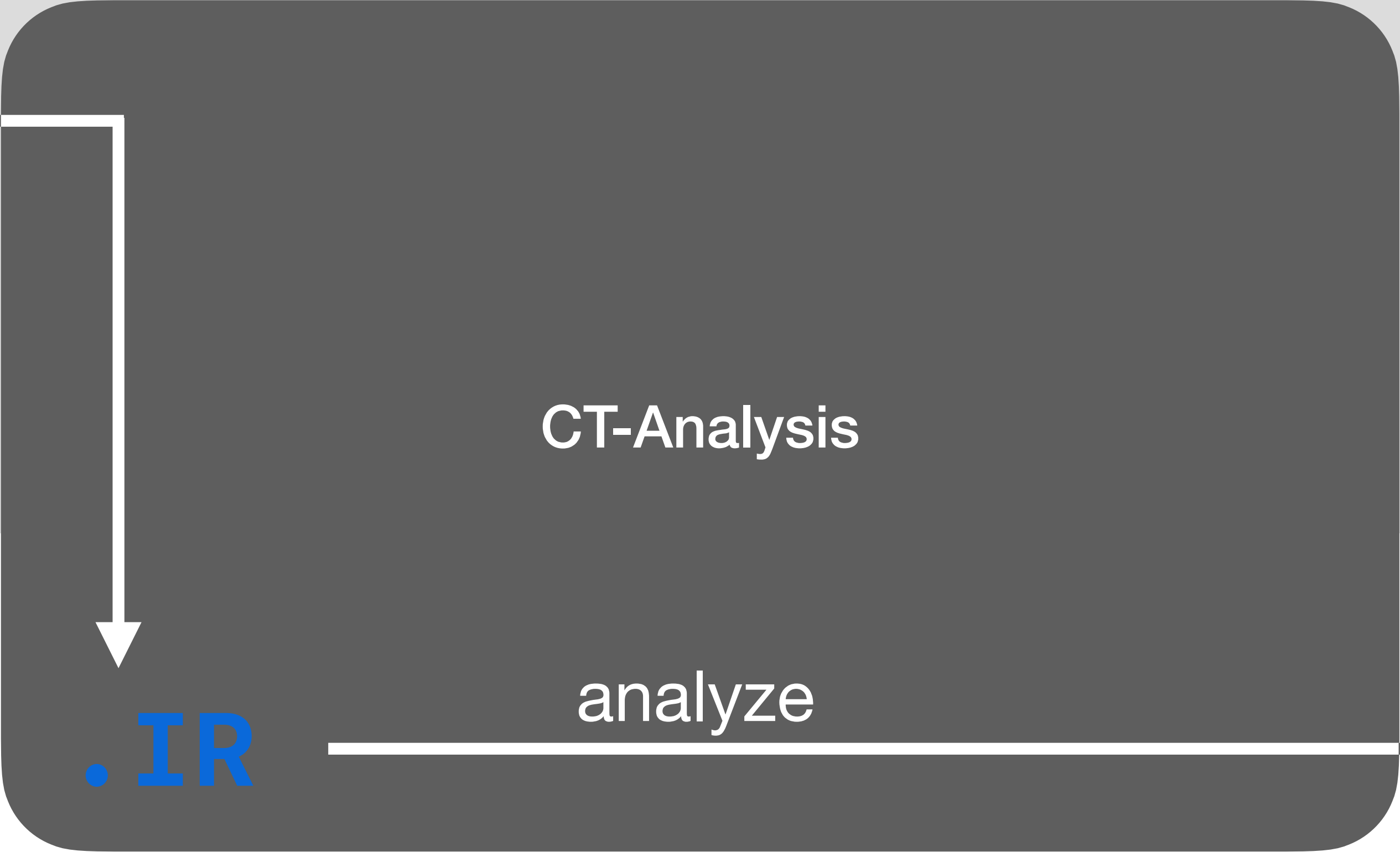
Decompile-then-Analyze

crypto.o

~~cr~~

binary level

IR level

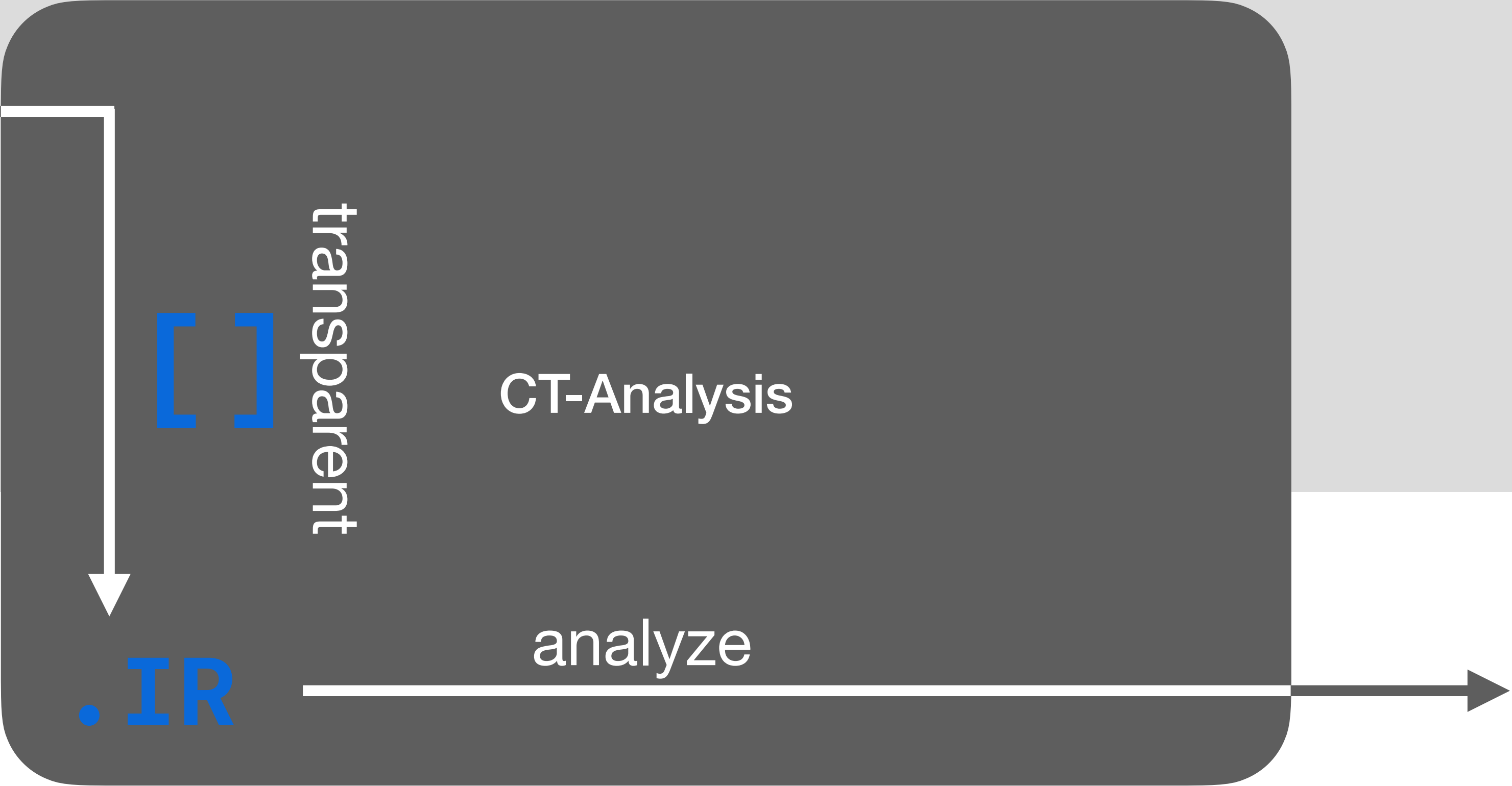


Decompile-then-Analyze

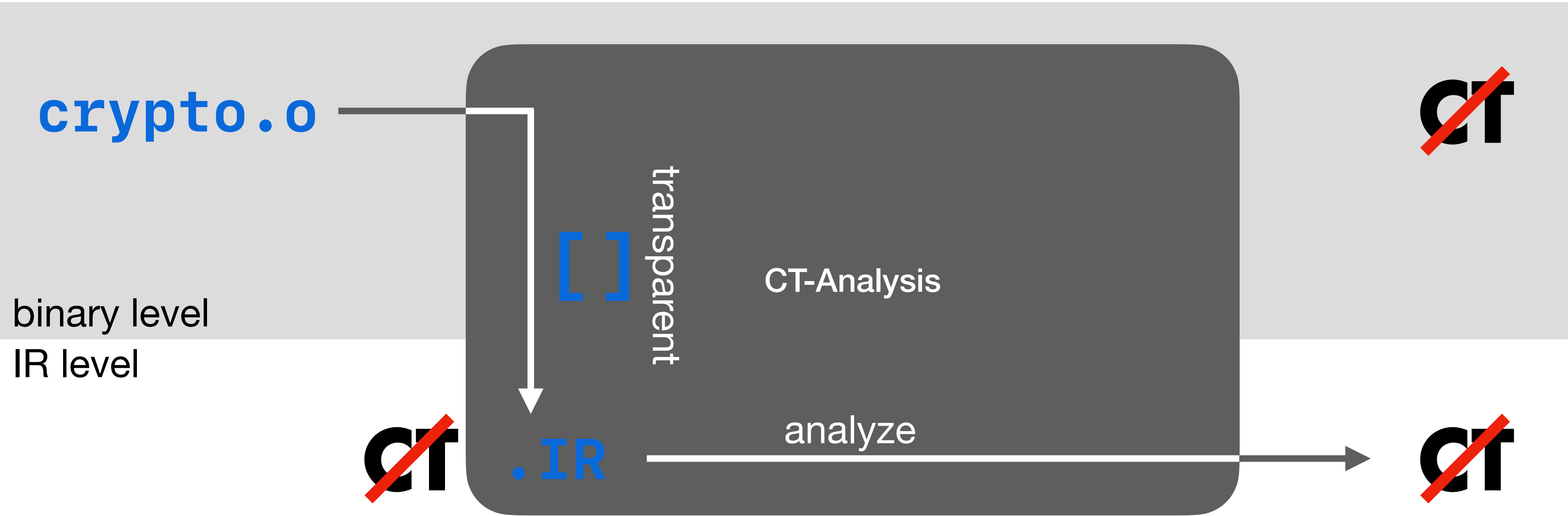
crypto.o

~~cr~~

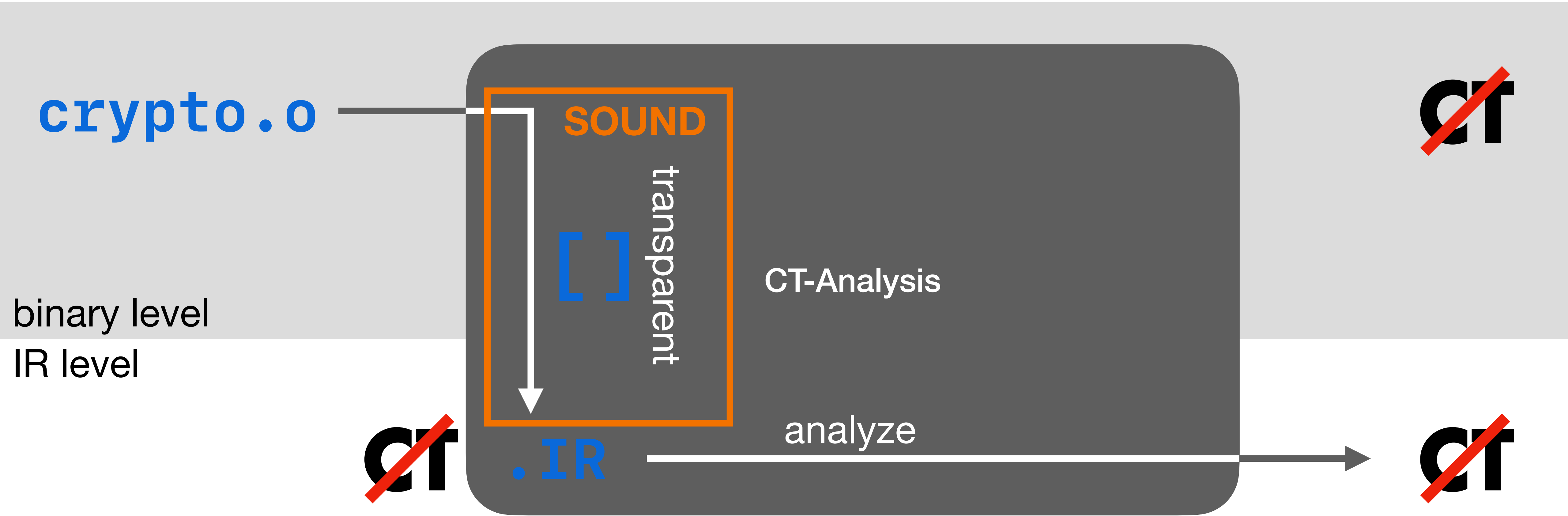
binary level
IR level



Decompile-then-Analyze



Decompile-then-Analyze



Proving “[] is CT transparent”

Proving “[] is CT transparent”

start with [] preserves **CT** and [] reflects **CT**

$$\text{prog CT} \implies [\text{prog}] \text{CT}$$

Proving “[] preserves CT”

Proving “[] preserves CT”

$\text{prog CT} \implies [\text{prog}] \text{CT}$

prog

[prog]

Cites!

Proving "[] preserves CT"

prog CT

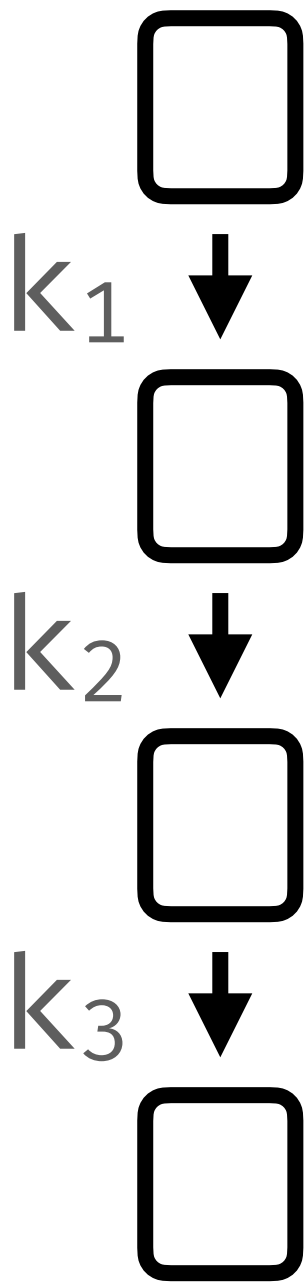
$\implies$

[prog] CT

prog

[prog]

Cites!



Proving "[] preserves CT"

prog CT

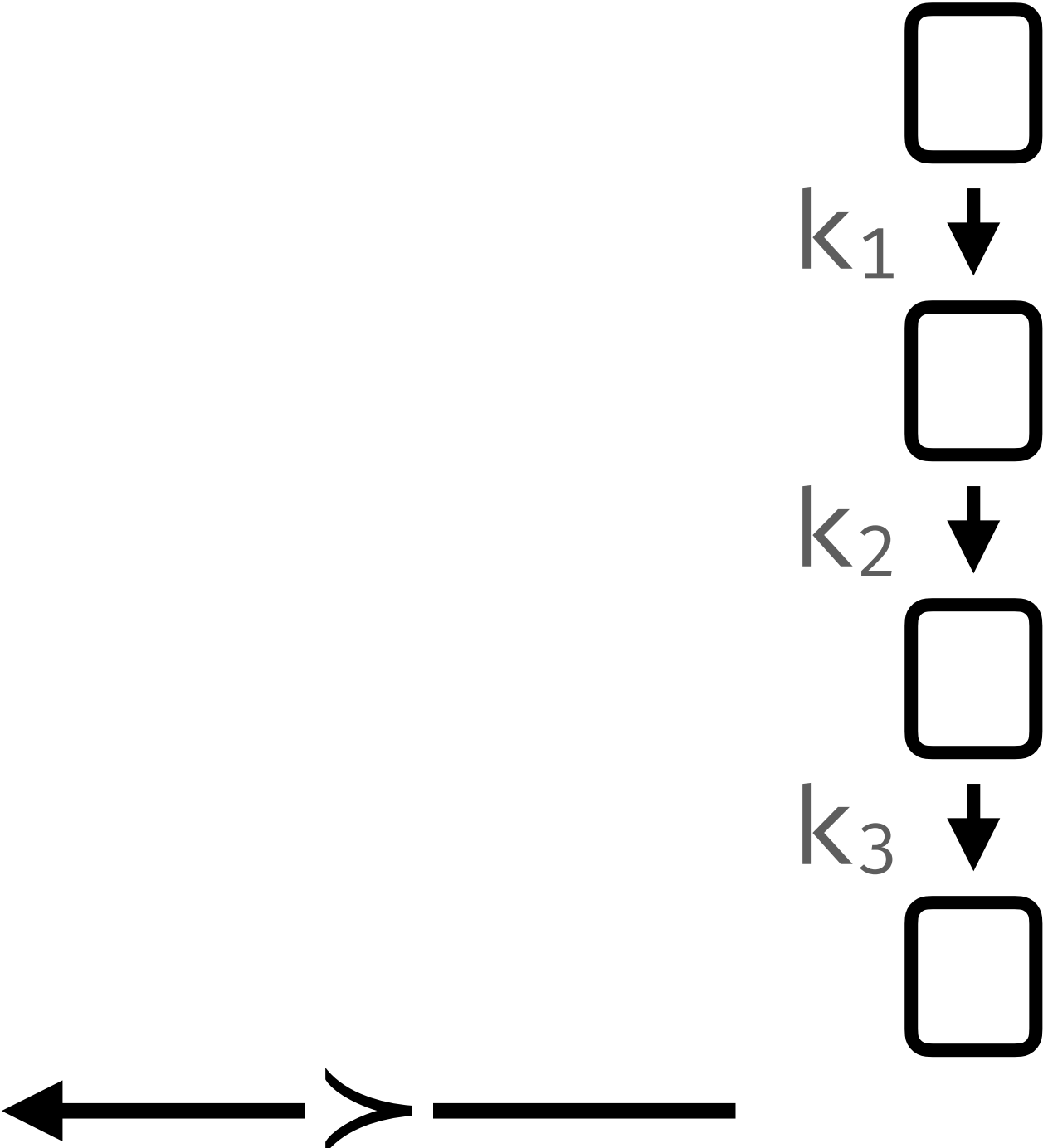
$\implies$

[prog] CT

Cites!

prog

[prog]



Proving "[] preserves CT"

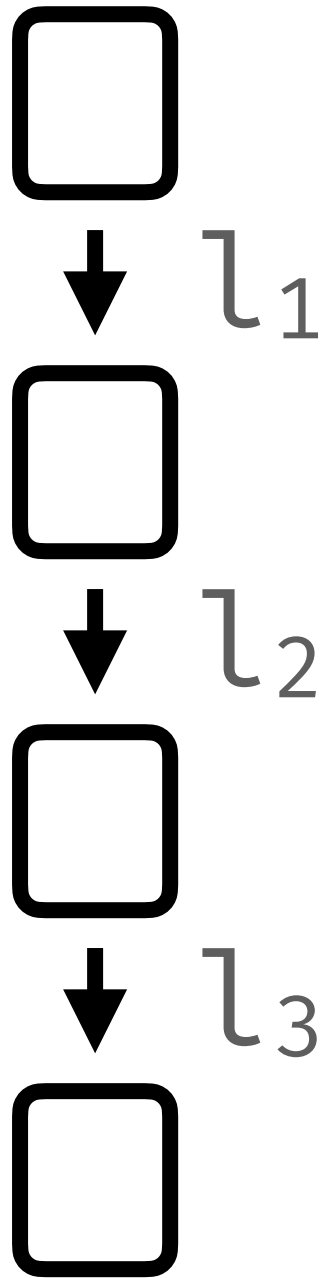
prog CT

$\implies$

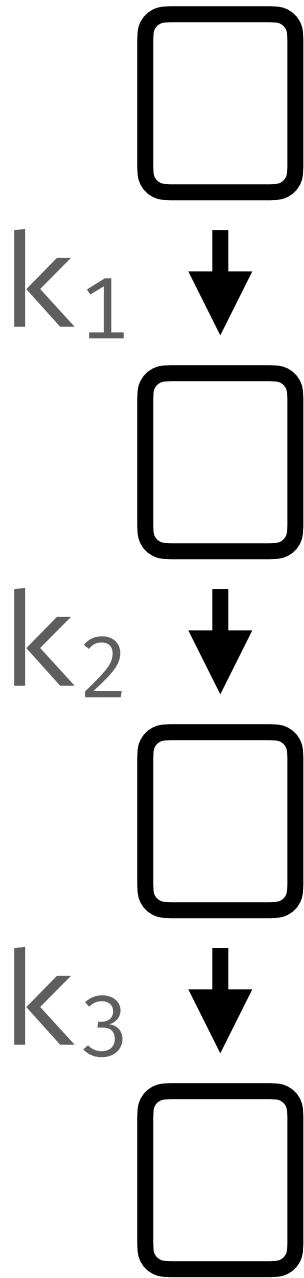
[prog] CT

Cites!

prog



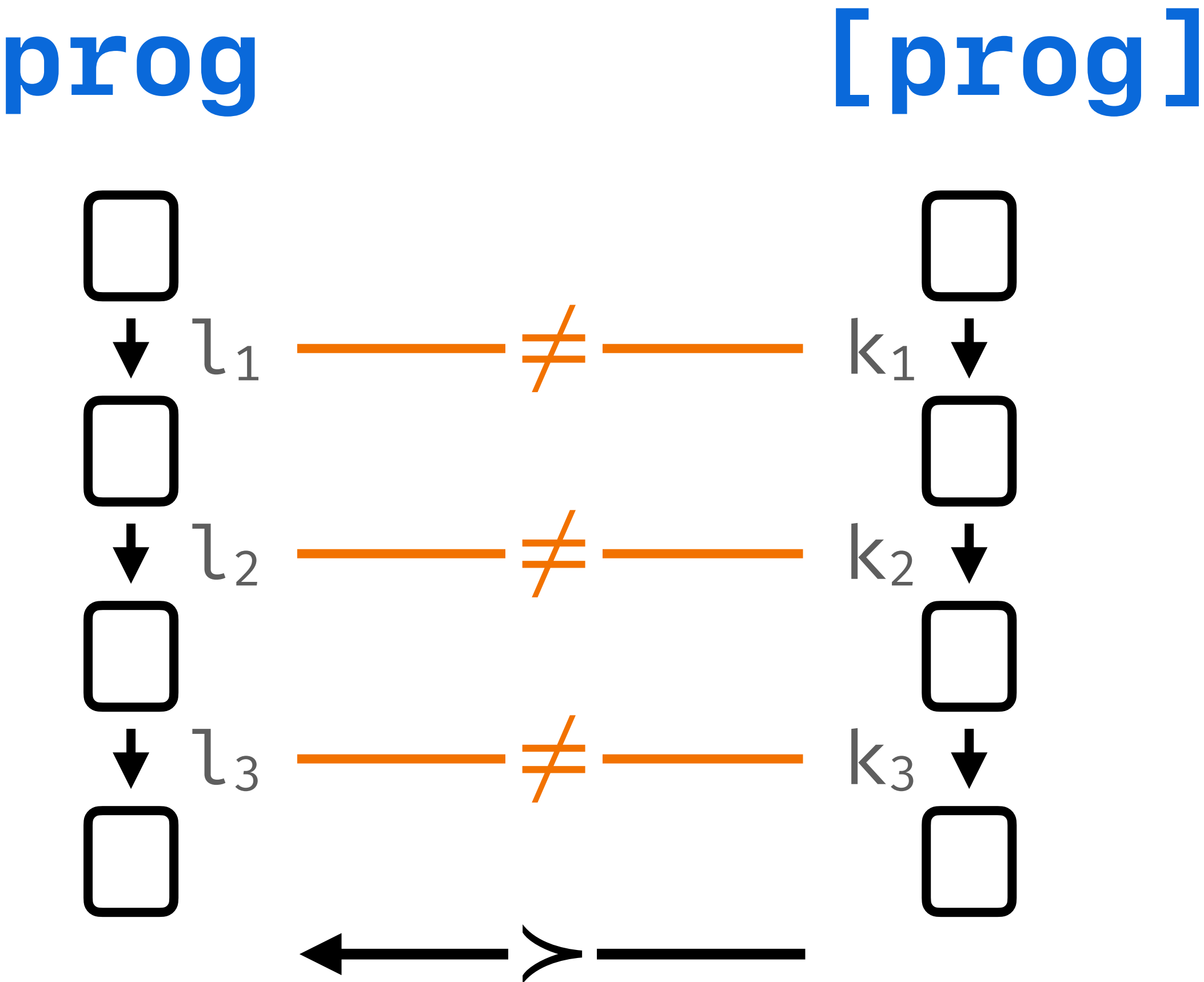
[prog]



Proving "[] preserves CT"

$\text{prog CT} \implies [\text{prog}] \text{CT}$

Cites!



Proving "[] preserves CT"

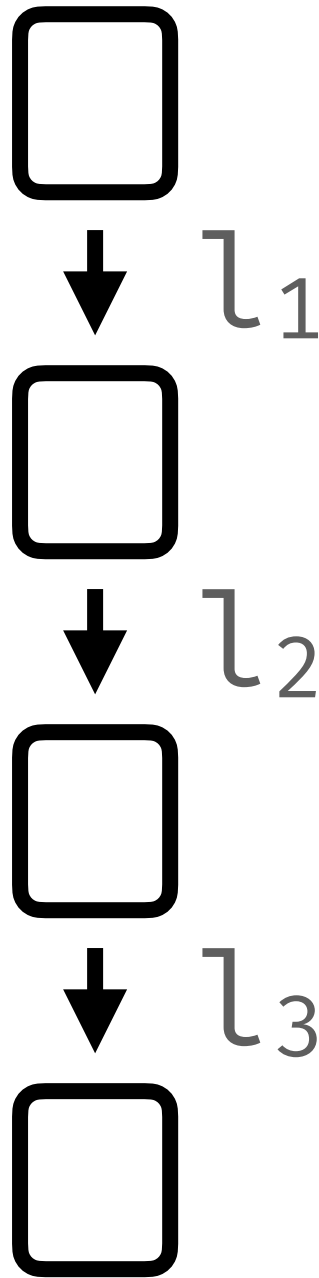
prog CT

$\implies$

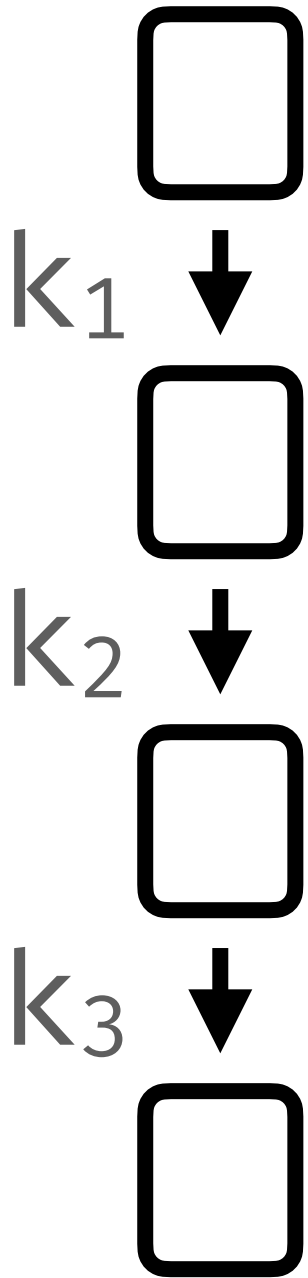
[prog] CT

Cites!

prog



[prog]



Proving "[] preserves CT"

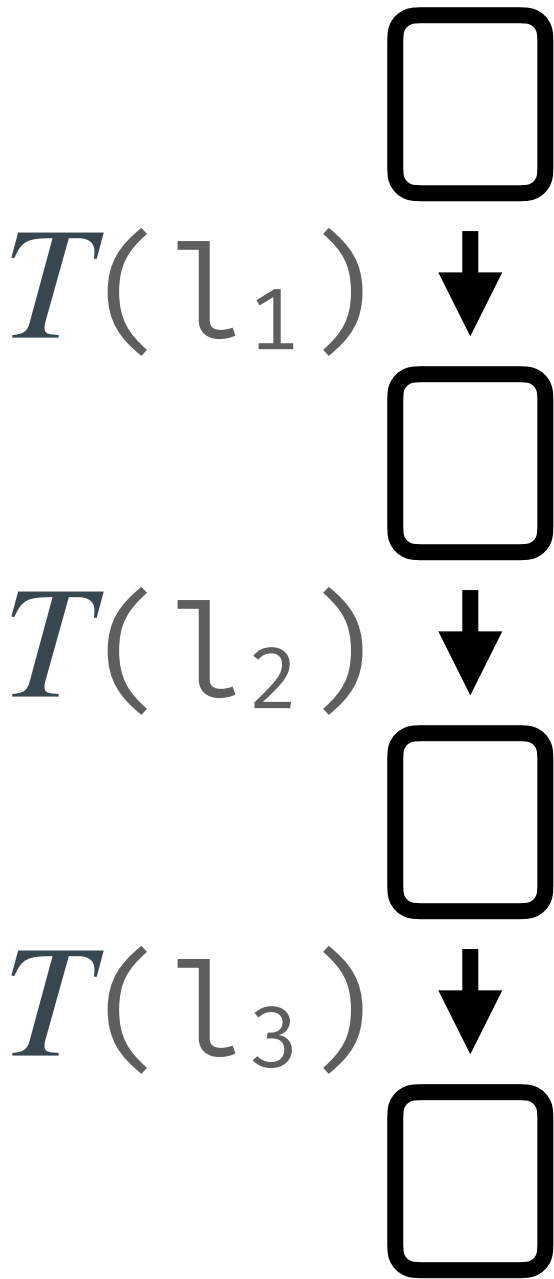
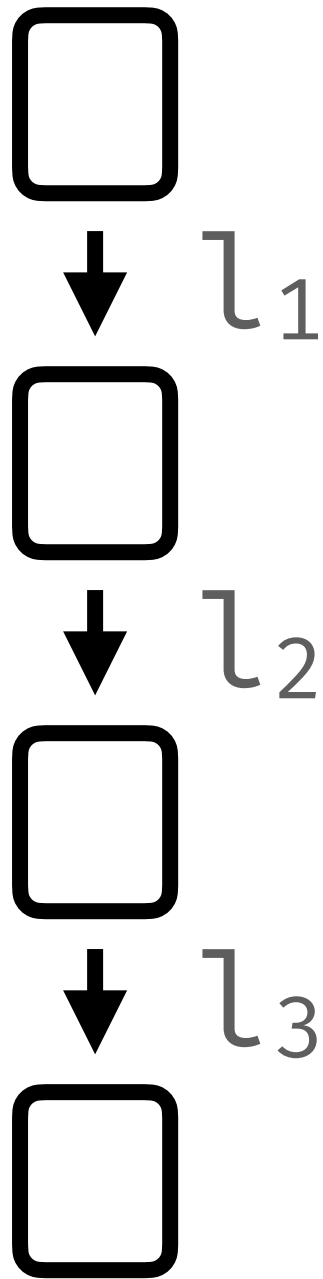
$\text{prog CT} \implies [\text{prog}] \text{CT}$

T : leakage
: transformer

Cites!

prog

[prog]



Proving "[] preserves CT"

prog CT

$\implies$

[prog] CT

T

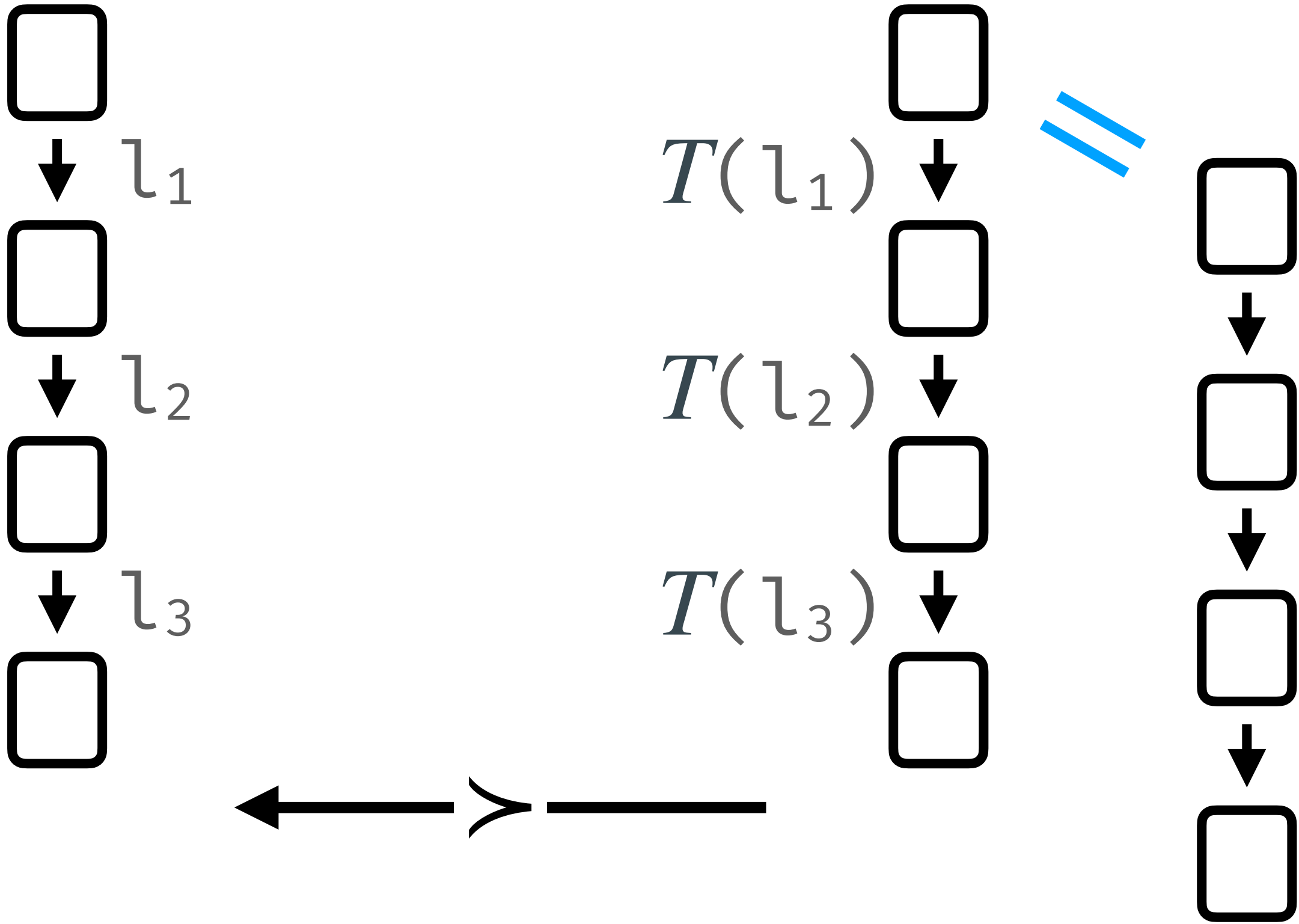
• leakage

• transformer

Cites!

prog

[prog]



Proving "[] preserves CT"

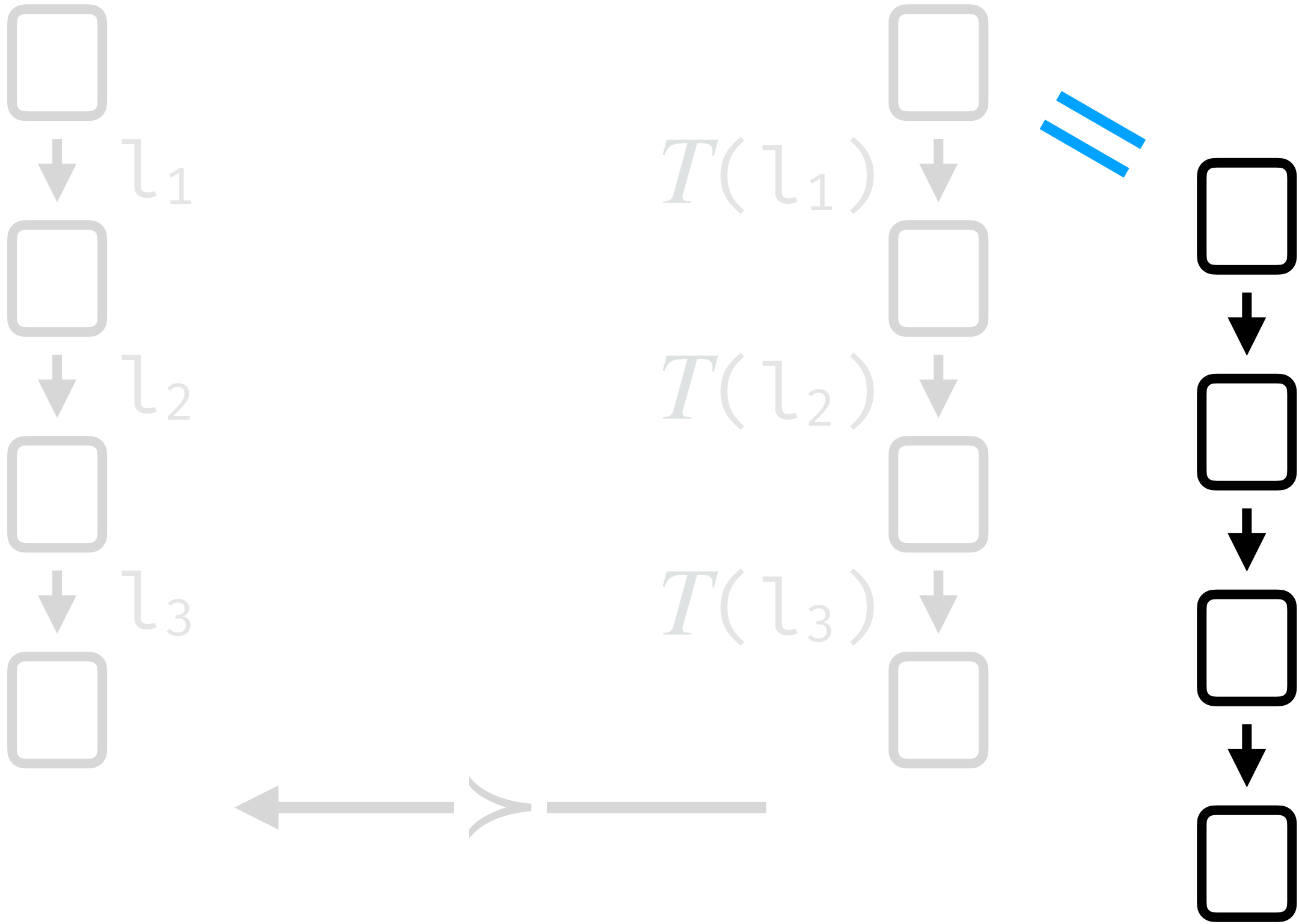
$\text{prog CT} \implies [\text{prog}] \text{CT}$

T : leakage
transformer

Cites!

prog

[prog]



Proving "[] preserves CT"

prog CT

⇒

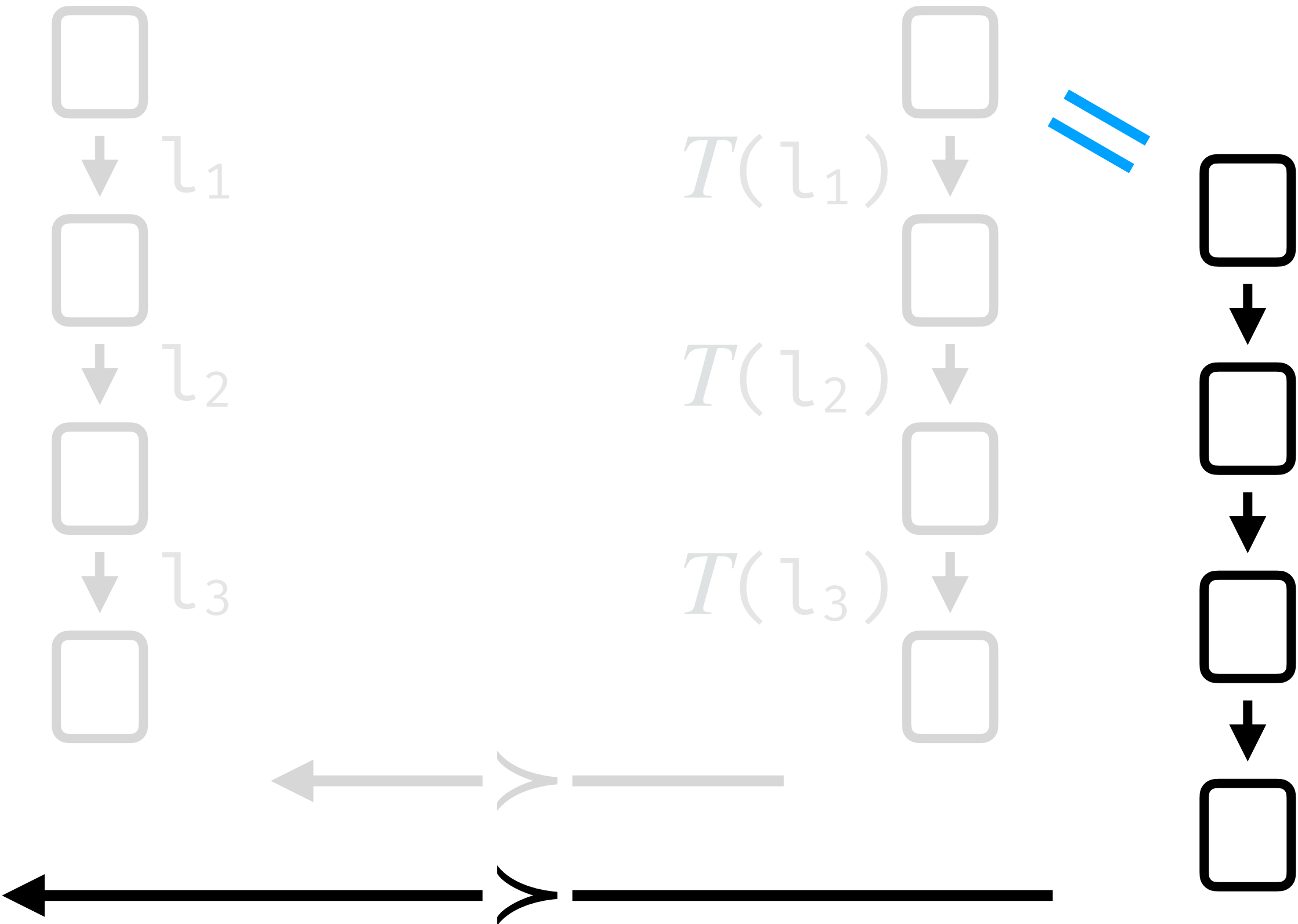
[prog] CT

T : leakage
transformer

Cites!

prog

[prog]



Proving "[] preserves CT"

prog CT

⇒

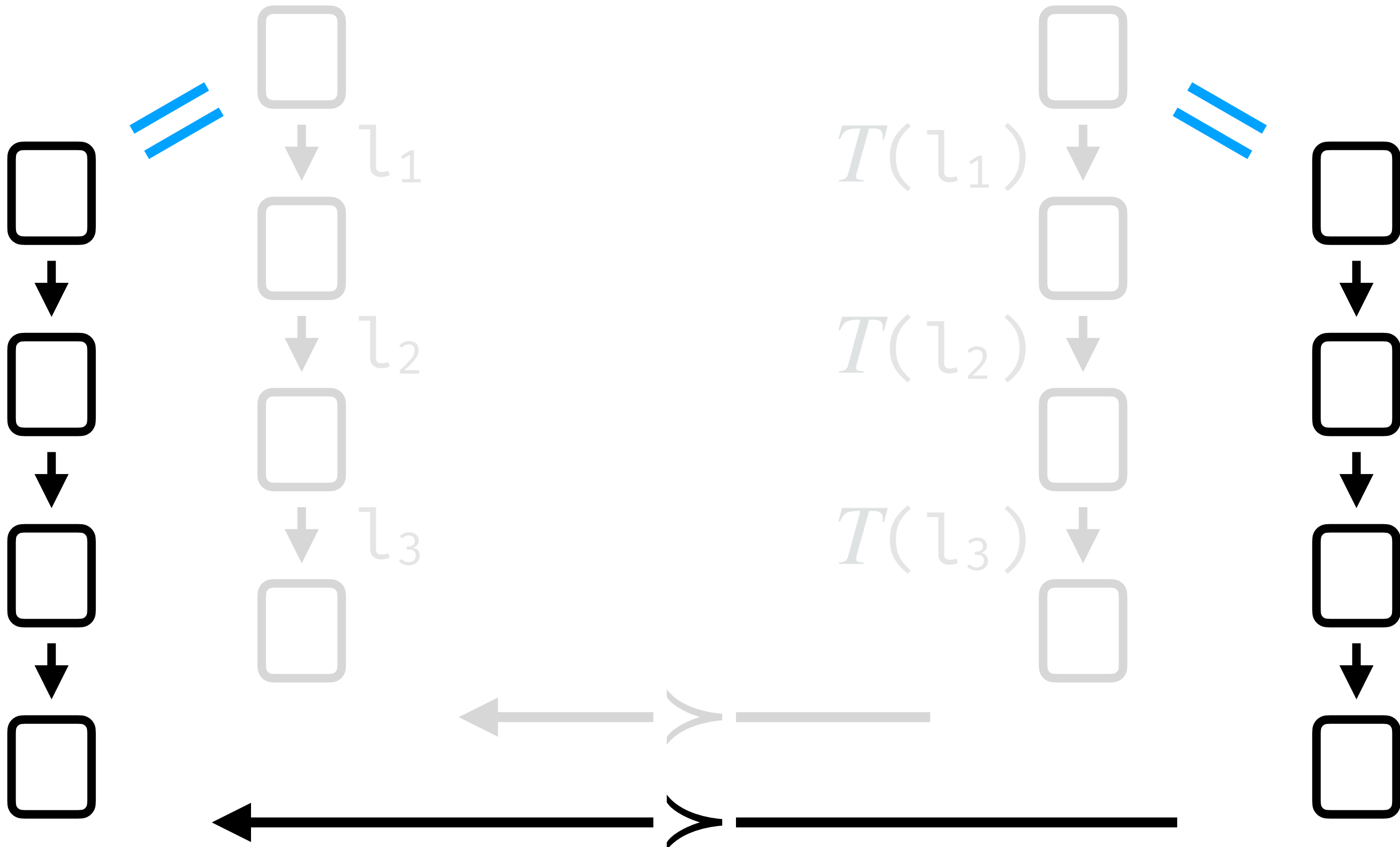
[prog] CT

T : leakage
transformer

Cites!

prog

[prog]



Proving "[] preserves CT"

prog CT

⇒

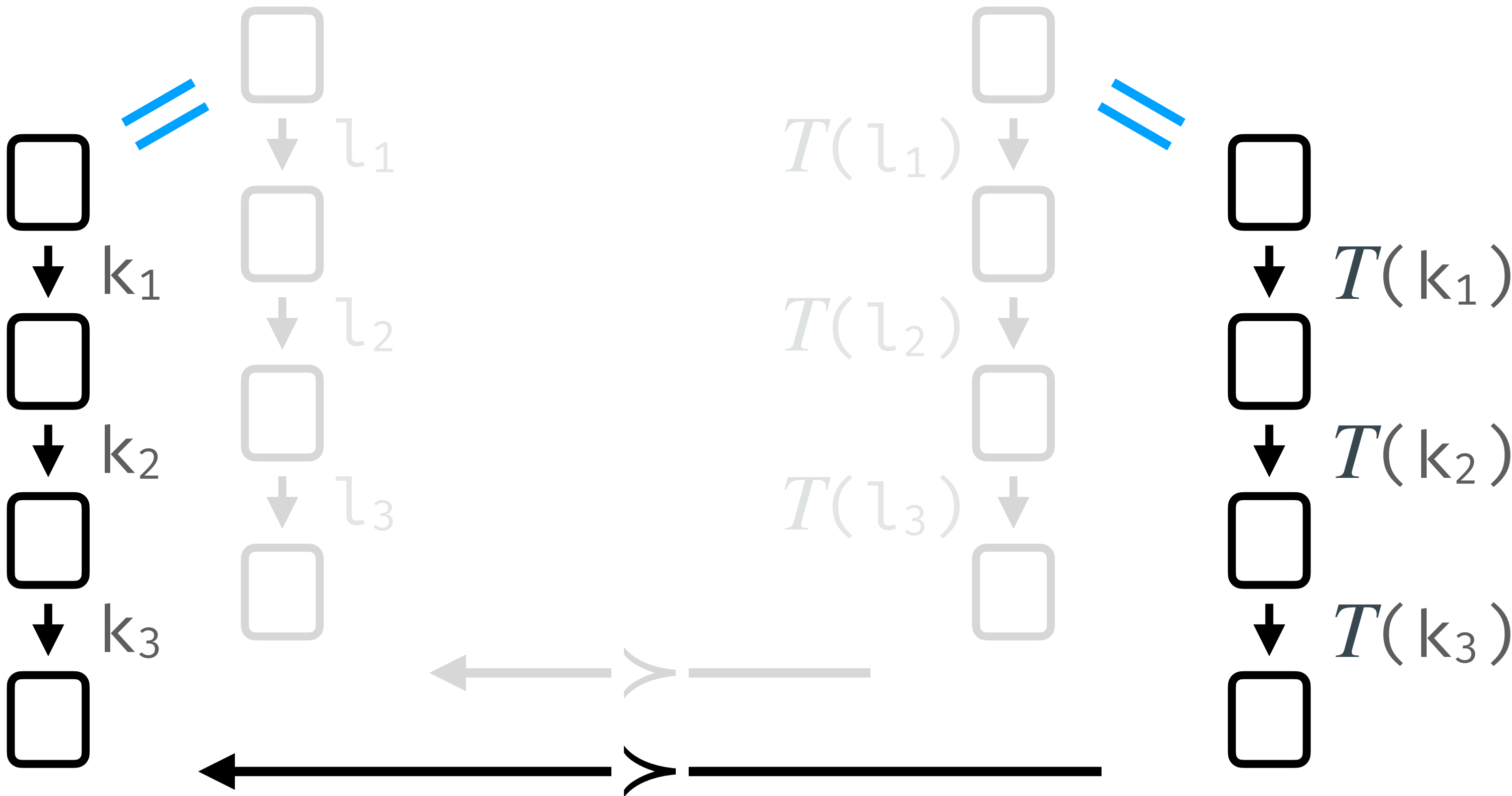
[prog] CT

T : leakage
transformer

Cites!

prog

[prog]



Proving "[] preserves CT"

prog CT

⇒

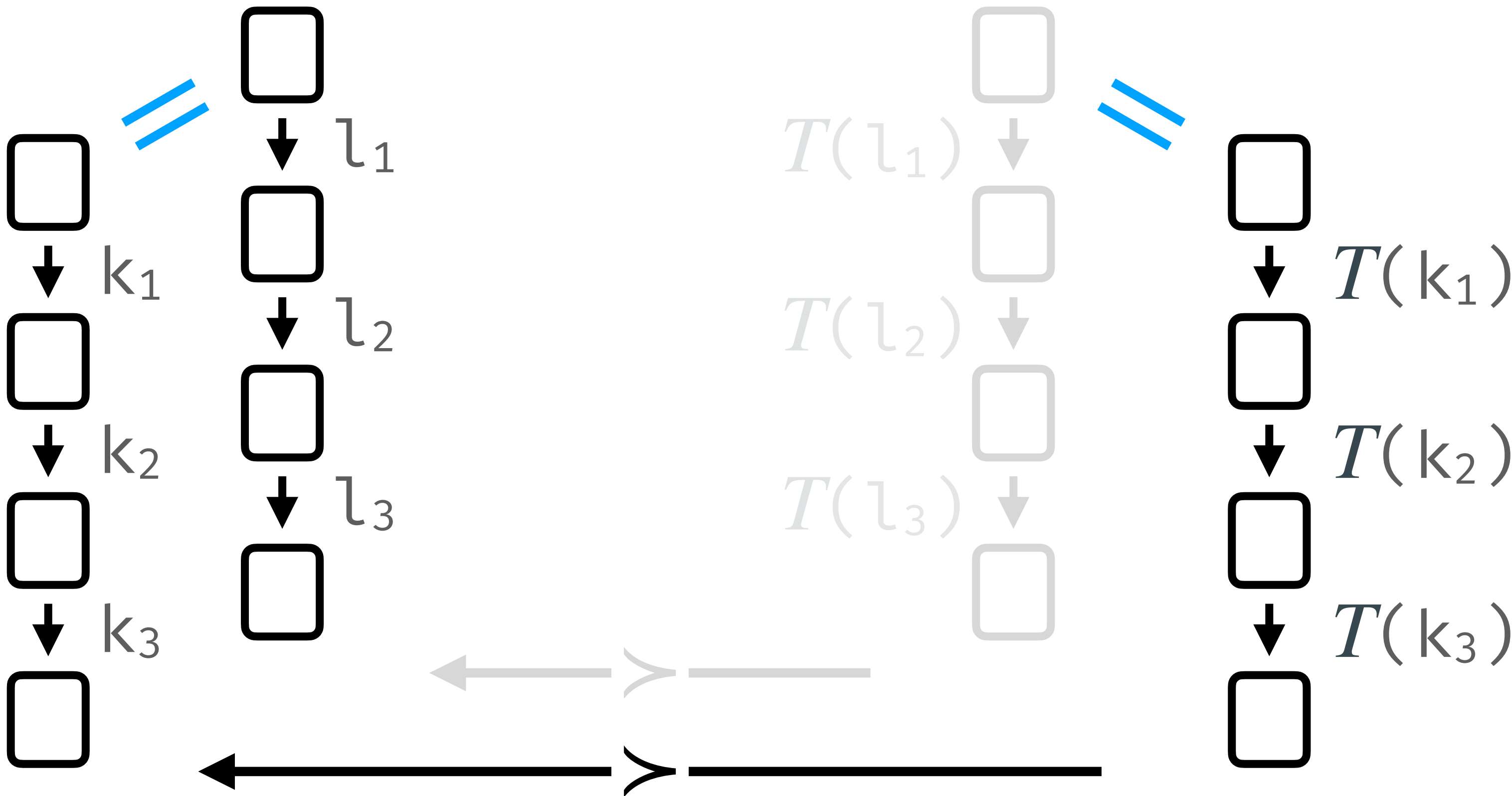
[prog] CT

T : leakage
: transformer

Cites!

CT prog

[prog]



Proving "[] preserves CT"

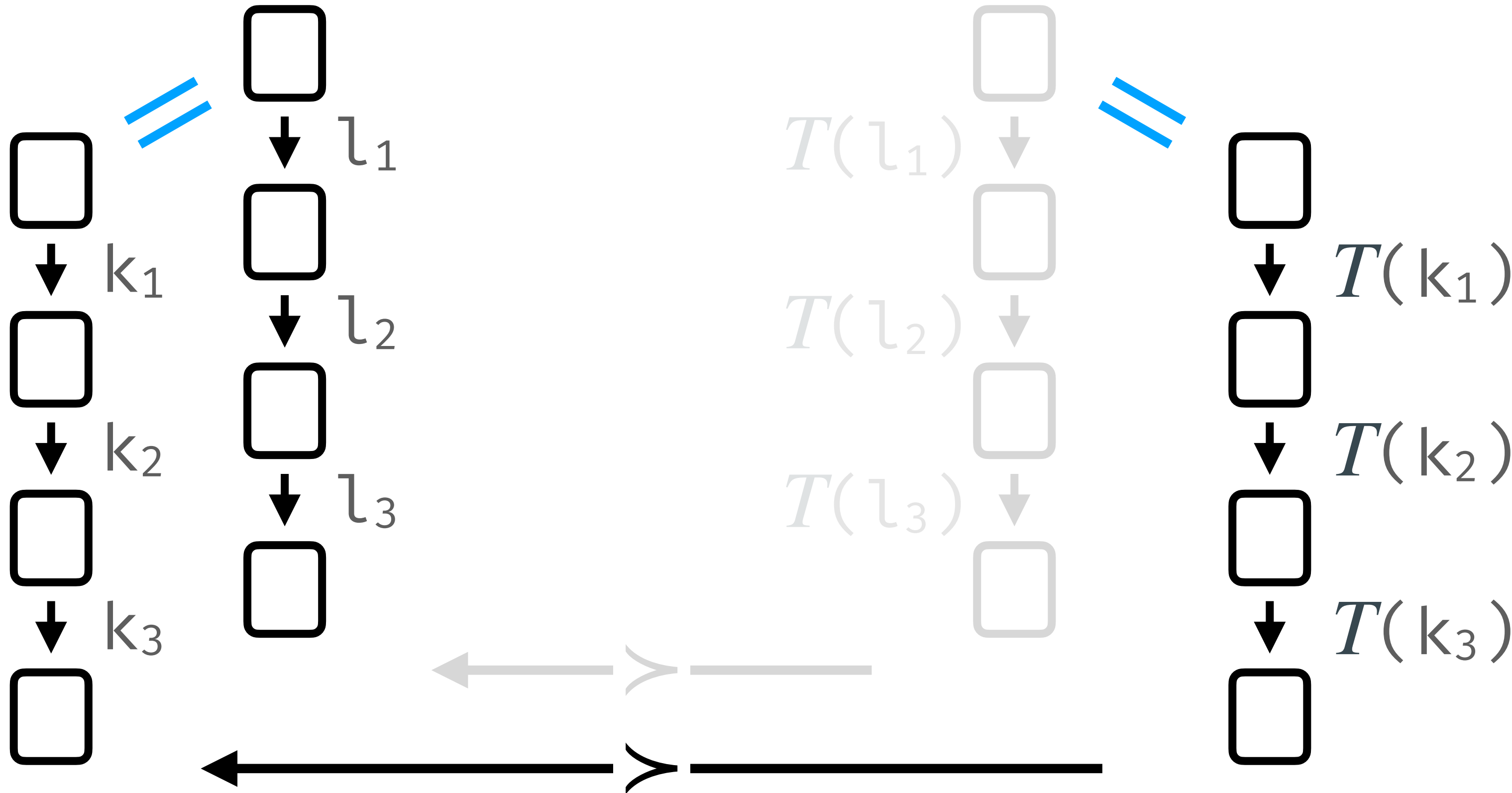
$\text{prog CT} \implies [\text{prog}] \text{CT}$

T : leakage
: transformer

Cites!

CT prog

[prog]



$k_1 k_2 k_3 \dots = l_1 l_2 l_3 \dots$

Proving "[] preserves CT"

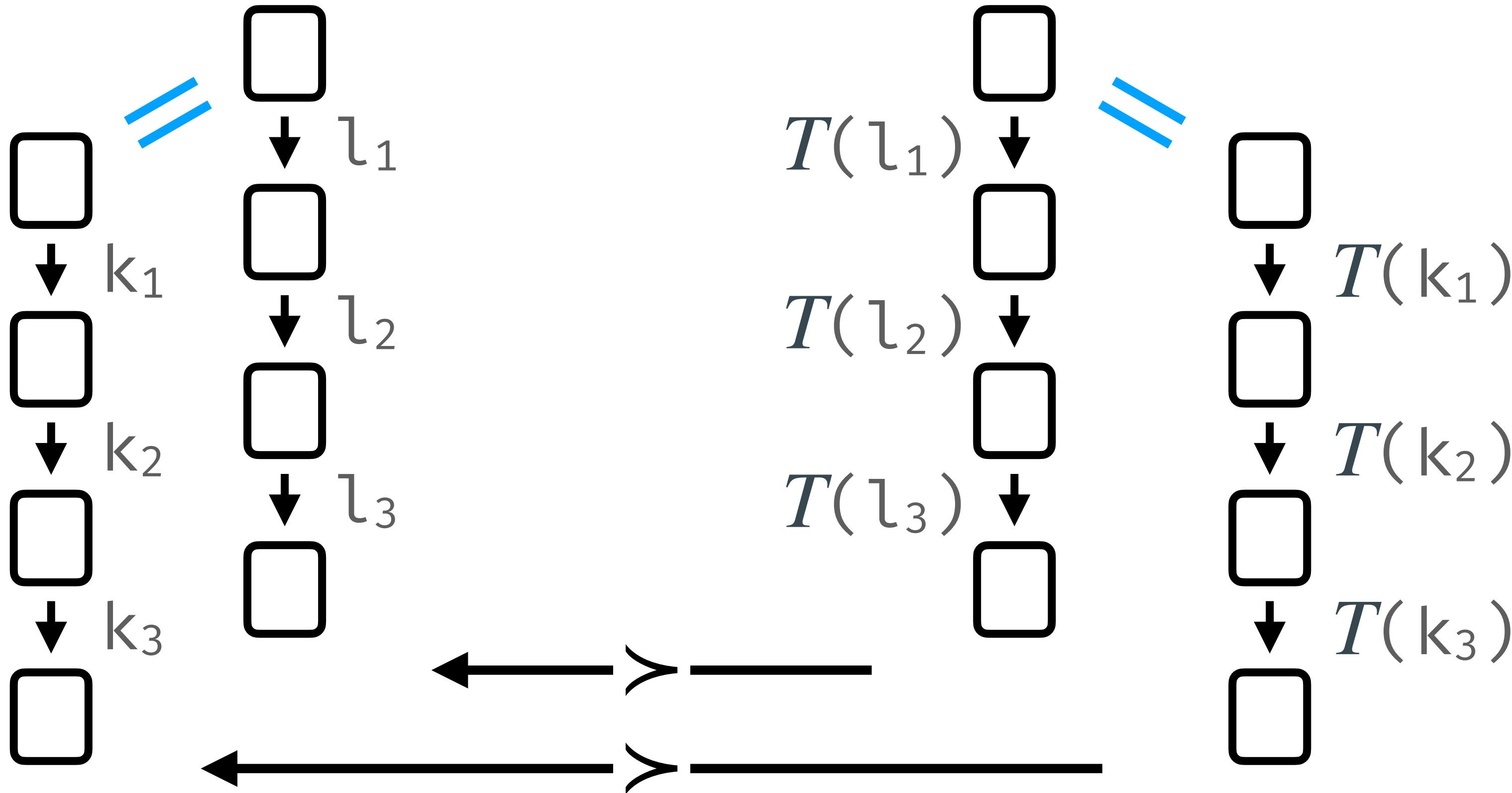
$\text{prog CT} \implies [\text{prog}] \text{CT}$

T : leakage
: transformer

Cites!

CT prog

[prog]



$$k_1 k_2 k_3 \dots = l_1 l_2 l_3 \dots$$

$$T(k_1 k_2 k_3 \dots) = T(l_1 l_2 l_3 \dots)$$

Proving "[] preserves CT"

prog CT

⇒

[prog] CT

T

• leakage

• transformer

Cites!

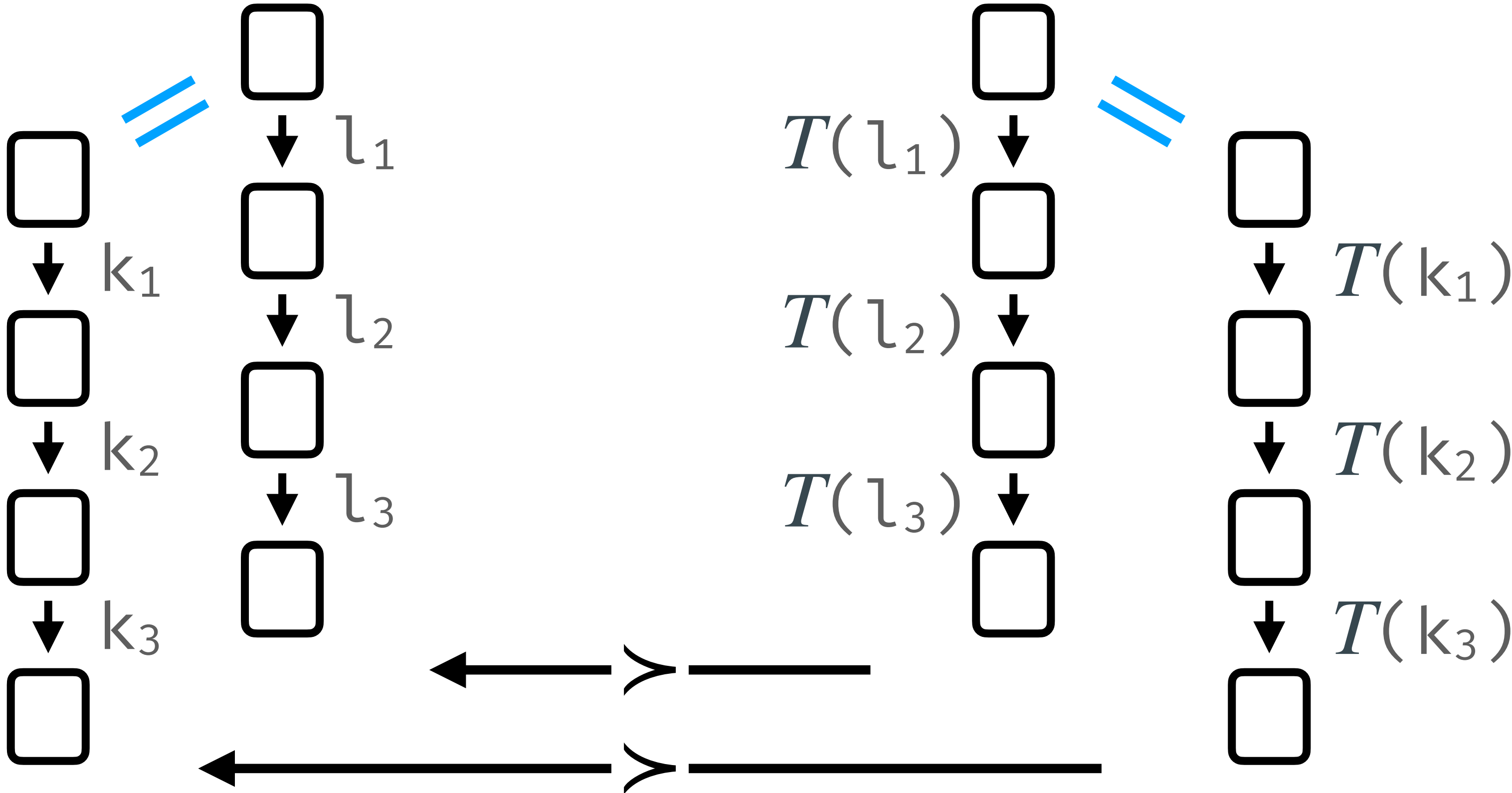
CT

prog

⇒

[prog]

CT



$k_1 k_2 k_3 \dots = l_1 l_2 l_3 \dots$

$T(k_1 k_2 k_3 \dots) = T(l_1 l_2 l_3 \dots)$

Proving “[] preserves CT”

$\text{prog CT} \implies [\text{prog}] \text{CT}$

T : Leakage
Transformer

$\text{CT prog} \implies [\text{prog}] \text{CT}$

THEOREM

[Barthe et al. '18, '21]

$\succ + T \implies [\] \text{ preserves CT.}$

$k_1 k_2 k_3 \dots = \iota_1 \iota_2 \iota_3 \dots$

$T(k_1 k_2 k_3 \dots) = T(\iota_1 \iota_2 \iota_3 \dots)$

$$\text{prog } \cancel{\text{cr}} \implies [\text{prog}] \cancel{\text{cr}}$$

Proving “[] reflects CT”

prog CT $\Leftarrow$ **[prog] CT**

Proving “[] reflects CT”

Proving “[] reflects CT”

$\text{prog CT} \Leftarrow [\text{prog}] \text{CT}$

prog

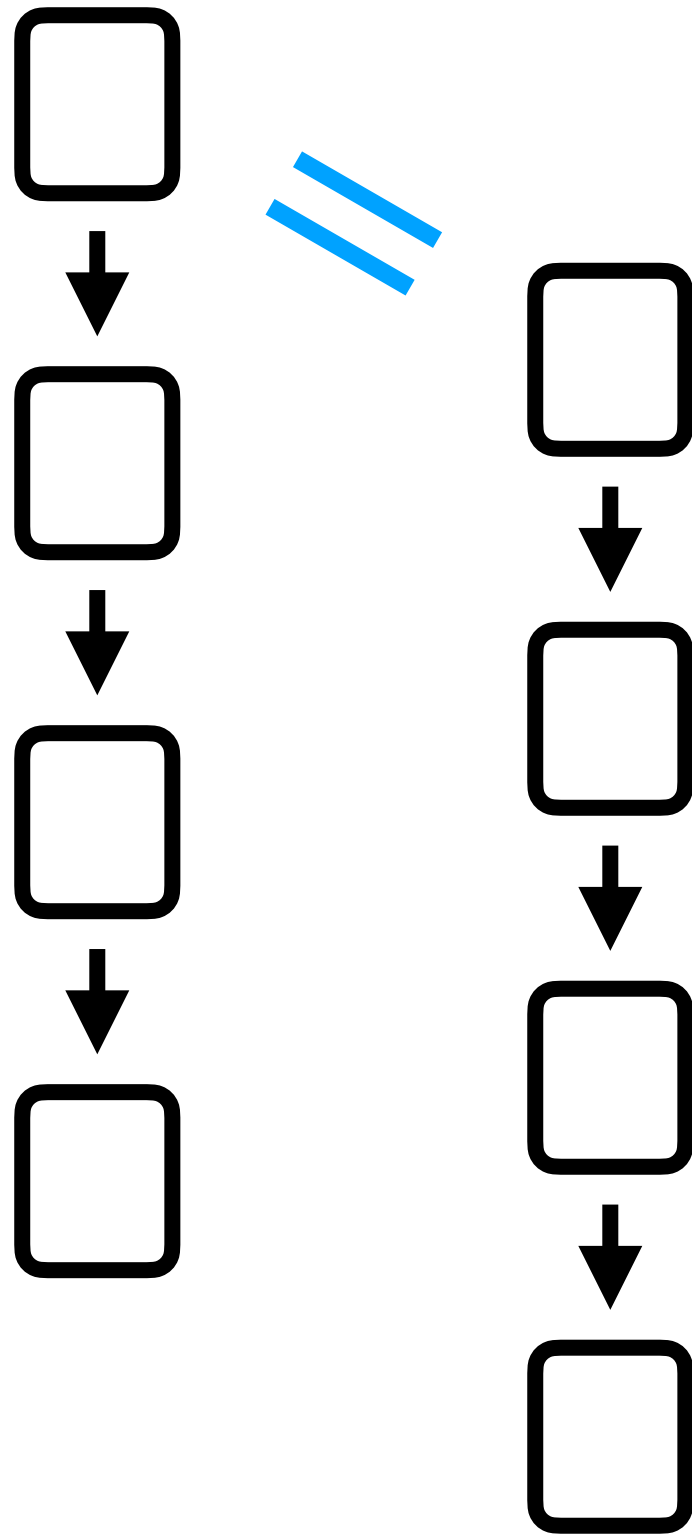
[prog]

Proving “[] reflects CT”

$\text{prog CT} \Leftarrow [\text{prog}] \text{CT}$

prog

[prog]

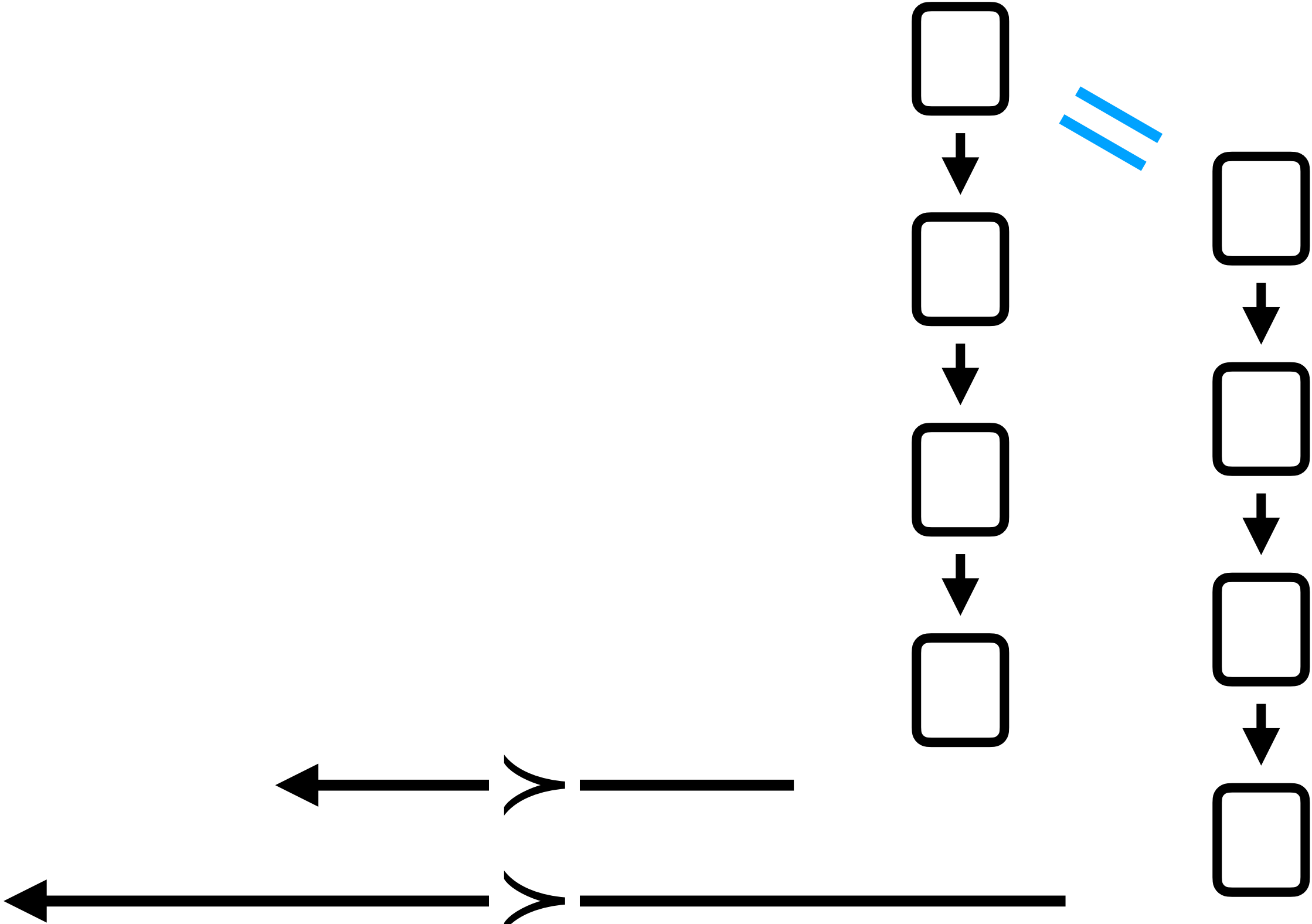


Proving "[] reflects CT"

$\text{prog CT} \Leftarrow [\text{prog}] \text{CT}$

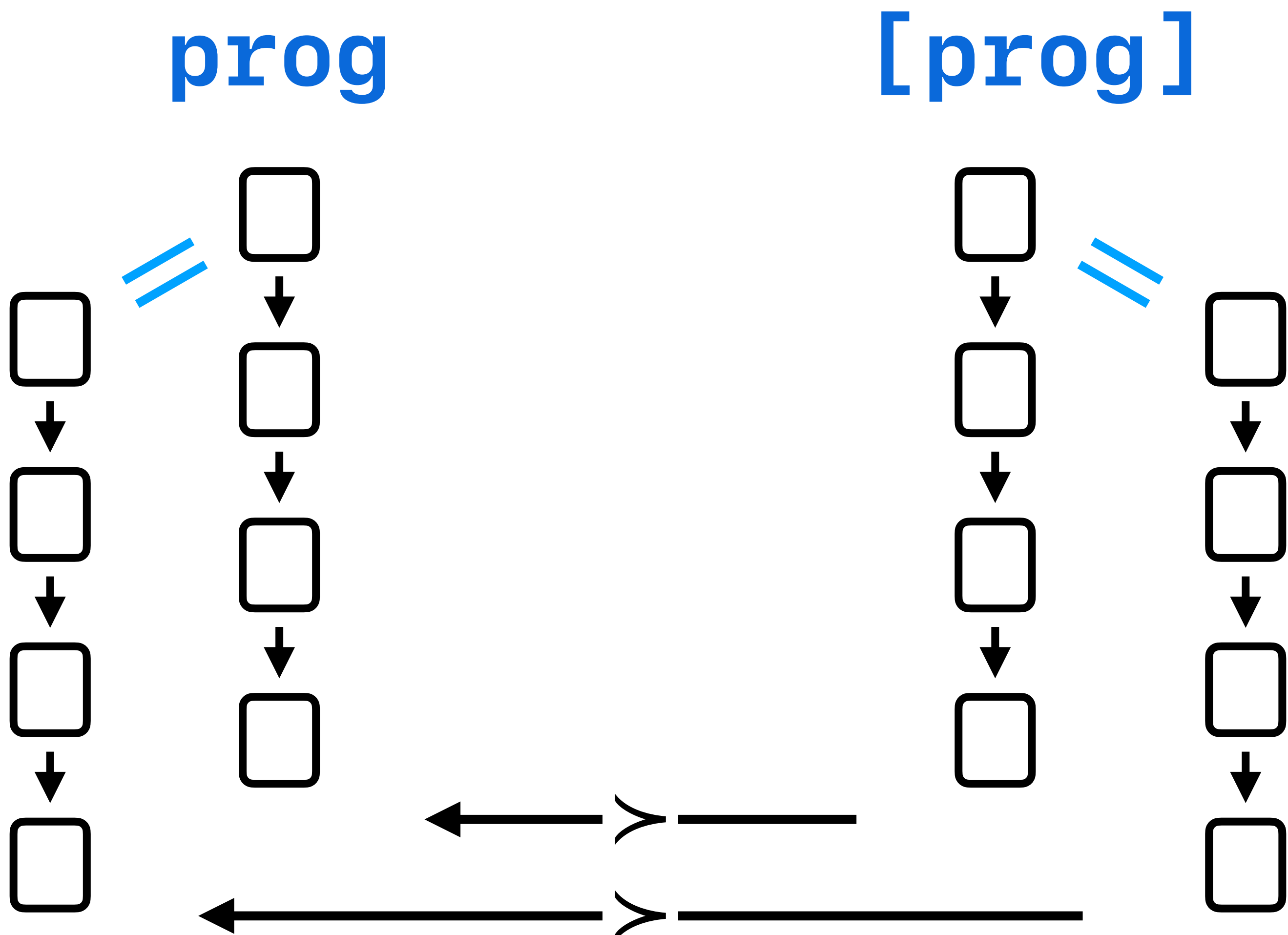
prog

[prog]



Proving "[] reflects CT"

$\text{prog CT} \Leftarrow [\text{prog}] \text{CT}$



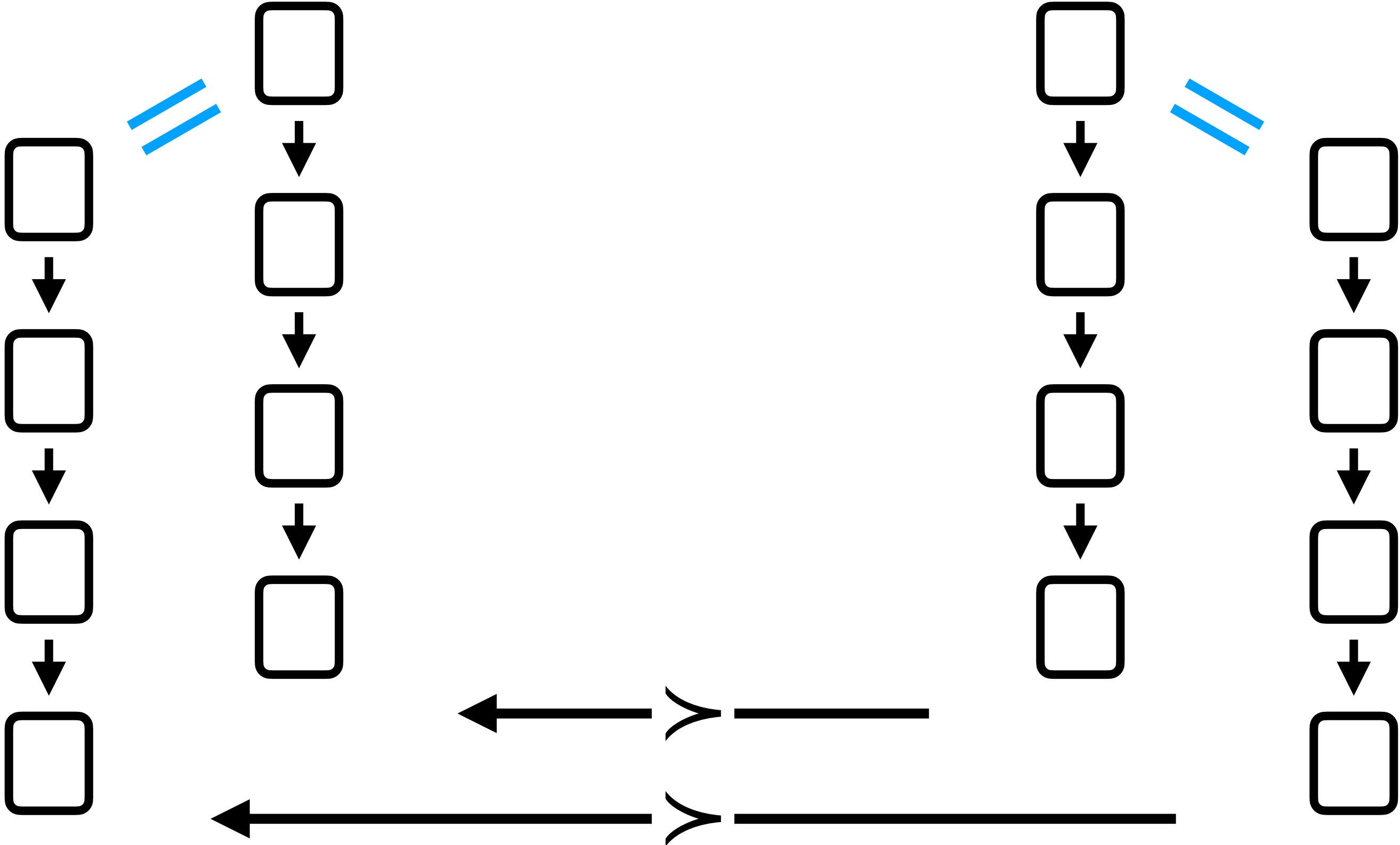
Proving "[] reflects CT"

$\text{prog CT} \Leftarrow [\text{prog}] \text{CT}$

U : Reverse
: Transformer

prog

[prog]



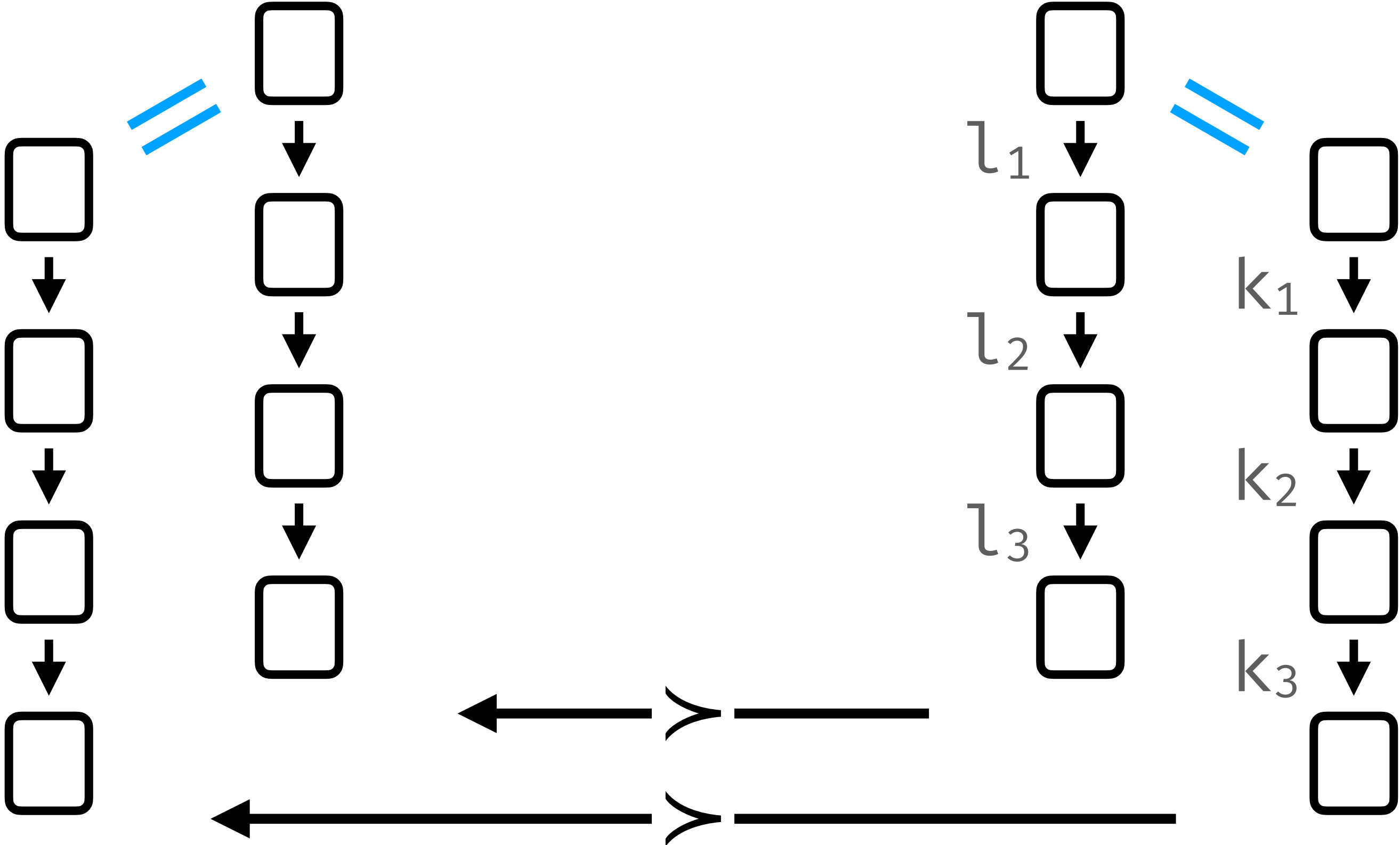
Proving "[] reflects CT"

$\text{prog CT} \Leftarrow [\text{prog}] \text{CT}$

U : Reverse
: Transformer

prog

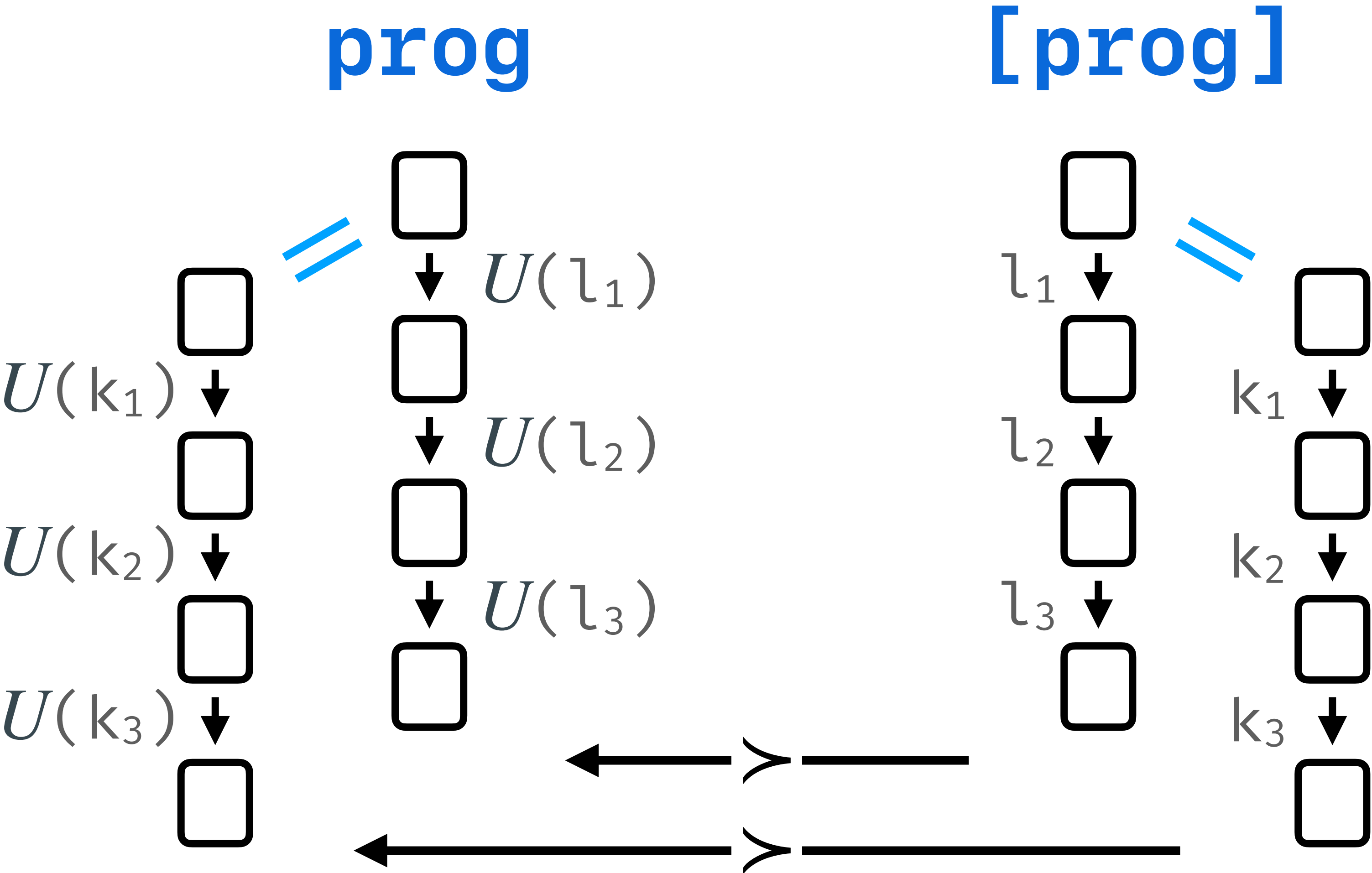
[prog]



Proving "[] reflects CT"

$\text{prog CT} \Leftarrow [\text{prog}] \text{CT}$

U : Reverse
Transformer



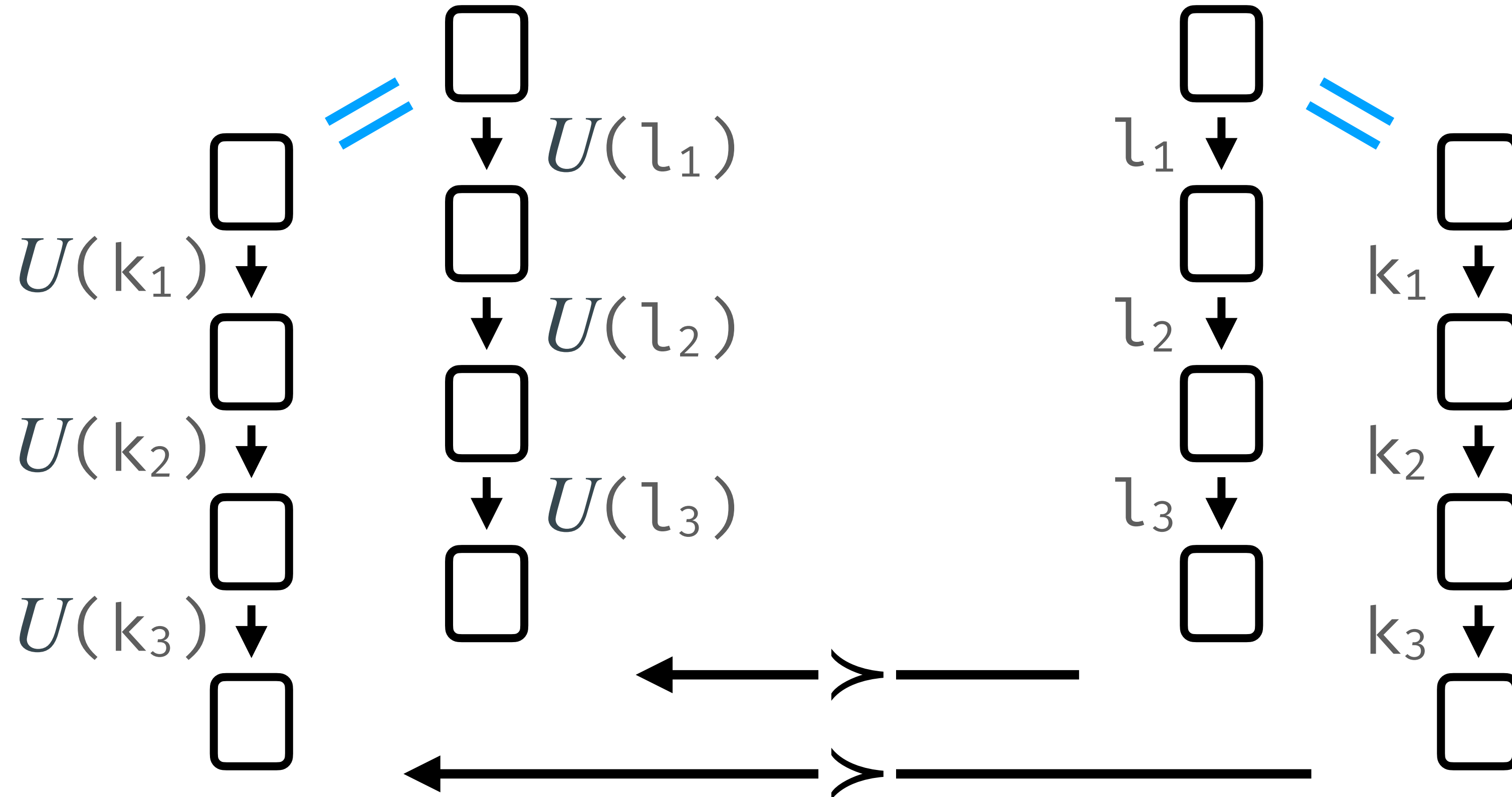
Proving “[] reflects CT”

prog CT $\Leftarrow$ **[prog] CT**

- Reverse
- Transformer

prog

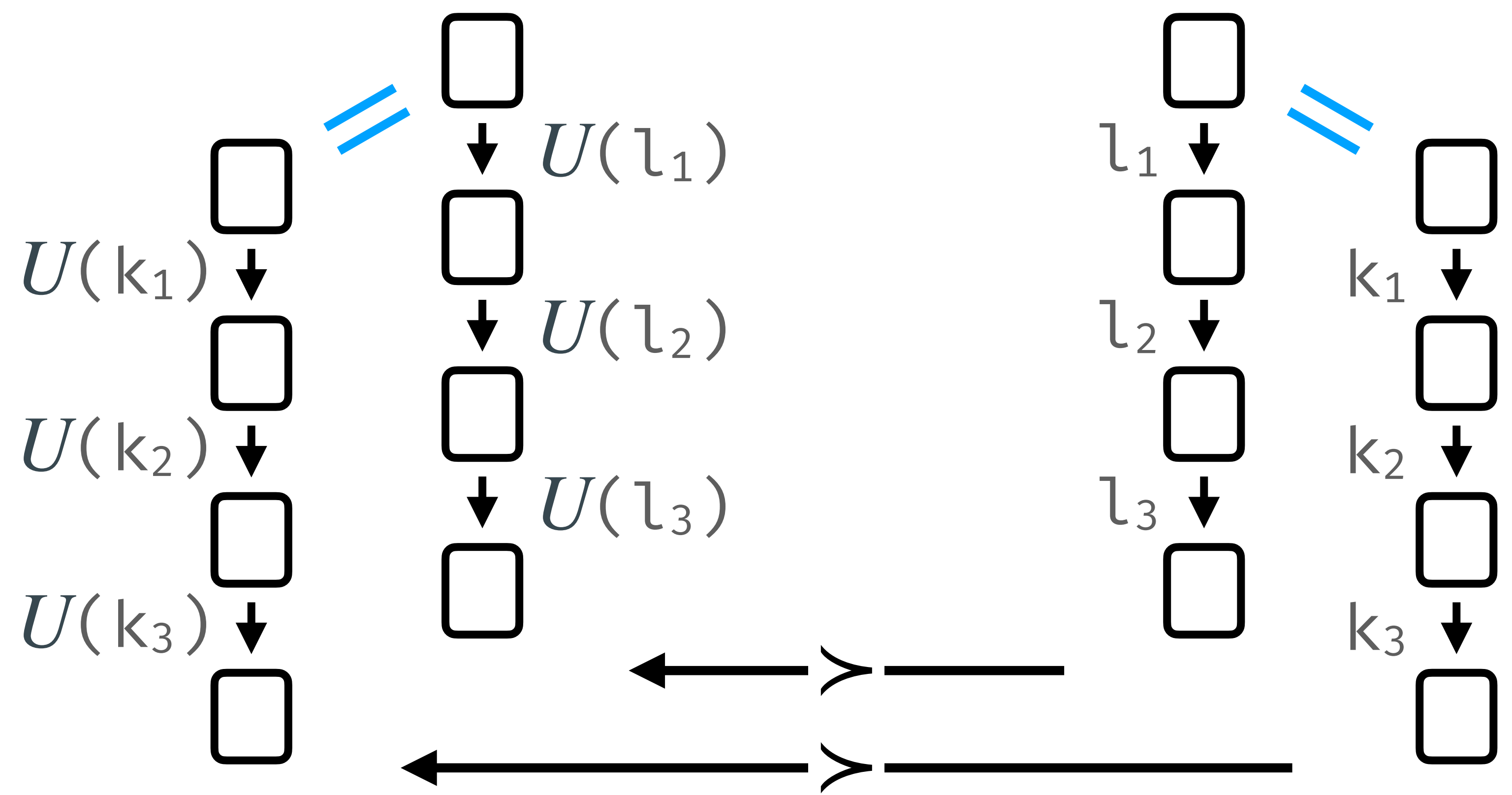
[prog] CT



$$k_1 k_2 k_3 \dots = l_1 l_2 l_3 \dots$$

U : Reverse
: Transformer

CT prog $\Leftarrow$ [prog] CT



$U(k_1 k_2 k_3 \dots) = U(l_1 l_2 l_3 \dots)$

$k_1 k_2 k_3 \dots = l_1 l_2 l_3 \dots$

Proving "[] reflects CT"

$\text{prog CT} \Leftarrow [\text{prog}] \text{CT}$

U : Reverse
Transformer

$\text{CT prog} \Leftarrow [\text{prog}] \text{CT}$

THEOREM

$\succ + \text{reverse } U \Rightarrow [] \text{ reflects CT.}$

$U(k_1k_2k_3...) = U(l_1l_2l_3...) \quad k_1k_2k_3... = l_1l_2l_3...$

Proving “[] is CT transparent”

Proving “[] is CT transparent”

$\text{prog CT} \iff [\text{prog}] \text{CT}$

[] preserves CT

+

[] reflects CT

Proving “**[]** is CT transparent”

$$\text{prog CT} \iff [\text{prog}] \text{CT}$$

[] preserves CT

+

[] reflects CT

transformer

reverse transformer

$$T : \begin{array}{l} \text{prog} \\ \text{leakage} \end{array} \longrightarrow \begin{array}{l} [\text{prog}] \\ \text{leakage} \end{array}$$

$$U : \begin{array}{l} [\text{prog}] \\ \text{leakage} \end{array} \longrightarrow \begin{array}{l} \text{prog} \\ \text{leakage} \end{array}$$

Proving “**[]** is CT transparent”

[] preserves CT

transformer

$$T : \begin{array}{c} \text{prog} \\ \bullet \text{ leakage} \end{array} \rightarrow \begin{array}{c} \text{[prog]} \\ \bullet \text{ leakage} \end{array}$$

exist for many **[]** already

+

[] reflects CT

reverse transformer

$$U : \begin{array}{c} \text{[prog]} \\ \bullet \text{ leakage} \end{array} \rightarrow \begin{array}{c} \text{prog} \\ \bullet \text{ leakage} \end{array}$$

$$\text{prog CT} \iff \text{[prog] CT}$$

Proving “**[]** is CT transparent”

[] preserves **CT**

transformer

$$T : \begin{array}{c} \text{prog} \\ \bullet \text{ leakage} \end{array} \rightarrow \begin{array}{c} \text{[prog]} \\ \bullet \text{ leakage} \end{array}$$

exist for many **[]** already

+

[] reflects **CT**

reverse transformer

$$U : \begin{array}{c} \text{[prog]} \\ \bullet \text{ leakage} \end{array} \rightarrow \begin{array}{c} \text{prog} \\ \bullet \text{ leakage} \end{array}$$

how?

$$\text{prog CT} \iff \text{[prog] CT}$$

Proving “[] is CT transparent”

$$\text{prog CT} \iff [\text{prog}] \text{CT}$$

[] preserves CT

+

[] reflects CT

transformer

$T : \text{prog leakage} \rightarrow [\text{prog}] \text{leakage}$

exist for many [] already

reverse transformer

$U : [\text{prog}] \text{leakage} \rightarrow \text{prog leakage}$

how?

Idea: Flip T to get U

Proving “[] is CT transparent”

$$\text{prog CT} \iff [\text{prog}] \text{CT}$$

[] preserves CT

+

[] reflects CT

transformer

$$T : \begin{array}{c} \text{prog} \\ \bullet \text{ leakage} \end{array} \rightarrow \begin{array}{c} [\text{prog}] \\ \bullet \text{ leakage} \end{array}$$

exist for many [] already

reverse transformer

$$T^{-1} : \begin{array}{c} [\text{prog}] \\ \bullet \text{ leakage} \end{array} \rightarrow \begin{array}{c} \text{prog} \\ \bullet \text{ leakage} \end{array}$$

how?

Idea: Flip T to get U

Proving “[] is CT transparent”

$$\text{prog CT} \iff [\text{prog}] \text{CT}$$

[] preserves CT

+

[] reflects CT

transformer

$$T : \begin{array}{l} \text{prog} \\ \bullet \text{ leakage} \end{array} \rightarrow \begin{array}{l} [\text{prog}] \\ \bullet \text{ leakage} \end{array}$$

exist for many [] already

reverse transformer

$$T^{-1} : \begin{array}{l} [\text{prog}] \\ \bullet \text{ leakage} \end{array} \rightarrow \begin{array}{l} \text{prog} \\ \bullet \text{ leakage} \end{array}$$

how?

Idea: Flip T to get U

Sufficient Condition

T must be **injective**!

Proving “[] is CT transparent”

$$\text{prog CT} \iff [\text{prog}] \text{CT}$$

[] preserves CT

+

[] reflects CT

transformer

$$T : \begin{array}{l} \text{prog} \\ \bullet \text{ leakage} \end{array} \rightarrow \begin{array}{l} [\text{prog}] \\ \bullet \text{ leakage} \end{array}$$

exist for many [] already

reverse transformer

$$T^{-1} : \begin{array}{l} [\text{prog}] \\ \bullet \text{ leakage} \end{array} \rightarrow \begin{array}{l} \text{prog} \\ \bullet \text{ leakage} \end{array}$$

how?

Idea: Flip T to get U

Sufficient Condition

T must be **injective!** ← + details

Proving “ $[\]$ is CT transparent”

$\text{prog CT} \iff [\text{prog}] \text{CT}$

THEOREM

$\succ + \text{injective } T \implies [\] \text{ is CT transparent.}$

Proving “[] is CT transparent”

$\text{prog CT} \iff [\text{prog}] \text{CT}$

THEOREM

$\succ + \text{injective } T \implies [\] \text{ is CT transparent.}$

[Loop Rotation]

[Untiling]

[Constant Folding]

[Unspilling]

[Structural Analysis]

[Dead Assignment Elimination]

[Dead Branch Elimination]

Making RETDEC transparent

Making RETDEC transparent

Making RETDEC transparent

Handcrafted
“likely to fail”
Snippets

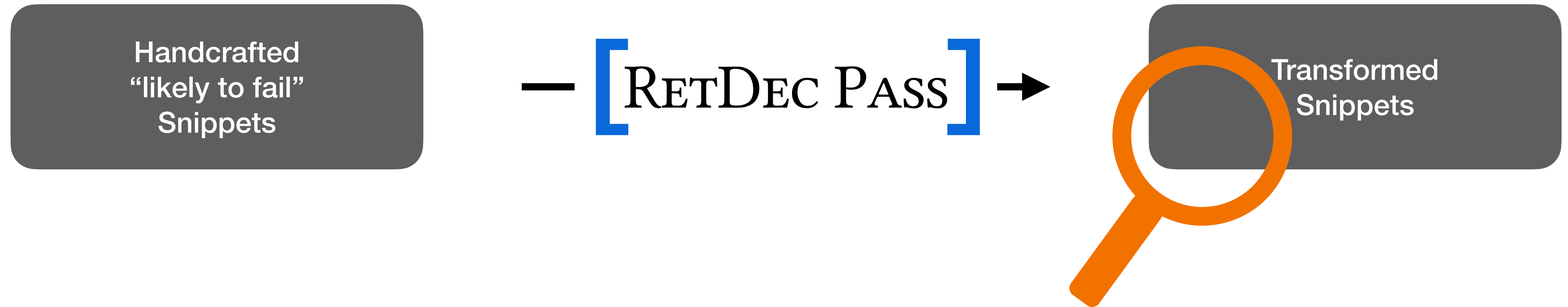
Making RETDEC transparent

Handcrafted
“likely to fail”
Snippets

— [RETDEC PASS] →

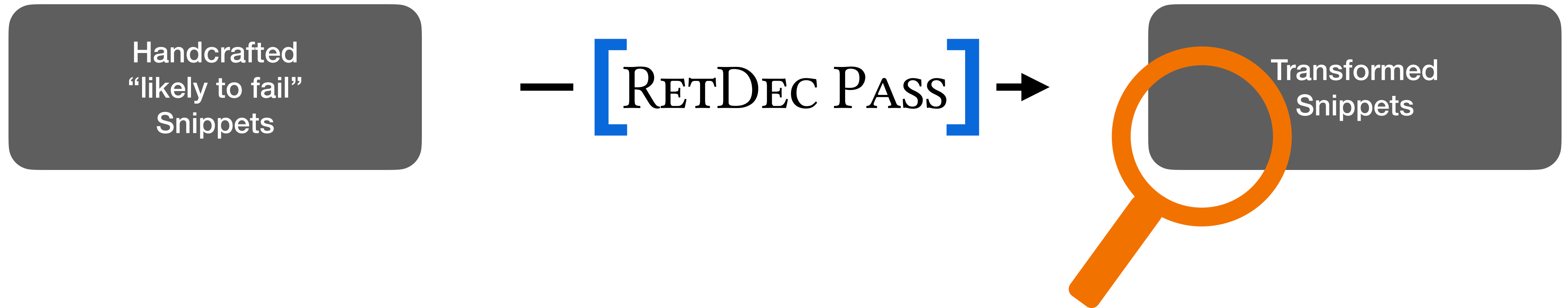
Transformed
Snippets

Making RETDEC transparent



Transparent on Snippets?

Making RETDEC transparent



Transparent on Snippets?

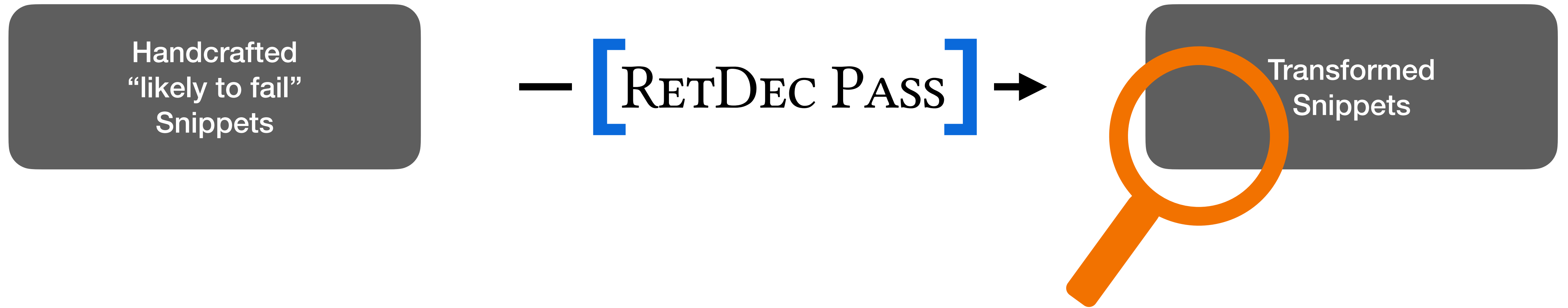


Keep pass (49)



Remove pass (13)

Making RETDEC transparent



Transparent on Snippets?



Keep pass (49)



Remove pass (13)



CT-RETDEC

Making RETDEC transparent

CT- RETDEC

Making RETDEC transparent

160 (10x16) Binaries
with and without
Real World Vulnerabilities

CT-RETDEC

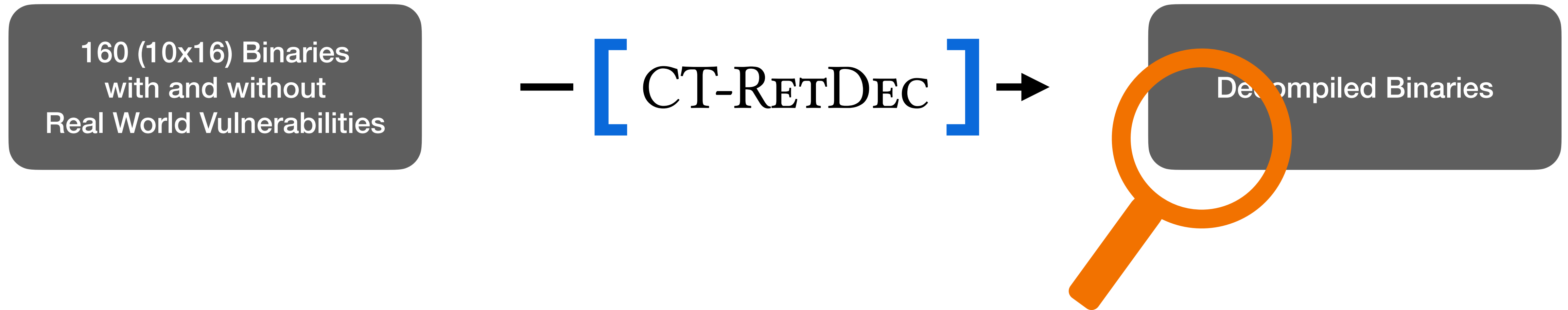
Making RETDEC transparent

160 (10x16) Binaries
with and without
Real World Vulnerabilities

— [CT-RETDEC] →

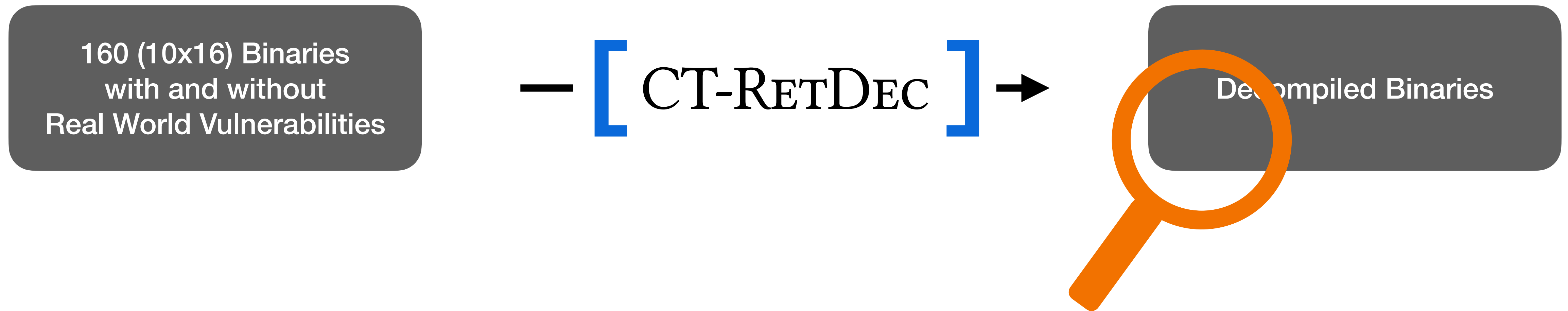
Decompiled Binaries

Making RETDEC transparent



transparent on all

Making RETDEC transparent



transparent on all

RETDEC removes vulnerabilities from 45

appearing in OOPSLA 26

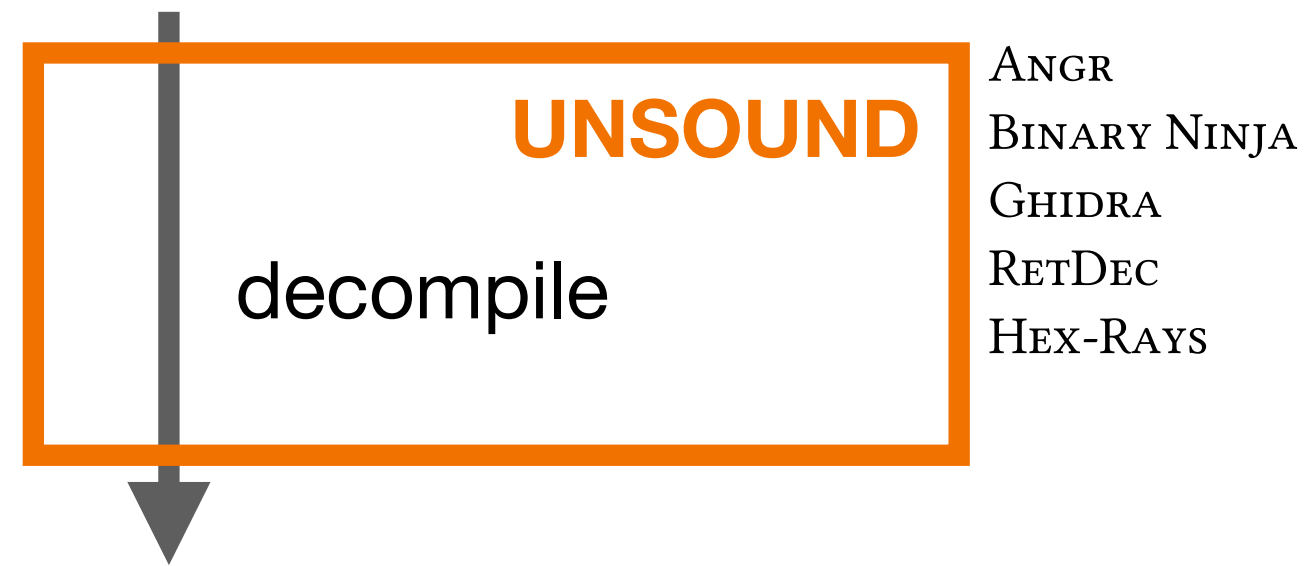
Decompiling for Constant Time Analysis

Santiago Arranz Olmos, Gilles Barthe, Lionel Blatter,
Youcef Bouzid, Sören van der Wall, Zhiyuan Zhang

s.van-der-wall@tu-bs.de



~~cr~~ crypto.o



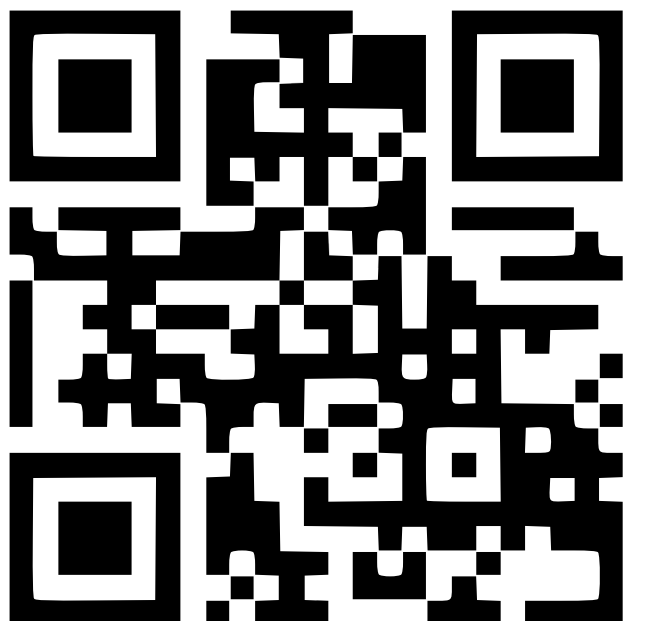
CT crypto.c

appearing in OOPSLA 26

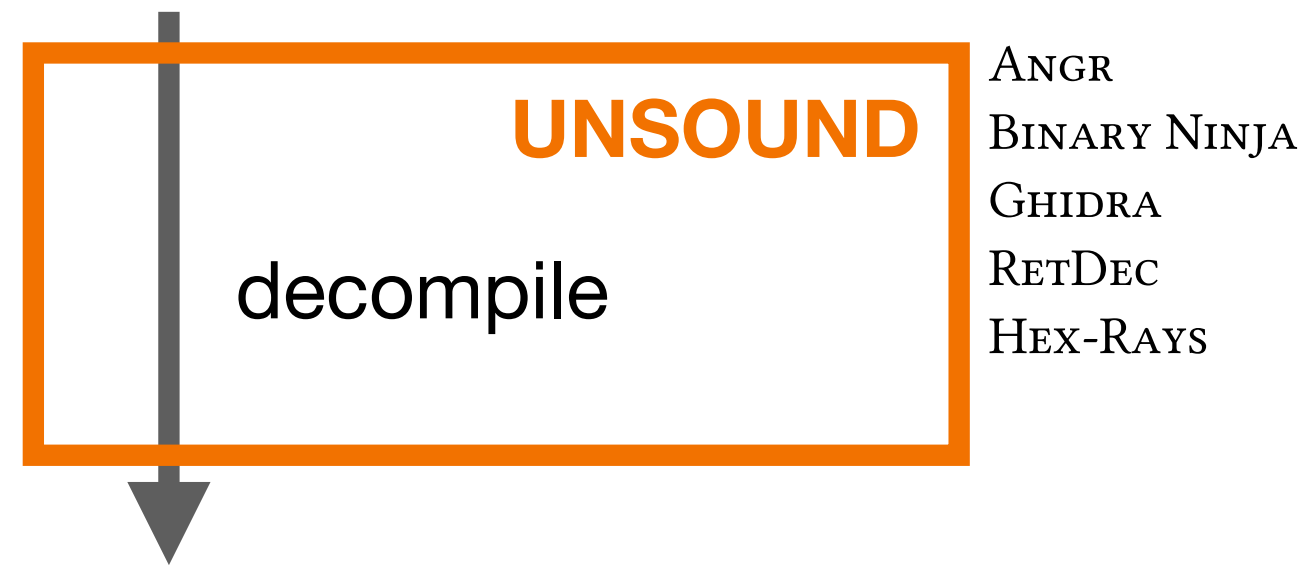
Decompiling for Constant Time Analysis

Santiago Arranz Olmos, Gilles Barthe, Lionel Blatter,
Youcef Bouzid, Sören van der Wall, Zhiyuan Zhang

s.van-der-wall@tu-bs.de



~~cr~~ crypto.o



CT crypto.c

[] CT transparent

WHEN

prog CT $\iff$ [prog] CT

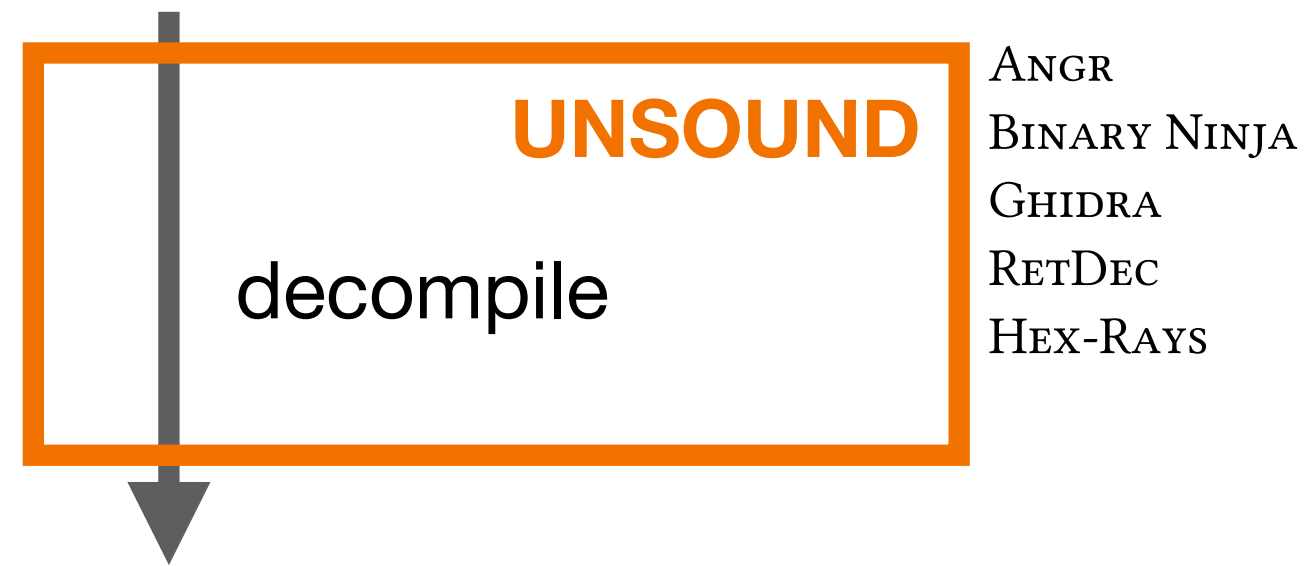
appearing in OOPSLA 26
Decompiling for Constant Time Analysis

Santiago Arranz Olmos, Gilles Barthe, Lionel Blatter,
Youcef Bouzid, Sören van der Wall, Zhiyuan Zhang

s.van-der-wall@tu-bs.de

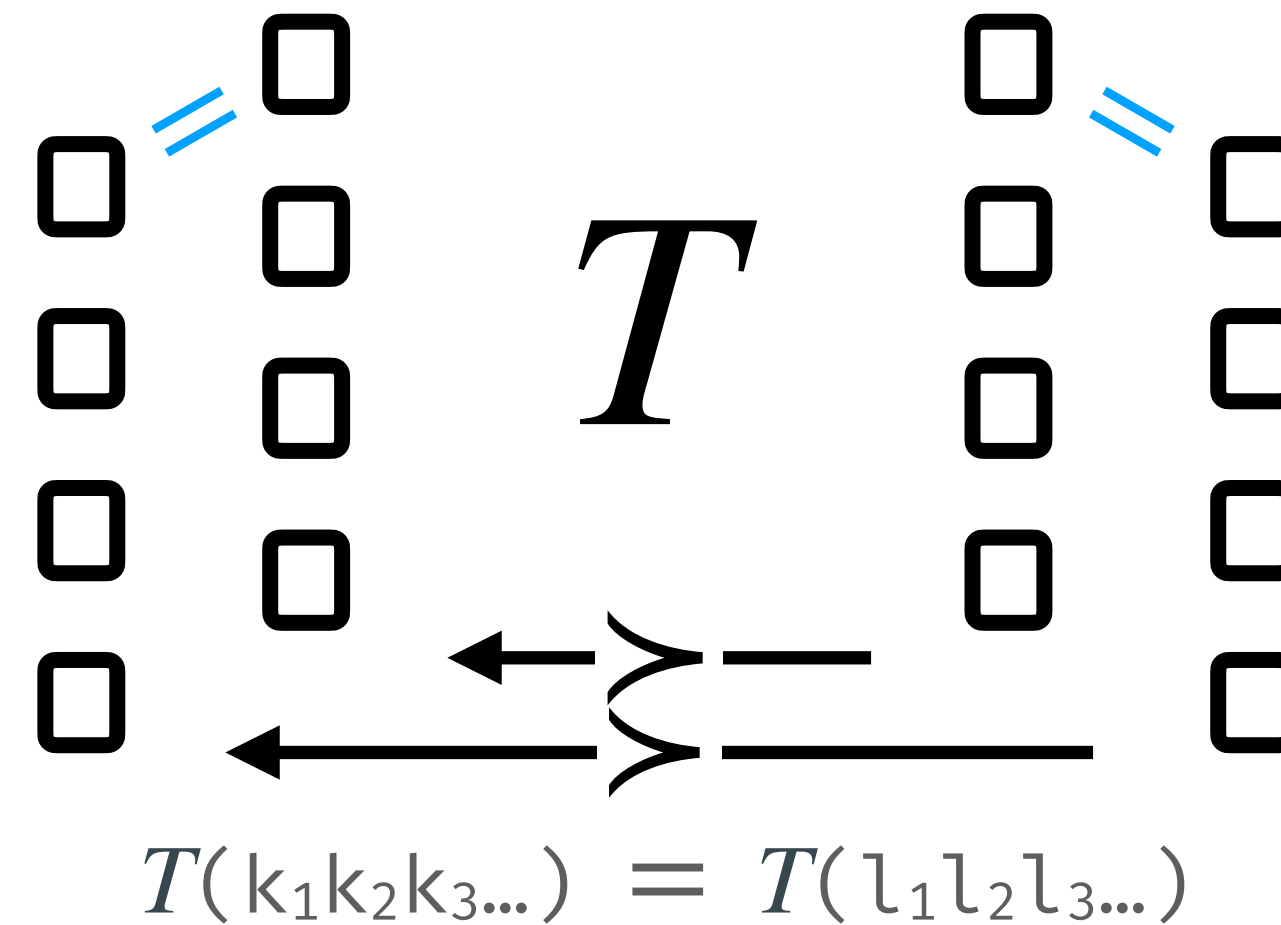


~~cr~~ crypto.o



CT crypto.c

Proving “[] preserves CT”



[] CT transparent

WHEN

prog CT $\iff$ [prog] CT

appearing in OOPSLA 26

Decompiling for Constant Time Analysis

Santiago Arranz Olmos, Gilles Barthe, Lionel Blatter,
Youcef Bouzid, Sören van der Wall, Zhiyuan Zhang

s.van-der-wall@tu-bs.de



~~cr~~ crypto.o



ANGR
BINARY NINJA
GHIDRA
RETDEC
HEX-RAYS

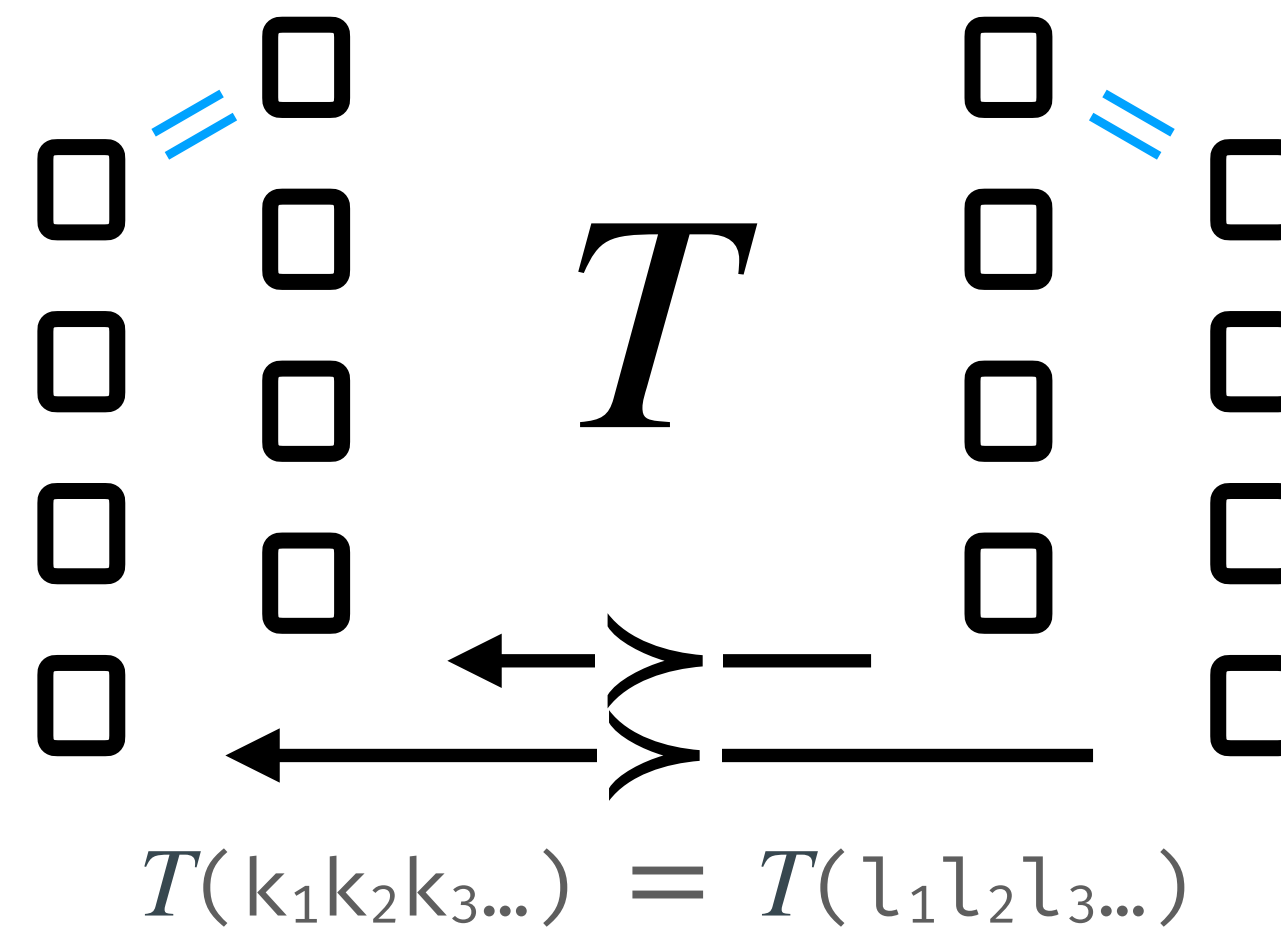
CT crypto.c

Proving “[] is CT transparent”

THEOREM

$\succ + \text{injective } T \implies [] \text{ is CT transparent.}$

Proving “[] preserves CT”



[] CT transparent

WHEN

prog CT $\iff$ [prog] CT

appearing in OOPSLA 26

Decompiling for Constant Time Analysis

Santiago Arranz Olmos, Gilles Barthe, Lionel Blatter,
Youcef Bouzid, Sören van der Wall, Zhiyuan Zhang

s.van-der-wall@tu-bs.de



~~cr~~ crypto.o



CT crypto.c

[] CT transparent

WHEN

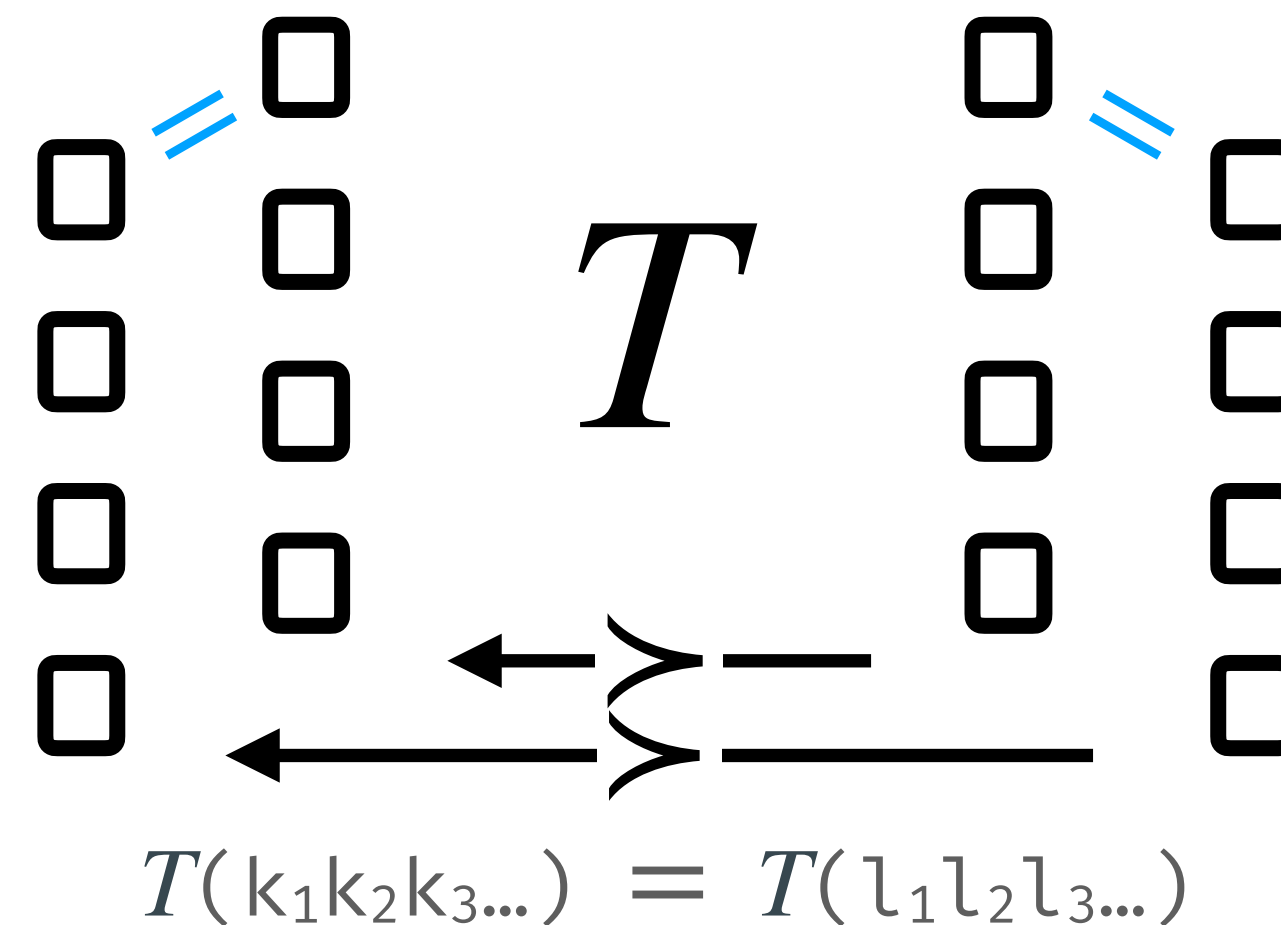
prog CT $\iff$ [prog] CT

Proving “[] is CT transparent”

THEOREM

$\succ +$ injective $T \implies$ [] is CT transparent.

Proving “[] preserves CT”



[CT-RETDEC]

appearing in OOPSLA 26

Decompiling for Constant Time Analysis

Santiago Arranz Olmos, Gilles Barthe, Lionel Blatter,
Youcef Bouzid, Sören van der Wall, Zhiyuan Zhang

s.van-der-wall@tu-bs.de

